

ORF522 – Linear and Nonlinear Optimization

17. Operator theory I

Today's lecture

[Chapter 4, FMO][PA][PMO][LSMO]

Operator theory I

- Proximal gradient method
- Operators
- Monotone and cocoercive operators
- Fixed-point Iterations

Proximal gradient method

Remember: gradient descent interpretation

Problem

$$\text{minimize } f(x)$$

Iterations

$$x^{k+1} = x^k - t \nabla f(x^k)$$

Quadratic approximation, replacing Hessian $\nabla^2 f(x^k)$ with $\frac{1}{t}I$

$$x^{k+1} = \underset{z}{\operatorname{argmin}} f(x^k) + \nabla f(x^k)^T (z - x^k) + \frac{1}{2t} \|z - x^k\|_2^2$$

Let's exploit the smooth part

$$\text{minimize } f(x) + g(x)$$

$f(x)$ convex and smooth

$g(x)$ convex (may be not differentiable)

Quadratic approximation of f while keeping g

$$x^{k+1} = \underset{z}{\operatorname{argmin}} g(z) + f(x^k) + \nabla f(x^k)^T (z - x^k) + \frac{1}{2t} \|z - x^k\|_2^2 \quad \leftarrow \text{same as gradient descent}$$

Equivalent to

Proximal operator

$$x^{k+1} = \underset{z}{\operatorname{argmin}} tg(z) + \frac{1}{2} \|z - (x^k - t\nabla f(x^k))\|_2^2 = \mathbf{prox}_{tg} (x^k - t\nabla f(x^k))$$

↑
make g
small

↑
stay close to
gradient update

Proximal gradient method

minimize $f(x) + g(x)$

$f(x)$ convex and smooth

$g(x)$ convex (may be not differentiable)

Iterations

$$x^{k+1} = \text{prox}_{tg} \left(x^k - t \nabla f(x^k) \right)$$

Properties

- Alternates between gradient updates of f and proximal updates on g
- Useful if prox_{tg} is inexpensive
- Can handle nonsmooth and constrained problems

Special cases

Generalized gradient descent

Smooth

$$g(x) = 0 \implies \text{prox}_{tg}(x) = x$$

Constraints

$$g(x) = \mathcal{I}_C(x) \implies \text{prox}_{tg}(x) = \Pi_C(x)$$

Non smooth

$$f(x) = 0$$

Problem

$$\text{minimize } f(x) + g(x)$$

Iterations

$$x^{k+1} = \text{prox}_{tg}(x^k - t\nabla f(x^k))$$

Gradient descent

$$\implies x^{k+1} = x^k - t\nabla f(x^k)$$

Projected gradient descent

$$\implies x^{k+1} = \Pi_C(x^k - t\nabla f(x^k))$$

Proximal minimization

$$\implies x^{k+1} = \text{prox}_{tg}(x^k)$$

Note: useful if prox_{tg} is cheap ⁷

What happens if we cannot evaluate the prox?

At every iteration, it can be very expensive to evaluate

$$\mathbf{prox}_g(x) = \operatorname{argmin}_z \left(g(z) + \frac{1}{2} \|z - x\|_2^2 \right)$$

Idea: solve it approximately!

If you precisely control the $\mathbf{prox}_g(x)$ evaluation errors
you can obtain the same convergence guarantees (and rates)
as the exact evaluations.

Example: Lasso

Iterative Soft Thresholding Algorithm (ISTA)

$$\text{minimize} \quad \underbrace{(1/2) \|Ax - b\|_2^2}_{f(x)} + \underbrace{\lambda \|x\|_1}_{g(x)}$$

Proximal gradient descent

$$x^{k+1} = \text{prox}_{tg} \left(x^k - t \nabla f(x^k) \right)$$

$$\nabla f(x) = A^T (Ax - b)$$

$$\text{prox}_{tg}(x) = S_{\lambda t}(x) \quad (\text{component wise soft-thresholding})$$

Closed-form iterations

$$x^{k+1} = S_{\lambda t} \left(x^k - t A^T (Ax^k - b) \right)$$

Example: Lasso

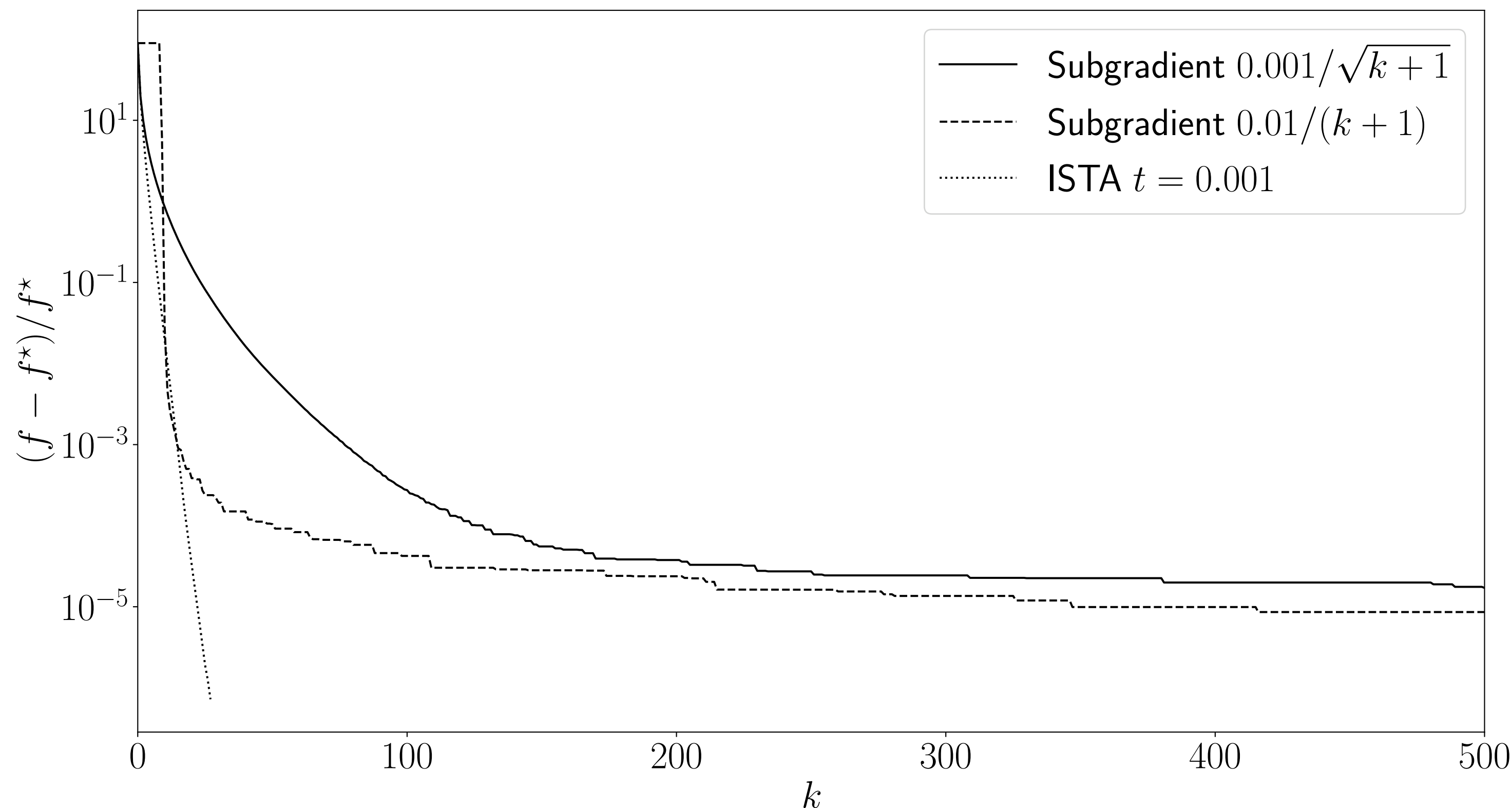
Iterative Soft Thresholding Algorithm (ISTA)

$$A \in \mathbf{R}^{500 \times 100}$$

$$\text{minimize} \quad (1/2)\|Ax - b\|_2^2 + \lambda\|x\|_1$$

Closed-form iterations

$$x^{k+1} = S_{\lambda t} \left(x^k - tA^T(Ax^k - b) \right)$$



Better convergence

Can we prove convergence generally?

