

# The Voice of Optimization

Bartolomeo Stellato

joint work with Dimitris Bertsimas

ORC IAP Seminar 2019





# Inventory management

minimize  $\sum_{t=0}^{T-1} h(x_t) + o(u_t)$

subject to  $x_{t+1} = x_t + u_t - d_t$   
 $x_0 = x_{\text{init}}$   
 $0 \leq u_t \leq M$

A photograph showing a close-up view of a metal shelving unit in a warehouse. The shelves are filled with numerous cardboard boxes, each with a white shipping label. The boxes are stacked in several layers across the visible height of the shelves.

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Inventory 

Order 



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Order

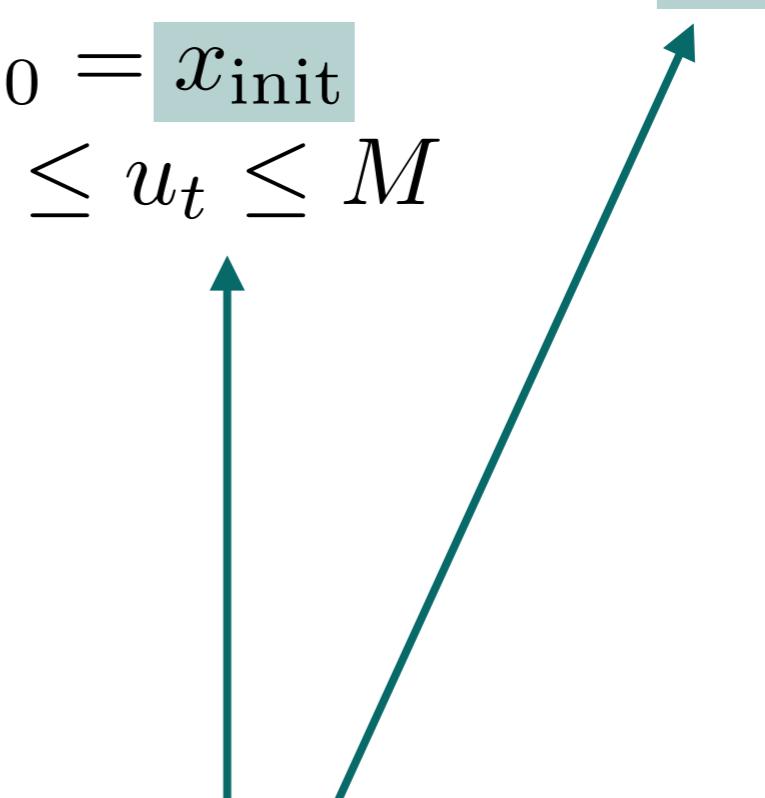
Demand

A photograph showing a close-up view of a metal shelving unit in a warehouse. The shelves are filled with numerous cardboard boxes, each with a white shipping label. The boxes are stacked in several layers across the visible height of the shelves.