Can we combine different operators?

Operators

Operators

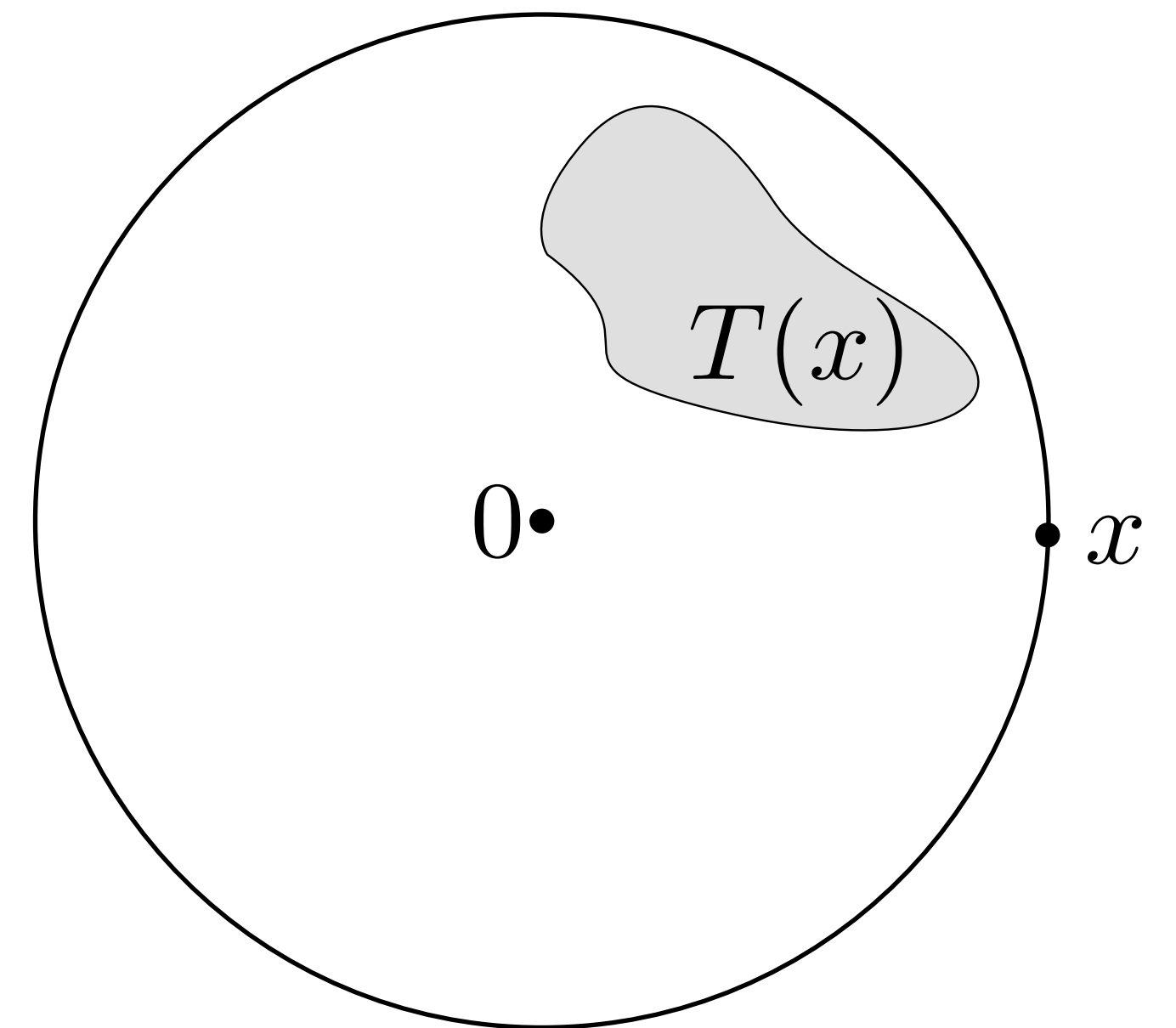
An operator T maps each point in \mathbf{R}^n to a subset of \mathbf{R}^n

- **set valued** $T(x)$ returns a set
- **single-valued** $T(x)$ (function) returns a singleton

The **domain** of T is the set $\text{dom } T = \{x \mid T(x) \neq \emptyset\}$

Example

- The subdifferential ∂f is a set-valued operator
- The gradient ∇f is a single-valued operator



Graph and inverse operators

Graph

The graph of an operator T is defined as

$$\mathbf{gph}T = \{(x, y) \mid y \in T(x)\}$$

In other words, all the pairs of points (x, y) such that $y \in T(x)$.

Inverse

The graph of the inverse operator T^{-1} is defined as

$$\mathbf{gph}T^{-1} = \{(y, x) \mid (x, y) \in \mathbf{gph}T\}$$

Therefore, $y \in T(x)$ if and only if $x \in T^{-1}(y)$.

Zeros

Zero

x is a **zero** of T if $0 \in T(x)$

Zero set

The set of all the zeros $T^{-1}(0) = \{x \mid 0 \in T(x)\}$

Example

If $T = \partial f$ and $f : \mathbf{R}^n \rightarrow \mathbf{R}$, then
 $0 \in T(x)$ means that x minimizes f

Many problems
can be posed as finding zeros
of an operator

Fixed points

\bar{x} is a **fixed-point** of a single-valued operator T if

$$\bar{x} = T(\bar{x})$$

Set of fixed points $\text{fix } T = \{x \in \text{dom } T \mid x = T(x)\} = (I - T)^{-1}(0)$

Examples

- **Identity** $T(x) = x$. Any point is a fixed point
- **Zero operator** $T(x) = 0$. Only 0 is a fixed point

Lipschitz operators

An operator T is L -Lipschitz if

$$\|T(x) - T(y)\| \leq L\|x - y\|, \quad \forall x, y \in \mathbf{dom} T$$

Fact If T is Lipschitz, then it is single-valued

Proof If $y = T(x)$, $z = T(x)$, then $\|y - z\| \leq L\|x - x\| = 0 \implies y = z$ ■

For $L = 1$ we say T is **nonexpansive**

For $L < 1$ we say T is **contractive** (with contraction factor L)

Lipschitz operators examples

Lipschitz affine functions

$$T(x) = Ax + b$$



maximum singular value

$$L = \|A\|_2 = \sqrt{\lambda_{\max}(A^T A)}$$

Lipschitz differentiable functions

T such that there exists derivative DT



derivative is bounded

$$\|DT\|_2 \leq L$$

Lipschitz operators and fixed points

Given a L -Lipschitz operator T and a fixed point $\bar{x} = T\bar{x}$,

$$\|Tx - \bar{x}\| = \|Tx - T\bar{x}\| \leq L\|x - \bar{x}\|$$

A contractive operator ($L < 1$) can have at most one fixed point, i.e., $\text{fix } T = \{\bar{x}\}$

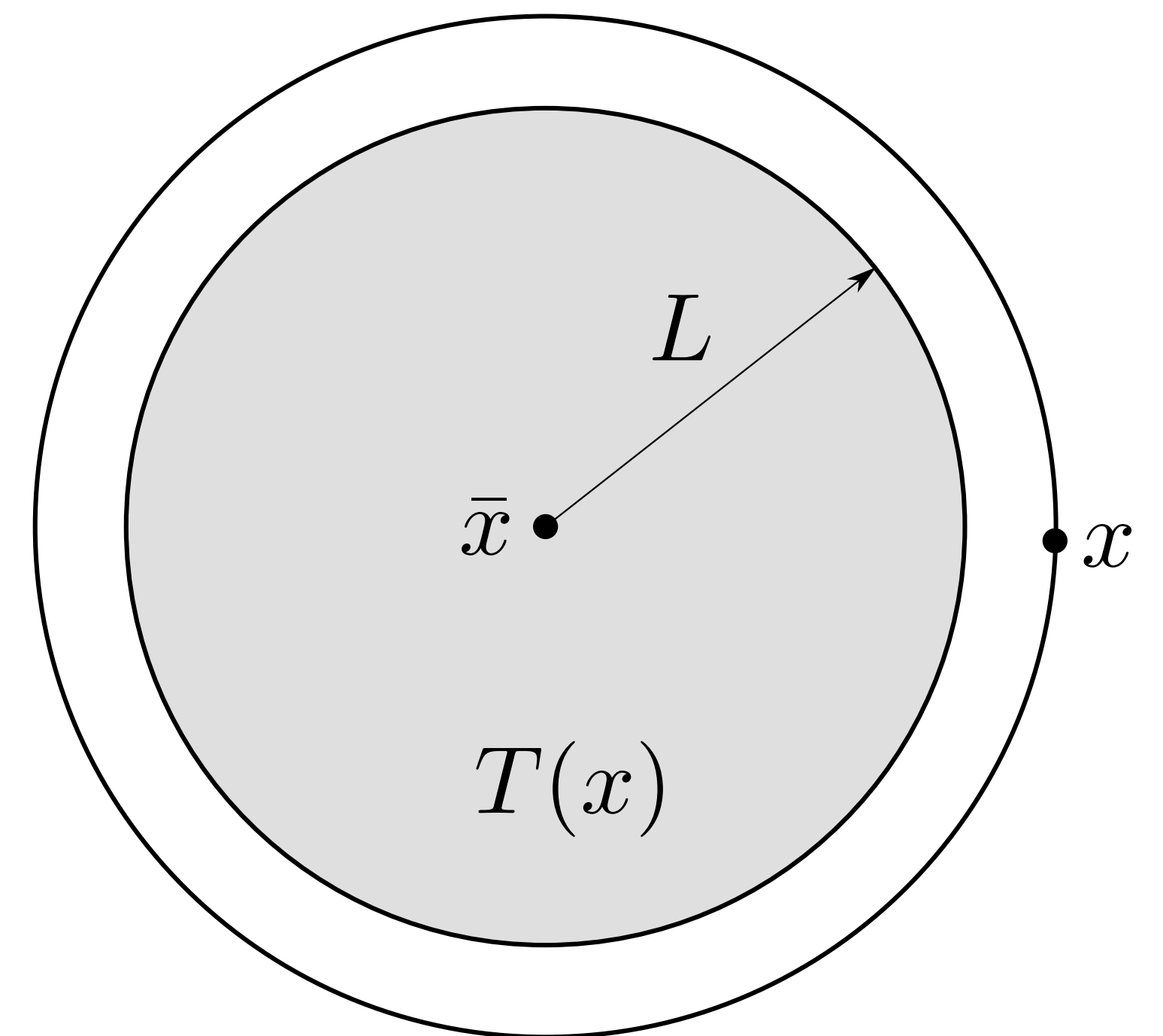
Proof

If $\bar{x}, \bar{y} \in \text{fix } T$ and $\bar{x} \neq \bar{y}$ then

$$\|\bar{x} - \bar{y}\| = \|T(\bar{x}) - T(\bar{y})\| < \|\bar{x} - \bar{y}\| \text{ (contradiction)} \blacksquare$$

A nonexpansive operator ($L = 1$) need not have a fixed point

Example $T(x) = x + 2$



Combining Lipschitz operators

T_1 is L_1 -Lipschitz and T_2 is L_2 -Lipschitz

The **composition** T_1T_2 is L_1L_2 -Lipschitz

Proof $\|T_1T_2x - T_1T_2y\|_2 \leq L_1\|T_2x - T_2y\|_2 \leq L_1L_2\|x - y\|_2$ ■

- Composition of *nonexpansive* is nonexpansive
- Composition of *nonexpansive* and *contractive* is contractive

The **weighted average** $\theta T_1 + (1 - \theta)T_2$, $\theta \in (0, 1)$ is $(\theta L_1 + (1 - \theta)L_2)$ -Lipschitz

Proof (exercise)

- Weighted average of *nonexpansive* is nonexpansive
- Weighted average of *nonexpansive* and *contractive* is contractive