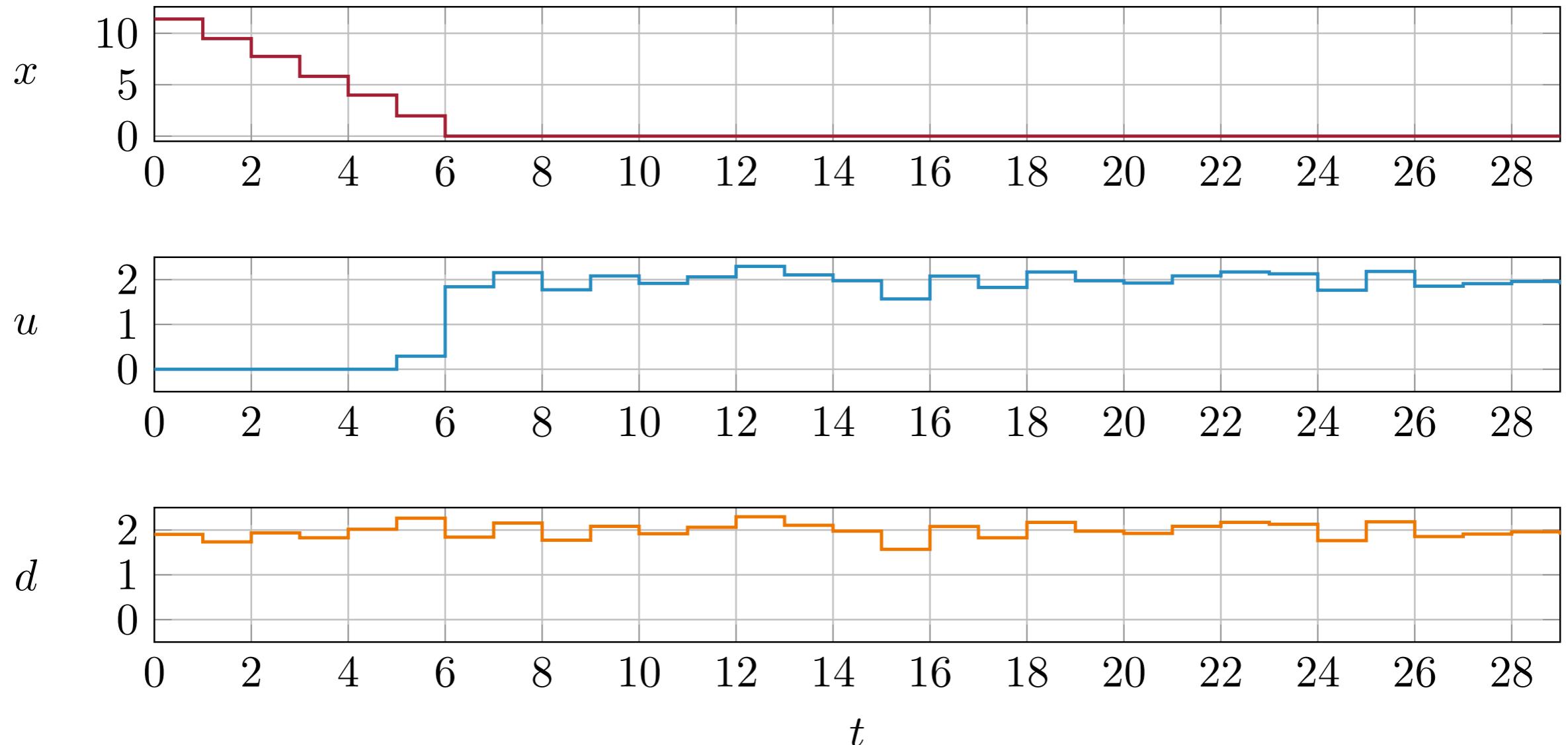
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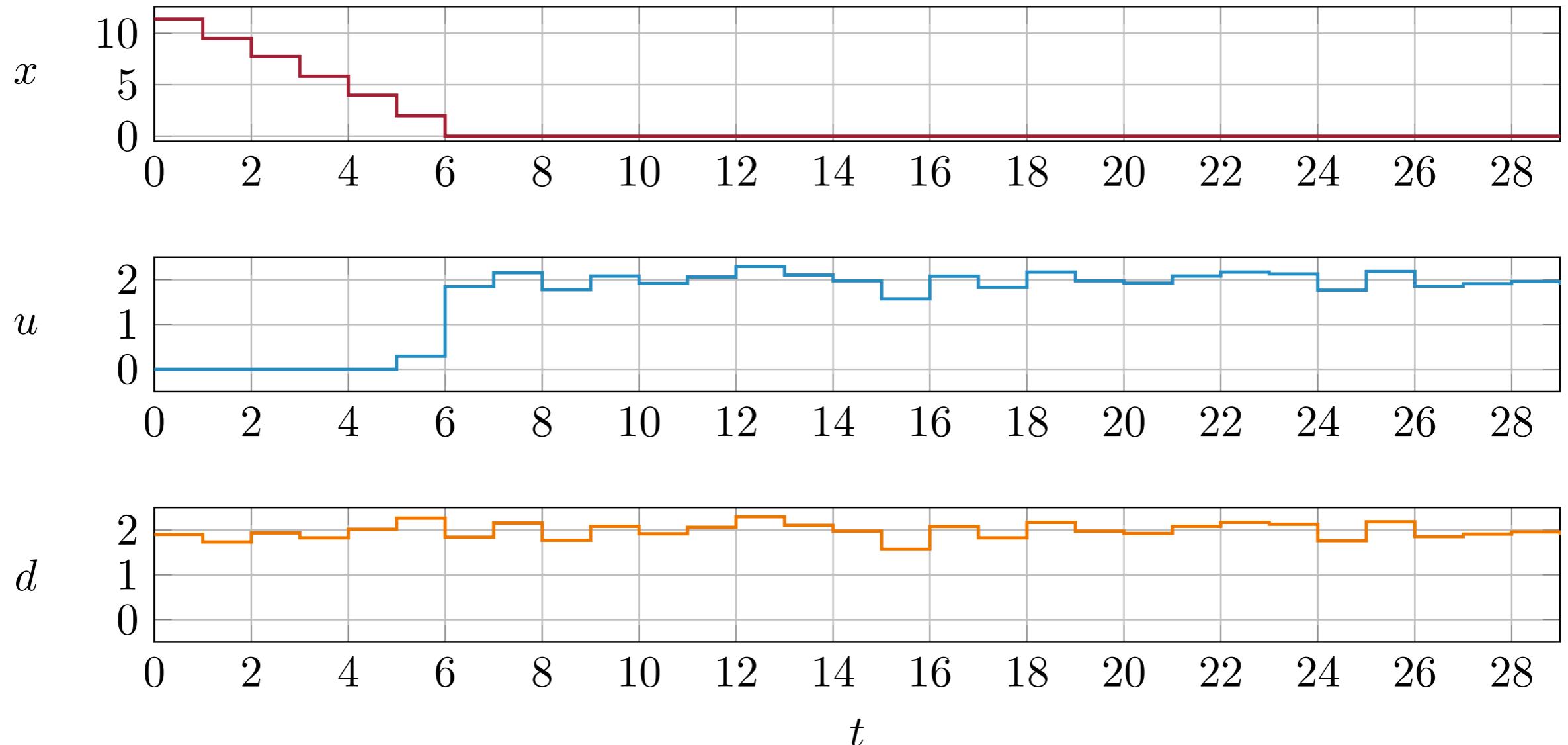
subject to  $x_{t+1} = x_t + u_t - d_t$   
 $x_0 = x_{\text{init}}$   
 $0 \leq u_t \leq M$