Monotone cocoercive operators

Monotone operators

An operator T on \mathbf{R}^n is **monotone** if

$$(u - v)^T (x - y) \geq 0, \quad \forall (x, u), (y, v) \in \text{gph} T$$

T is **maximal monotone** if

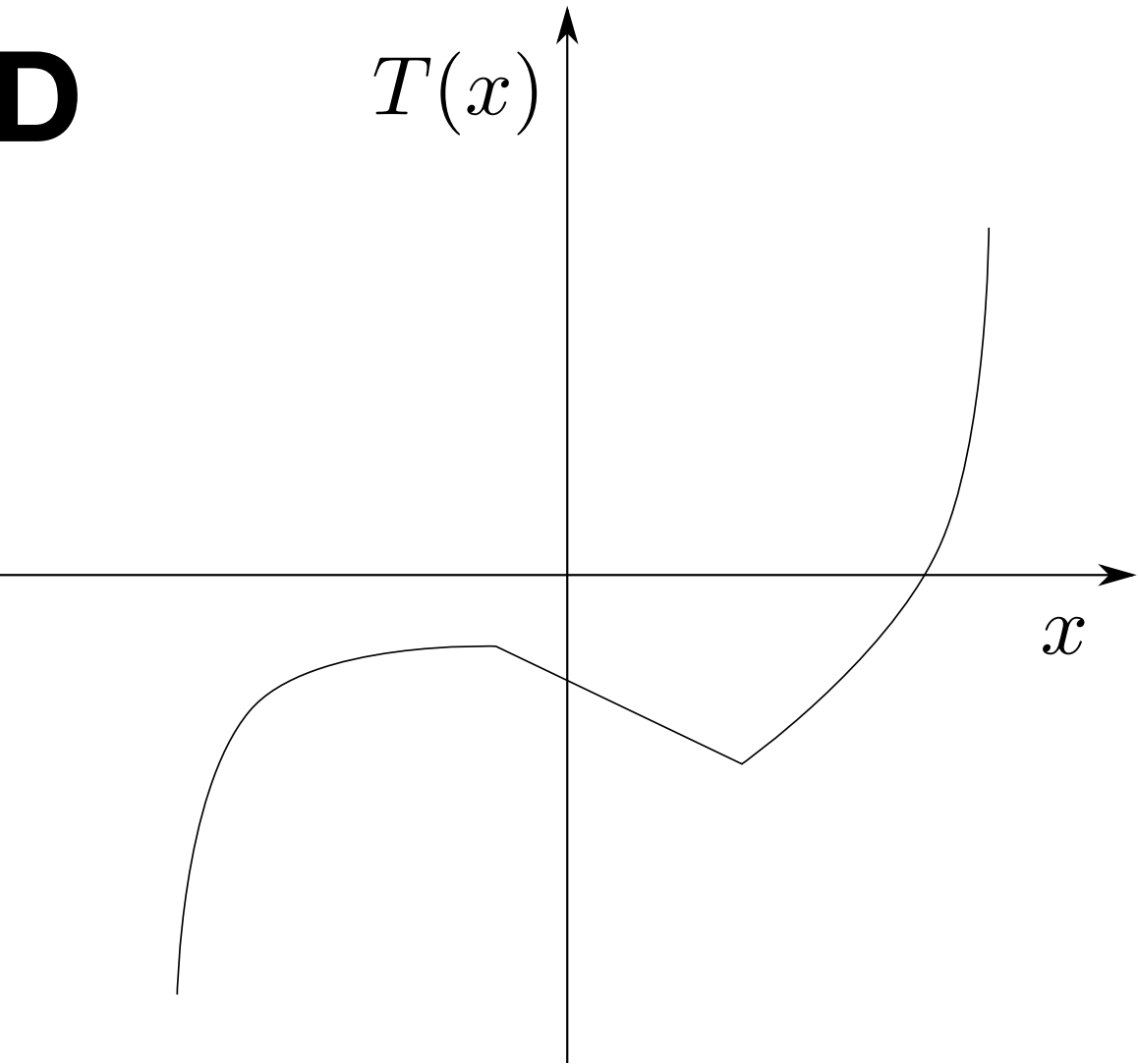
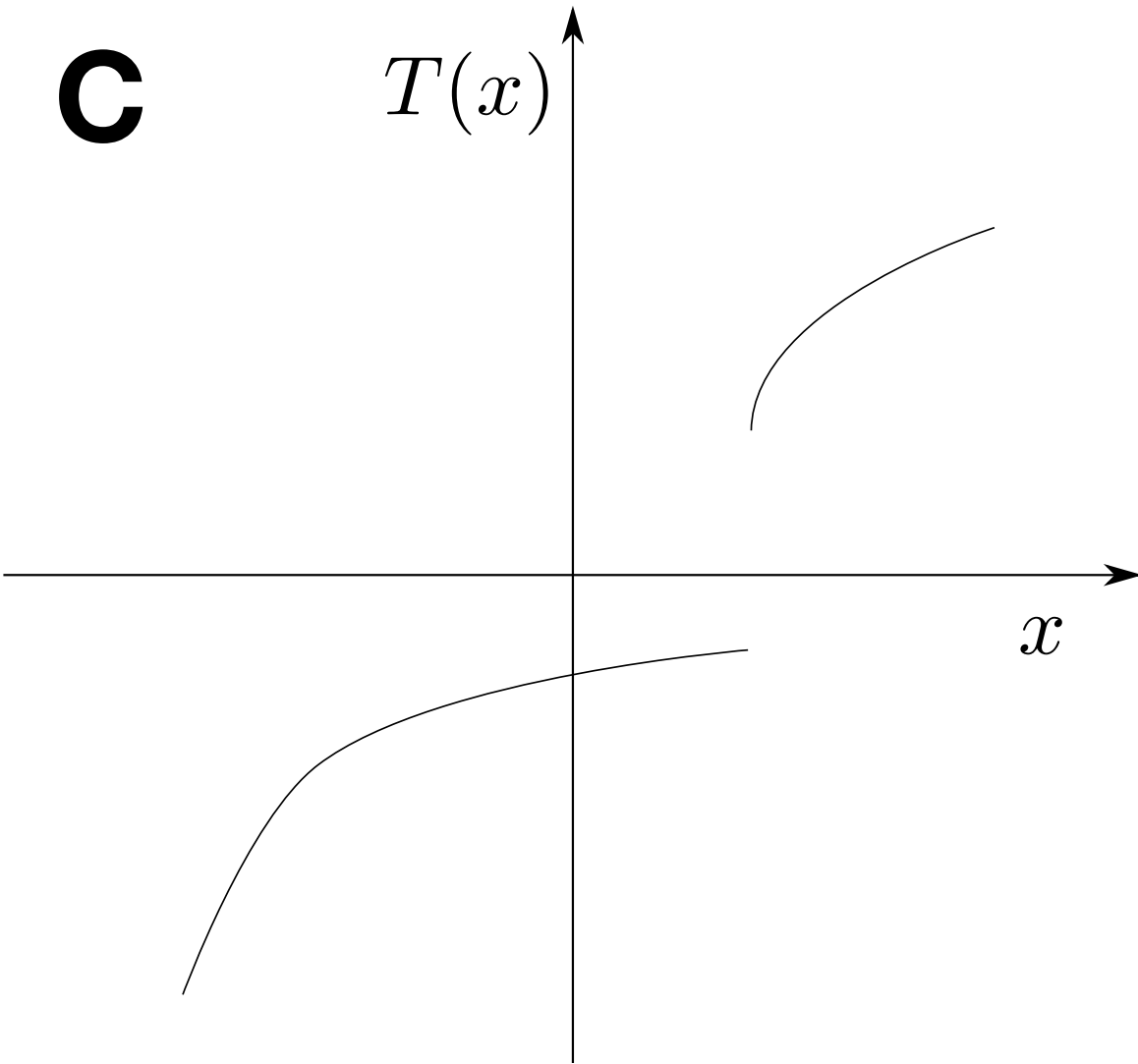
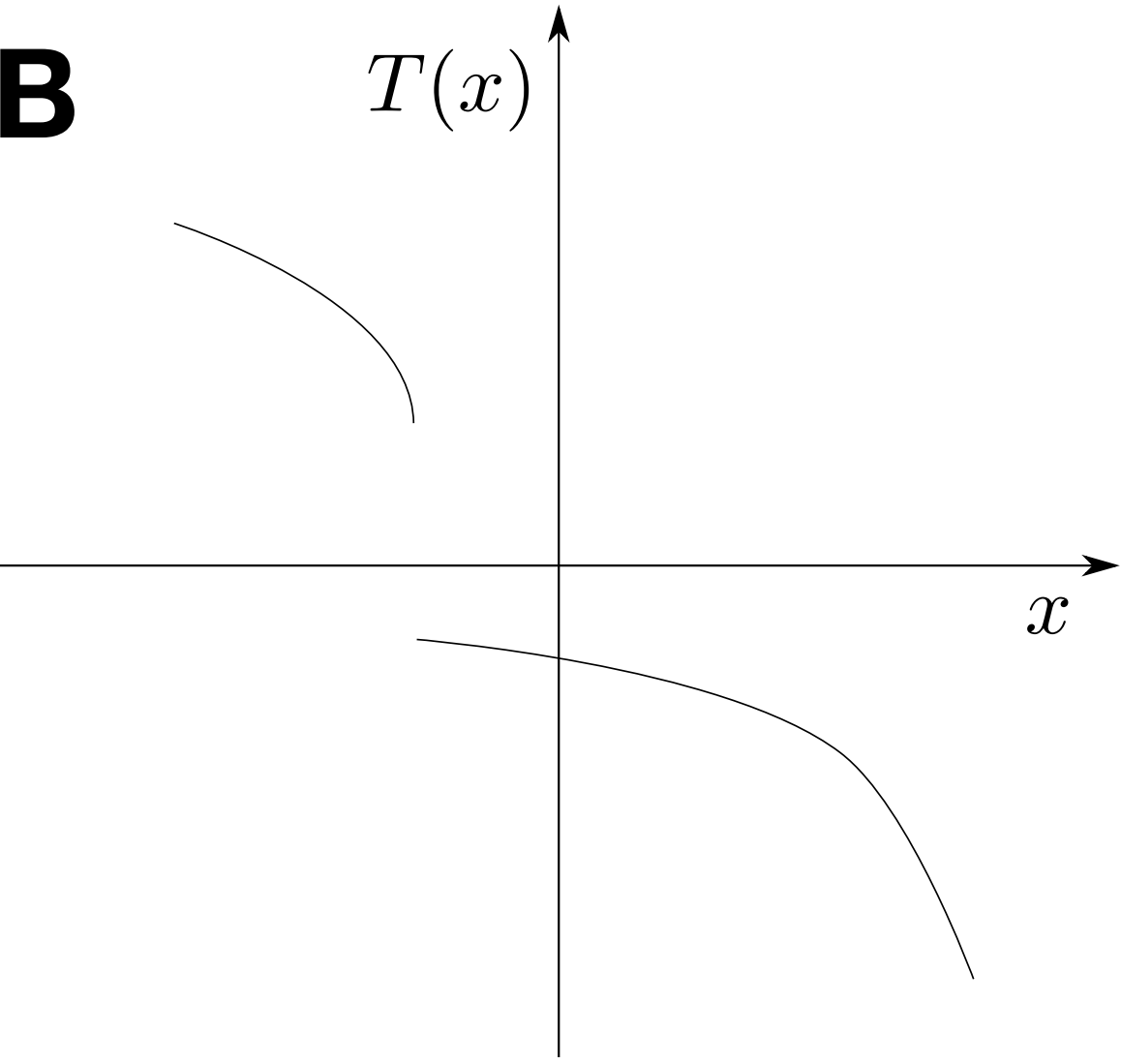
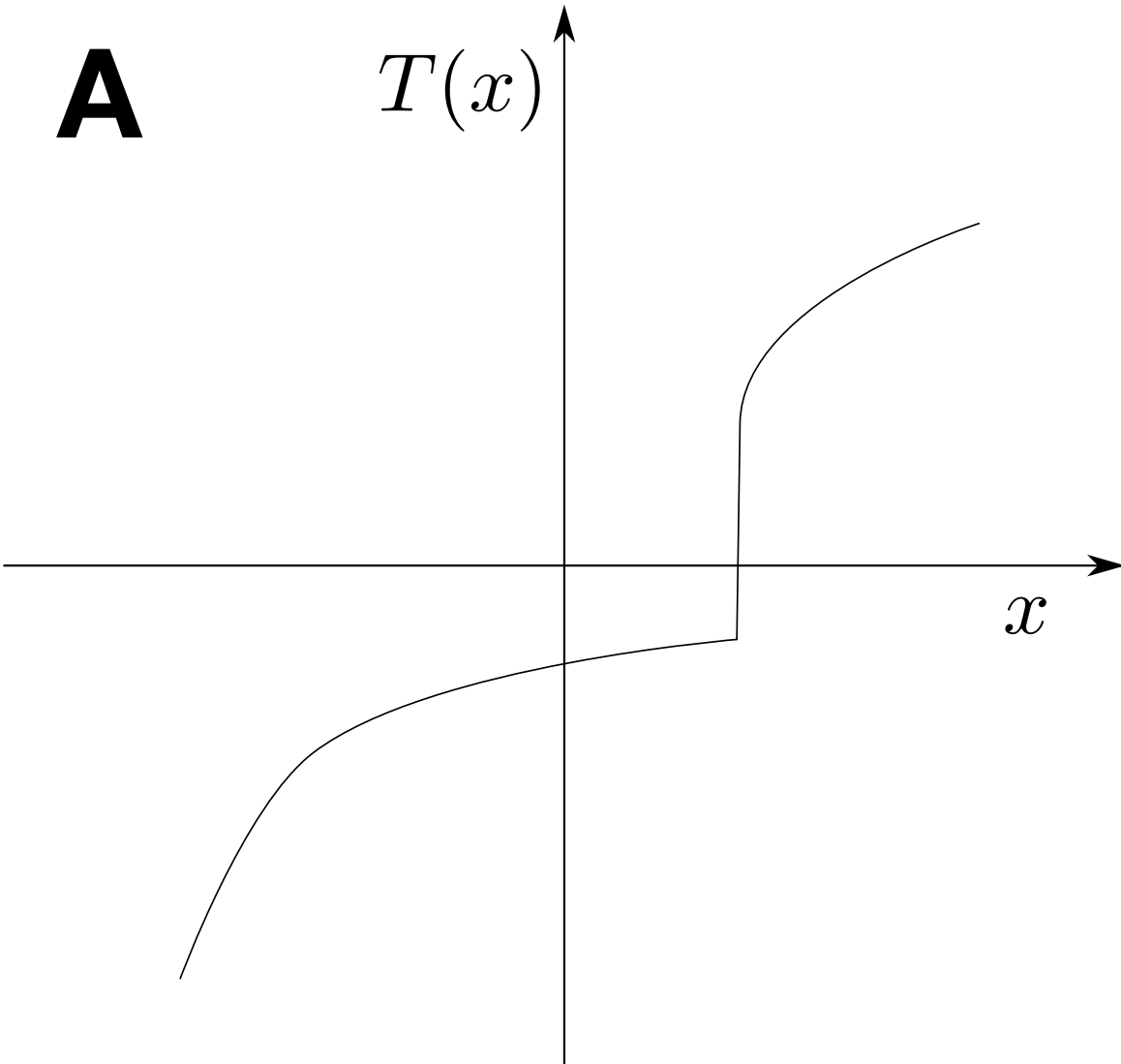
$\nexists (\bar{x}, \bar{u}) \notin \text{gph} T$ such that

$$(\bar{u} - u)^T (\bar{x} - x) \geq 0$$

Equivalently: \nexists monotone R
such that $\text{gph} T \subset \text{gph} R$

Monotone operators in 1D

Let's fill the table



	Monotone	Max Monotone
A		
B		
C		
D		

Monotonicity

$$y > x \quad \Rightarrow \quad T(y) \geq T(x)$$

Continuity

If T single-valued,
continuous and monotone,
then it's maximal monotone²²

Monotone operator properties

- **sum** $T + R$ is monotone
- **nonnegative scaling** αT with $\alpha \geq 0$ is monotone
- **inverse** T^{-1} is monotone
- **congruence** for $M \in \mathbf{R}^{n \times m}$, then $M^T T(Mz)$ is monotone on \mathbf{R}^m

Affine function $T(x) = Ax + b$ is maximal monotone
 $\iff A + A^T \succeq 0$

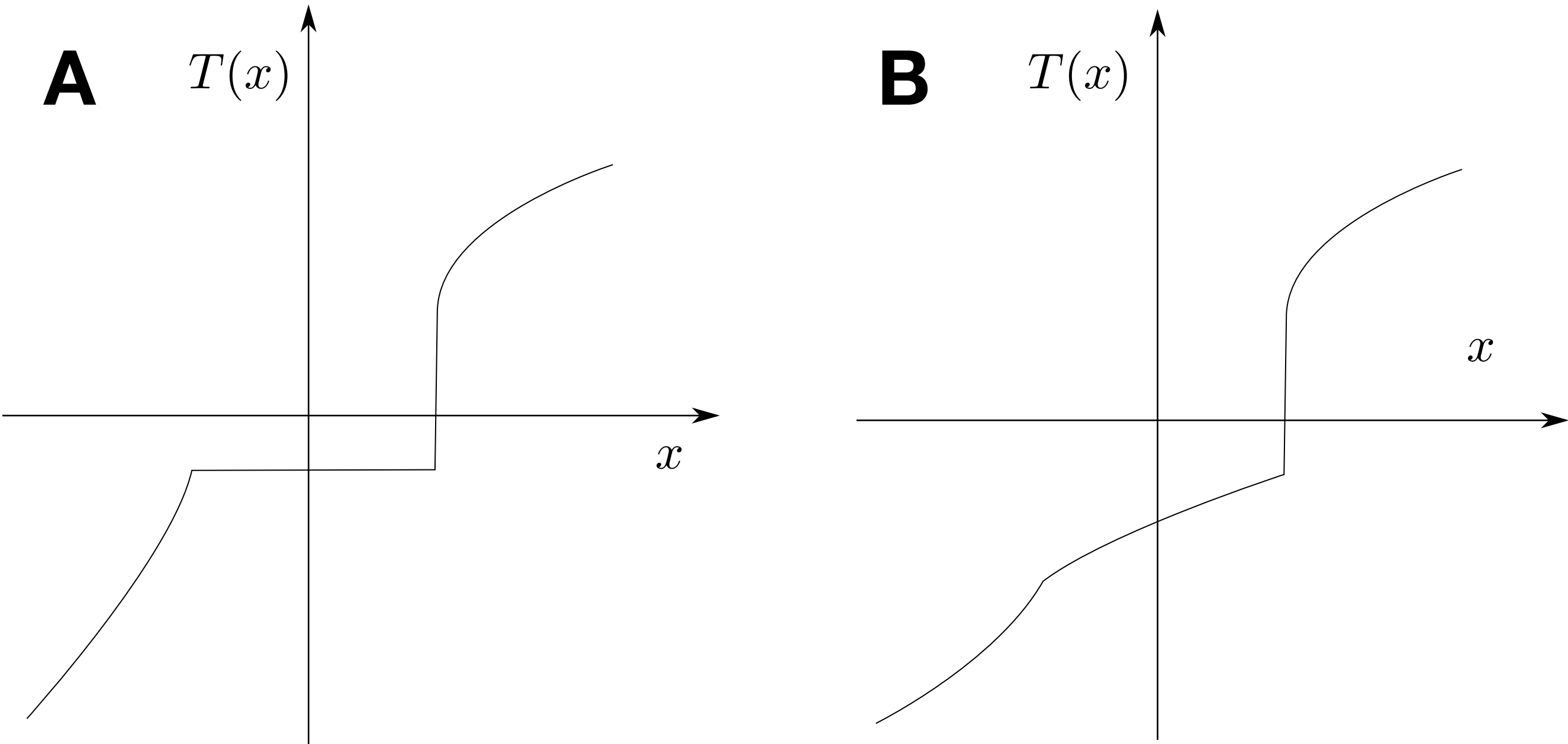
Strongly monotone operators

An operator T on \mathbb{R}^n is μ -strongly monotone if

$$(u - v)^T (x - y) \geq \mu \|x - y\|^2, \quad \mu > 0 \qquad \text{(also called } \mu\text{-coercive)}$$

$$\forall (x, u), (y, v) \in \text{gph} T$$

Let's fill the table



	Monotone	Strongly Monotone
A		
B		

The slope is at least μ

Cocoercive operators

An operator T is β -**cocoercive**, $\beta > 0$, if

$$(T(x) - T(y))^T (x - y) \geq \beta \|T(x) - T(y)\|^2$$

If T is β -**cocoercive**, then T is $(1/\beta)$ -**Lipschitz**

Proof $\beta \|T(x) - T(y)\|^2 \leq (T(x) - T(y))^T (x - y) \leq \|T(x) - T(y)\| \|x - y\|$
 $\implies \|T(x) - T(y)\| \leq (1/\beta) \|x - y\|$ ■