# What does it mean?



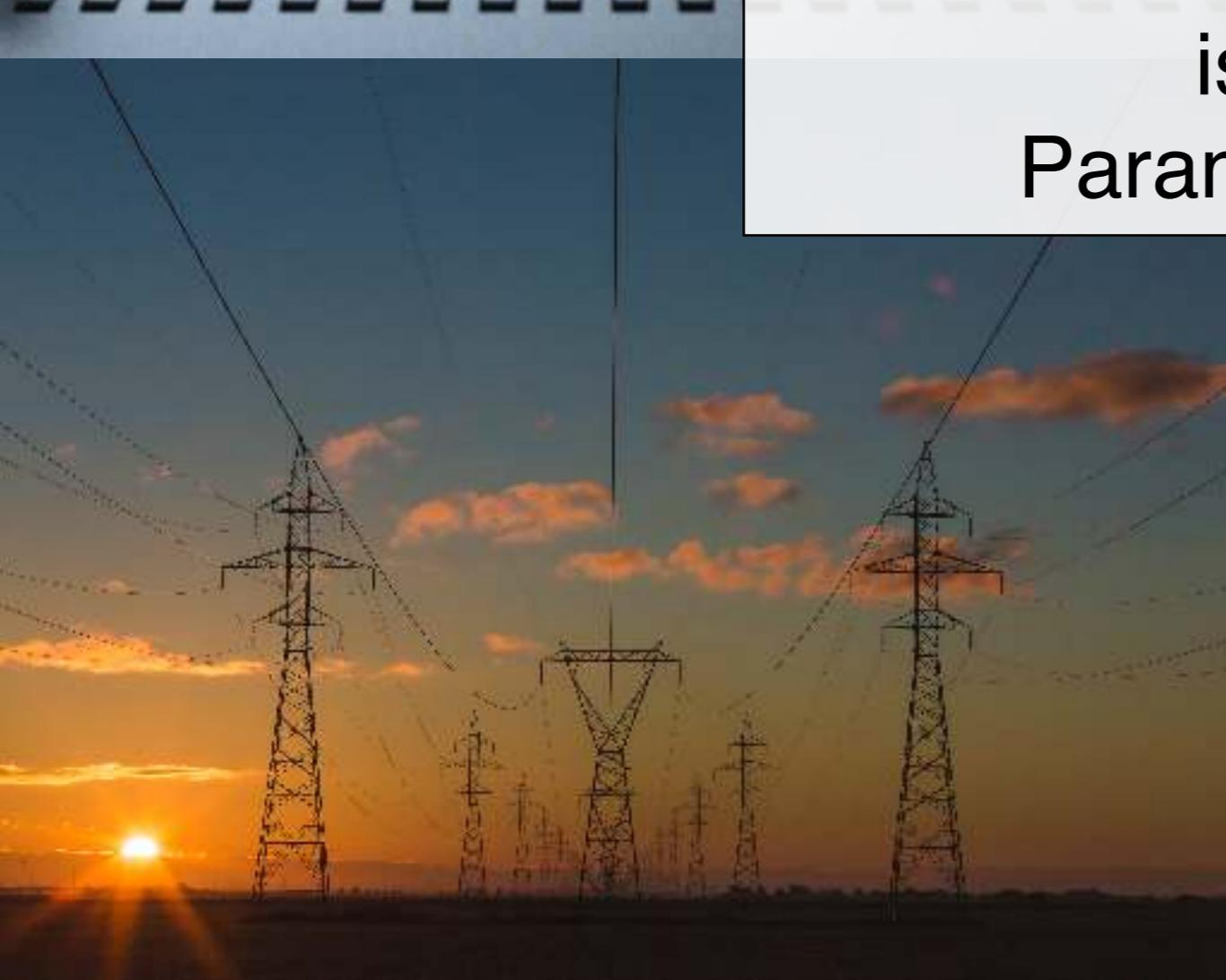
# What does it mean?



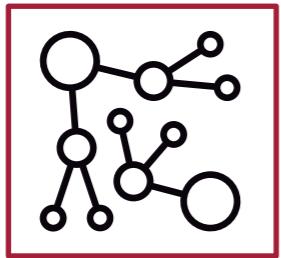
$x_{\text{init}}$  ?       $d_t$  ?



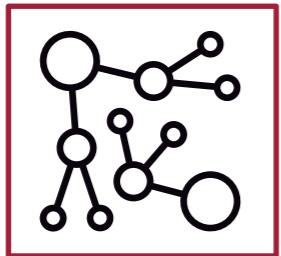
Real-world Optimization  
is  
Parametric