If T is μ -**strongly monotone** if and only if T^{-1} is μ -**cocoercive**

Proof $(T(x) - T(y))^T (x - y) \geq \mu \|x - y\|^2$

Inverse: $u = T(x)$ and $v = T(y)$ if and only if $x \in T^{-1}(u)$ and $y \in T^{-1}(v)$

$$(u - v)^T (T^{-1}(u) - T^{-1}(v)) \geq \mu \|T^{-1}(u) - T^{-1}(v)\|^2$$
 ■

Cocoercive and nonexpansive operators

If T is β -**cocoercive** if and only if $I - 2\beta T$ is **nonexpansive**

Proof

$$\begin{aligned} & \| (I - 2\beta T)(y) - (I - 2\beta T)(x) \|^2 = \\ &= \| y - 2\beta T(y) - x + 2\beta T(x) \|^2 \\ &= \| y - x \|^2 - 4\beta (T(y) - T(x))^T (y - x) + 4\beta^2 \| T(y) - T(x) \|^2 \\ &= \| y - x \|^2 - 4\beta \left((T(y) - T(x))^T (y - x) - \beta \| T(y) - T(x) \|^2 \right) \\ &\leq \| y - x \|^2 \quad \blacksquare \quad \text{(cocoercive)} \end{aligned}$$

Summary of monotone and cocoercive operators

Monotone

$$(T(x) - T(y))^T (x - y) \geq 0$$

$$\uparrow \mu = 0$$

Strongly monotone

$$(T(x) - T(y))^T (x - y) \geq \mu \|x - y\|^2$$

$$\longleftrightarrow_{F = T^{-1}}$$

Lipschitz

$$\|F(x) - F(y)\| \leq L \|x - y\|$$

$$\uparrow L = 1/\mu$$

Cocoercive

$$(F(x) - F(y))^T (x - y) \geq \mu \|F(x) - F(y)\|^2$$

$$\updownarrow G = I - 2\mu F$$

Nonexpansive

$$\|G(x) - G(y)\| \leq \|x - y\|$$

Fixed point iterations

Fixed point iteration

Apply operator

$$x^{k+1} = T(x^k)$$

until you reach $\bar{x} \in \text{fix } T$

Main approach

1. Find a suitable T such that $\bar{x} \in \text{fix } T$ solve your problem
2. Show that the fixed point iteration converges

Fixed point residual to terminate

$$r^k = T(x^k) - x^k$$

Contractive fixed point iterations

Contraction mapping theorem

If T is L -Lipschitz with $L < 1$ (contraction), the iteration

$$x^{k+1} = T(x^k)$$

converges to \bar{x} , the unique fixed point of T

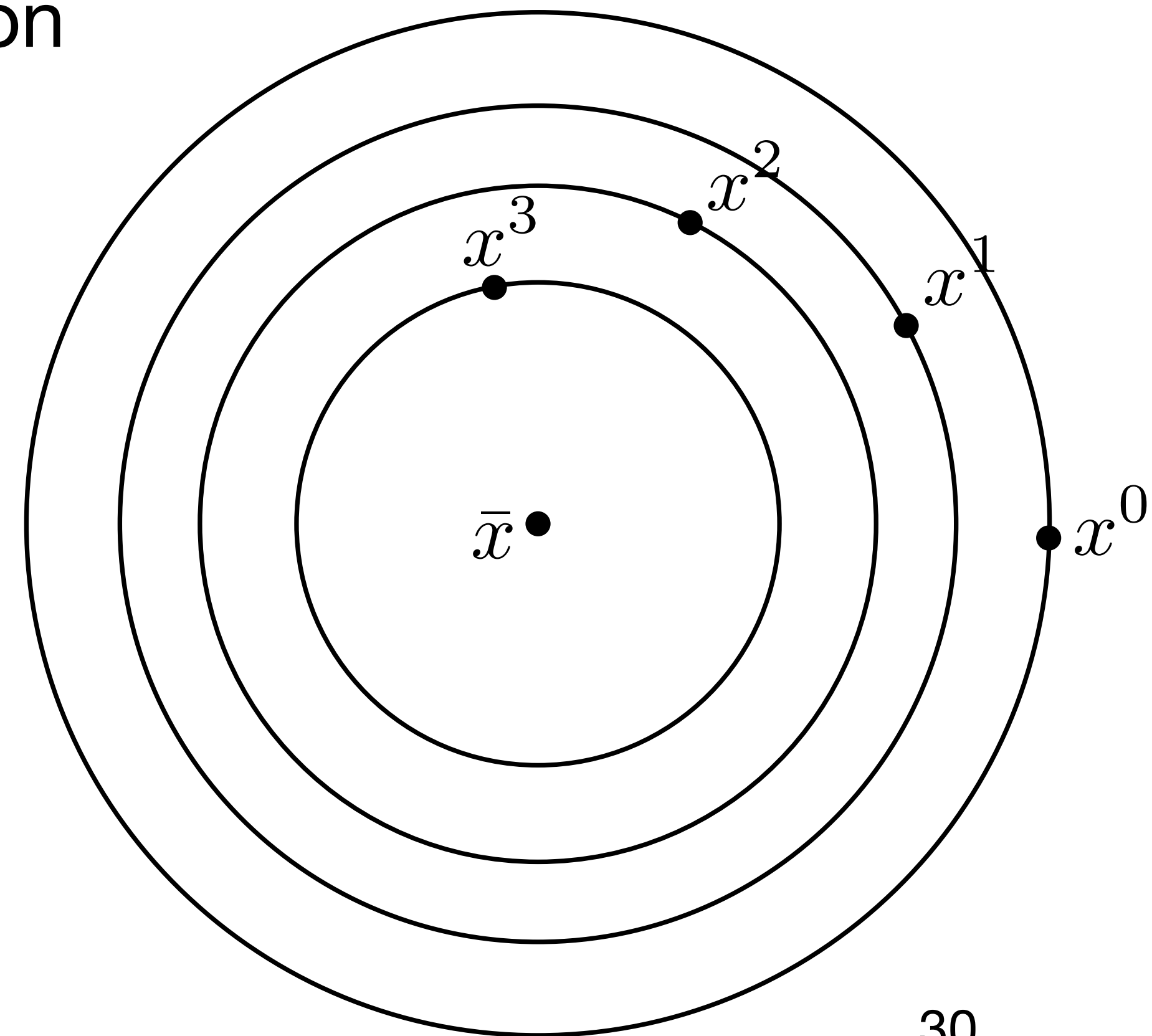
Properties

- Distance to \bar{x} decreases at each step

$$\|x^{k+1} - \bar{x}\| \leq L \|x^k - \bar{x}\|$$

(iteration is **Fejer monotone**)

- Linear convergence rate L



Contraction mapping theorem

Proof

The sequence x^k is Cauchy

$$\begin{aligned}\|x^{k+\ell} - x^k\| &\leq \|x^{k+\ell} - x^{k+\ell-1}\| + \dots + \|x^{k+1} - x^k\| \\ &\leq (L^{\ell-1} + \dots + 1)\|x^{k+1} - x^k\| \\ &\leq \frac{1}{1-L}\|x^{k+1} - x^k\| \\ &\leq \frac{L^k}{1-L}\|x^1 - x^0\|\end{aligned}$$

(Lipschitz constant)

(geometric series)

(Lipschitz constant)

Therefore it converges to a point \bar{x} which must be the (unique) fixed point of T

The convergence is linear (geometric) with rate L

$$\|x^k - \bar{x}\| = \|T(x^{k-1}) - T(\bar{x})\| \leq L\|x^{k-1} - \bar{x}\| \leq L^k\|x^0 - x^*\|$$



Nonexpansive fixed point iterations

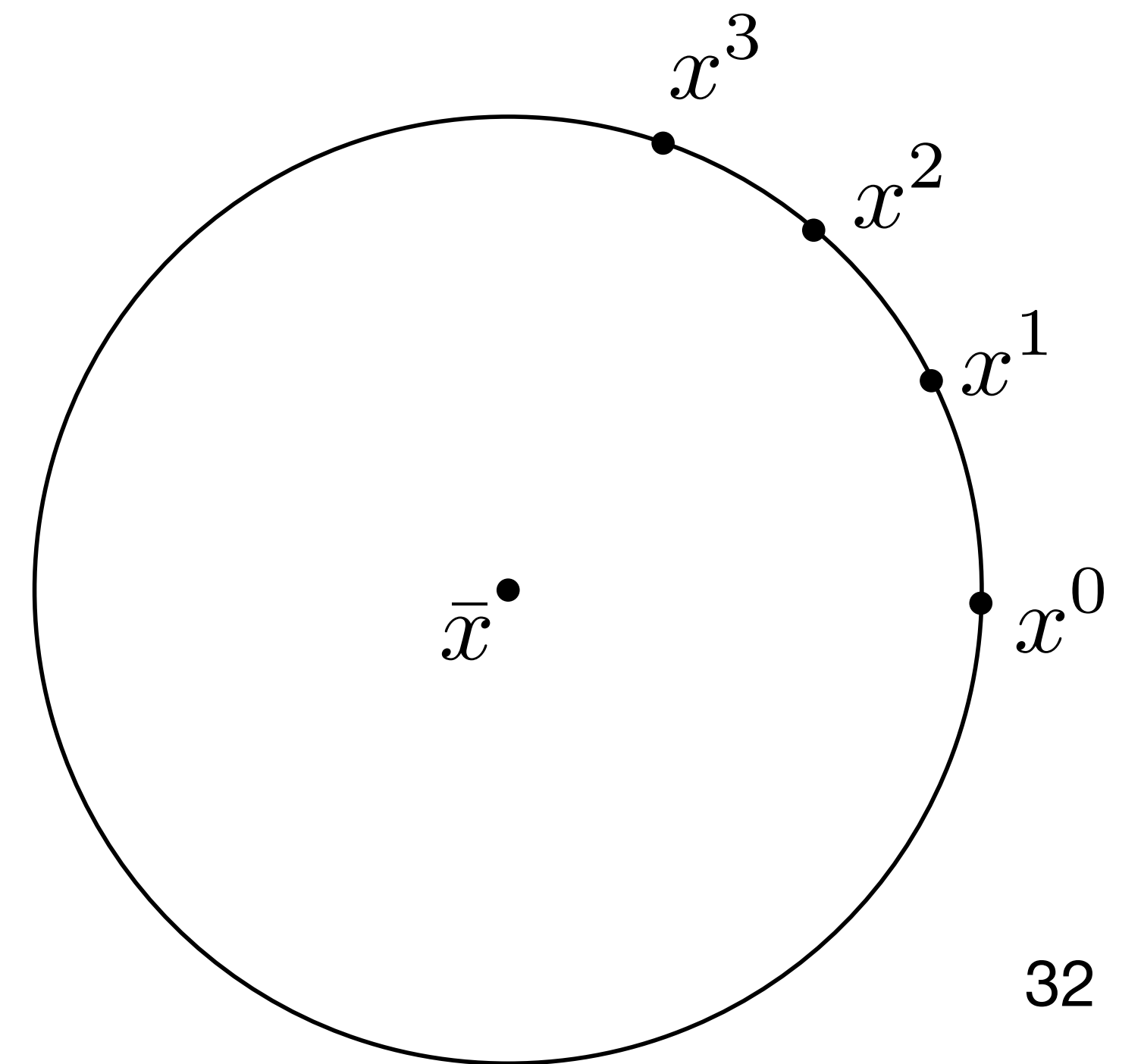
If T is L -Lipschitz with $L = 1$ (nonexpansive), the iteration

$$x^{k+1} = T(x^k)$$

need not converge to a fixed point, even if one exists.

Example

- Let T be a rotation around the origin
- T is nonexpansive and has a fixed point $\bar{x} = 0$
- $\|x^k\|$ never decreases

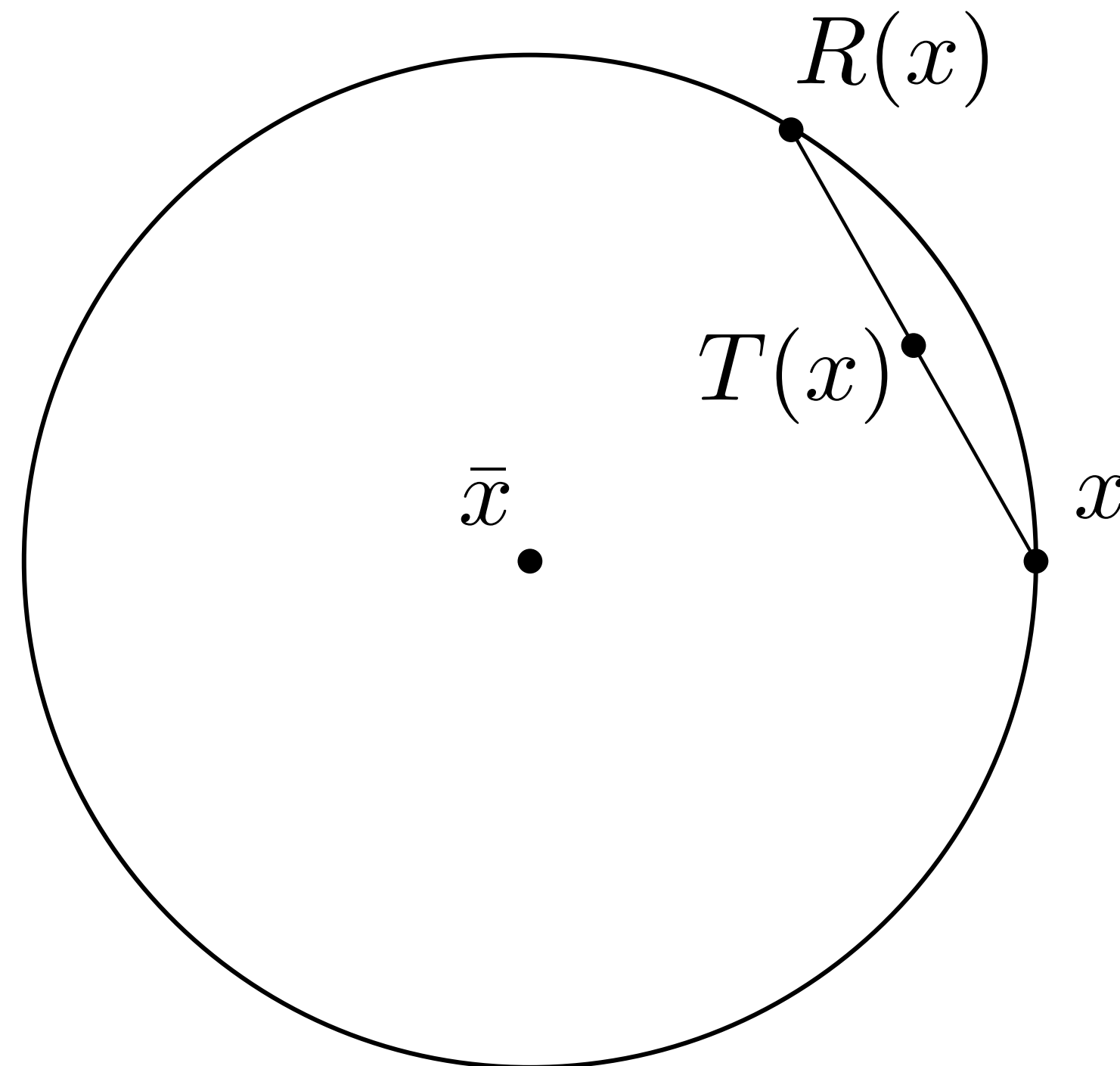


Averaged operators

We say that an operator T is α -**averaged** with $\alpha \in (0, 1)$ if

$$T = (1 - \alpha)I + \alpha R$$

and R is nonexpansive.




Averaged operators fixed points

We say that an operator T is α -**averaged** with $\alpha \in (0, 1)$ if

$$T = (1 - \alpha)I + \alpha R$$

Fact If T is α -averaged, then $\text{fix } T = \text{fix } R$

Proof $\bar{x} = T(\bar{x}) = (1 - \alpha)I(\bar{x}) + \alpha R(\bar{x})$
 $= (1 - \alpha)\bar{x} + \alpha R(\bar{x})$
 $\iff \alpha\bar{x} = \alpha R(\bar{x})$
 $\iff \bar{x} = R(\bar{x})$ 

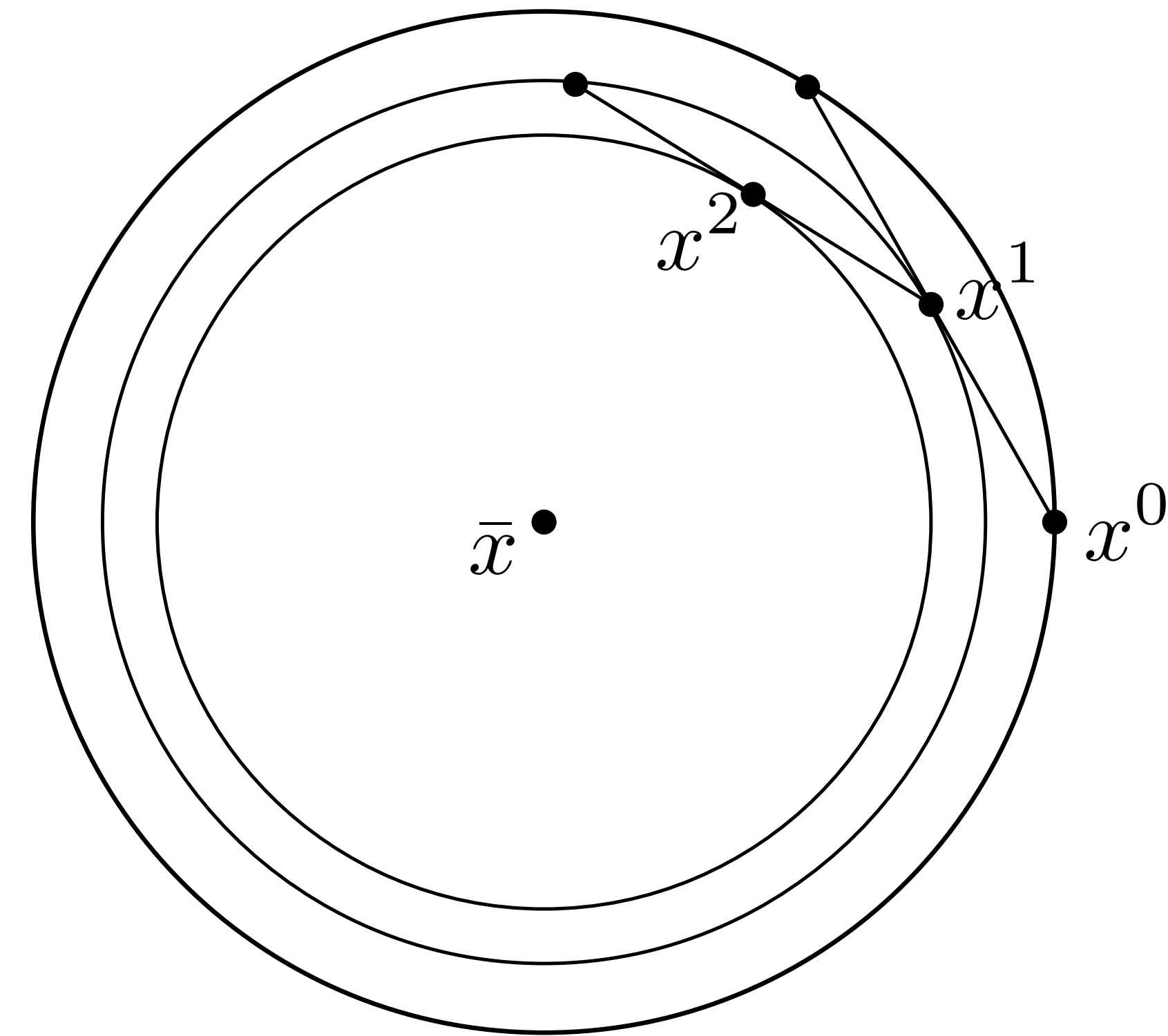
Averaged fixed point iterations

If $T = (1 - \alpha)I + \alpha R$ is α -averaged
($\alpha \in (0, 1)$ and R nonexpansive), the iteration

$$x^{k+1} = T(x^k)$$

converges to $\bar{x} \in \mathbf{fix} T$

(also called damped, averaged
or Mann-Krasnosel'skii iteration)



Properties

- Distance to \bar{x} decreases at each step (**Fejer monotone**)
- Sublinear convergence to fixed-point residual

$$\|R(x^k) - x^k\| \leq \frac{1}{\sqrt{(k+1)\alpha(1-\alpha)}} \|x^0 - \bar{x}\|$$

Averaged fixed point iterations

Proof

Use the identity (proof by expanding)

$$\|(1 - \alpha)a + \alpha b\|^2 = (1 - \alpha)\|a\|^2 + \alpha\|b\|^2 - \alpha(1 - \alpha)\|a - b\|^2$$

and apply it to

$$x^{k+1} - \bar{x} = (1 - \alpha) \underbrace{(x^k - \bar{x})}_a + \alpha \underbrace{(R(x^k) - \bar{x})}_b$$

obtaining

$$\begin{aligned} \|x^{k+1} - \bar{x}\|^2 &= (1 - \alpha)\|x^k - \bar{x}\|^2 + \alpha\|R(x^k) - \bar{x}\|^2 - \alpha(1 - \alpha)\|x^k - R(x^k)\|^2 \\ &\leq (1 - \alpha)\|x^k - \bar{x}\|^2 + \alpha\|x^k - \bar{x}\|^2 - \alpha(1 - \alpha)\|x^k - R(x^k)\|^2 \quad (\text{nonexpansive}) \\ &= \|x^k - \bar{x}\|^2 - \alpha(1 - \alpha)\|x^k - R(x^k)\|^2 \\ &\leq 0 \end{aligned}$$

Iterations are Fejer monotone

Averaged fixed point iterations

Proof (continued)

iterate righthand side over k steps

$$\|x^{k+1} - \bar{x}\|^2 \leq \|x^0 - \bar{x}\|^2 - \alpha(1 - \alpha) \sum_{i=0}^k \|x^i - R(x^i)\|^2$$

Since $\|x^{k+1} - \bar{x}\|^2 \geq 0$, we have
$$\sum_{i=0}^k \|x^i - R(x^i)\|^2 \leq \frac{1}{\alpha(1 - \alpha)} \|x^0 - \bar{x}\|^2$$

Using $\sum_{i=0}^k \|x^i - R(x^i)\|^2 \geq (k + 1) \min_{i=0, \dots, k} \|x^i - R(x^i)\|^2$, we obtain

$$\min_{i=0, \dots, k} \|x^i - R(x^i)\|^2 \leq \frac{1}{(k + 1)\alpha(1 - \alpha)} \|x^0 - \bar{x}\|^2$$

(R is nonexpansive \rightarrow min at k)
$$\|x^k - R(x^k)\|^2 \leq \frac{1}{(k + 1)\alpha(1 - \alpha)} \|x^0 - \bar{x}\|^2 \quad \blacksquare \quad 37$$

Average fixed point iteration convergence rates

$$\|R(x^k) - x^k\| \leq \frac{1}{\sqrt{(k+1)\alpha(1-\alpha)}} \|x^0 - \bar{x}\|$$

Righthand side minimized when $\alpha = 1/2$

$$\|R(x^k) - x^k\| \leq \frac{2}{\sqrt{k+1}} \|x^0 - \bar{x}\|$$

Iterations

$$x^{k+1} = (1/2)x^k + (1/2)R(x^k)$$

Remarks

- Sublinear convergence (same as subgrad method), in general not the actual rate
- $\alpha = 1/2$ is very common for averaged operators

How to design an algorithm

Problem

minimize $f(x)$

Algorithm (operator) construction

1. Find a suitable T such that $\bar{x} \in \text{fix } T$ solve your problem
2. Show that the fixed point iteration converges

If T is contractive \implies **linear convergence**

If T is averaged \implies **sublinear convergence**

Most first order algorithms can be constructed in this way

Operator theory

Today, we learned to:

- **Define** and **evaluate** proximal operators for various common functions
- **Apply** proximal operators to generalize gradient descent (vanilla, projected, proximal)
- **Define** monotone and cocoercive operators and their relations
- **Use operator theory** to construct general fixed-point iterations and prove their convergence

Next lecture

- Operators in optimization algorithms