# Data



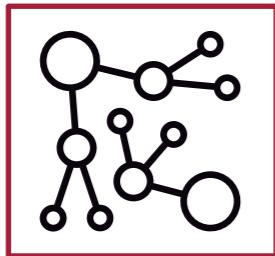
Data



Practitioner



Data

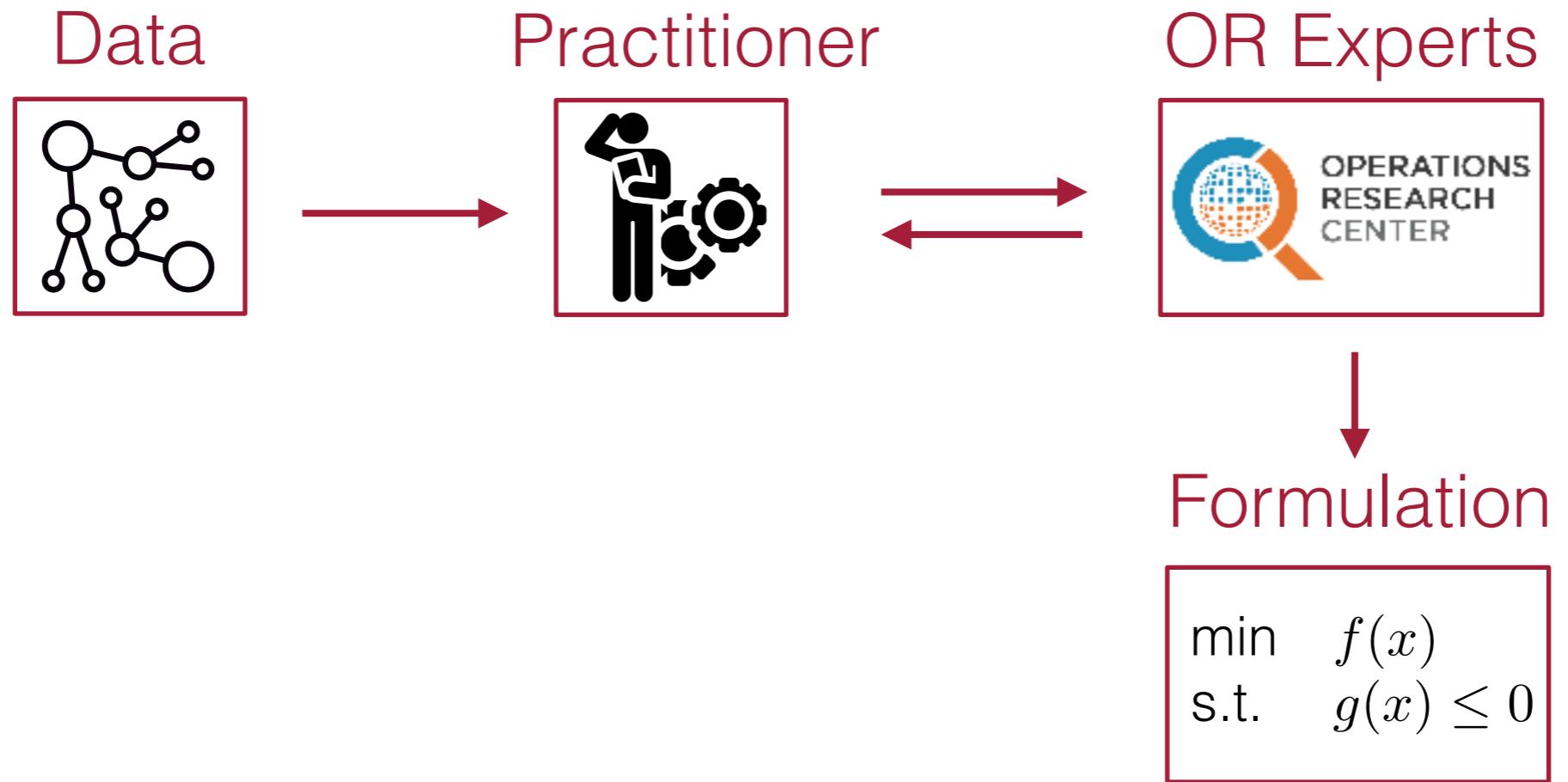


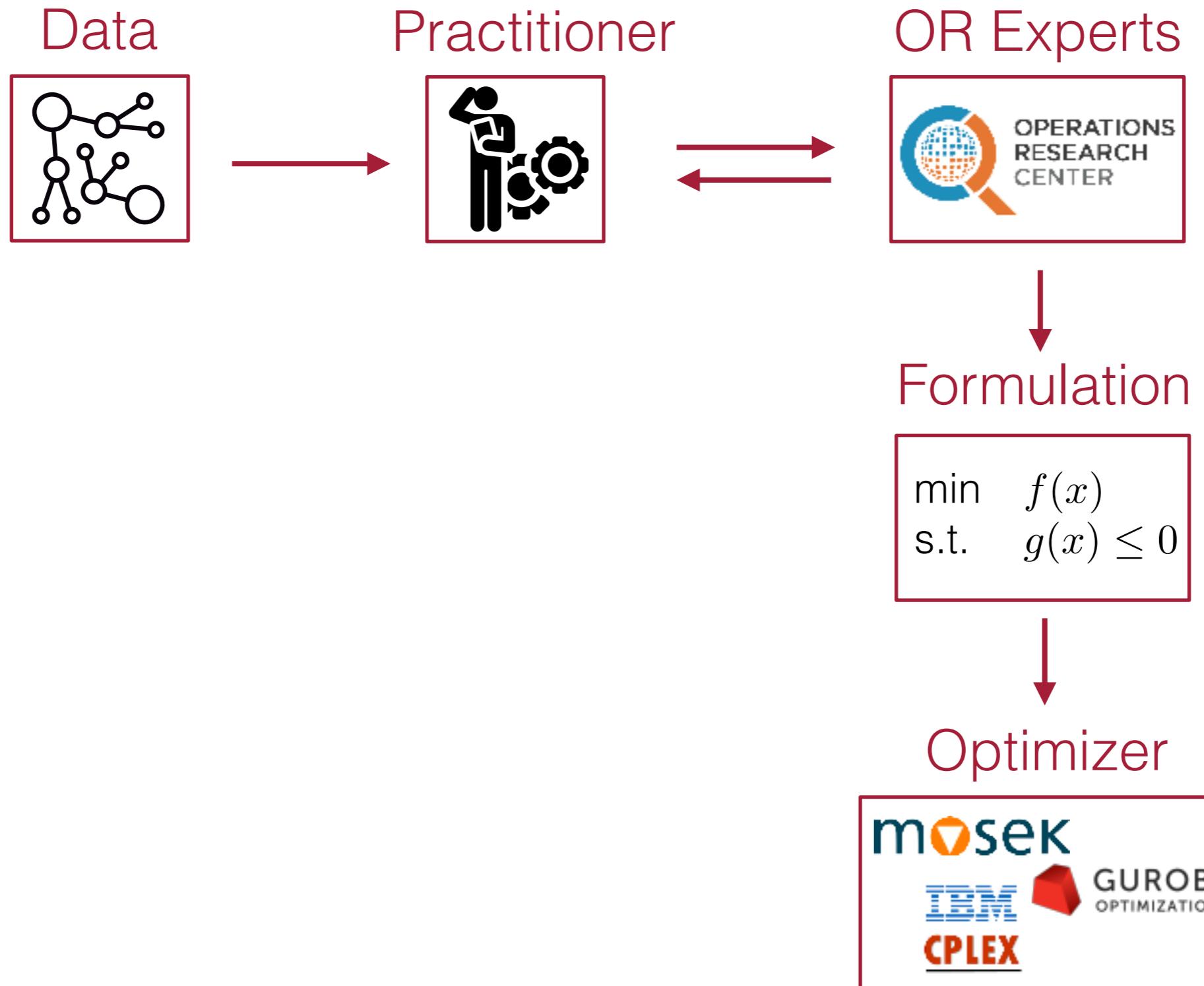
Practitioner



OR Experts









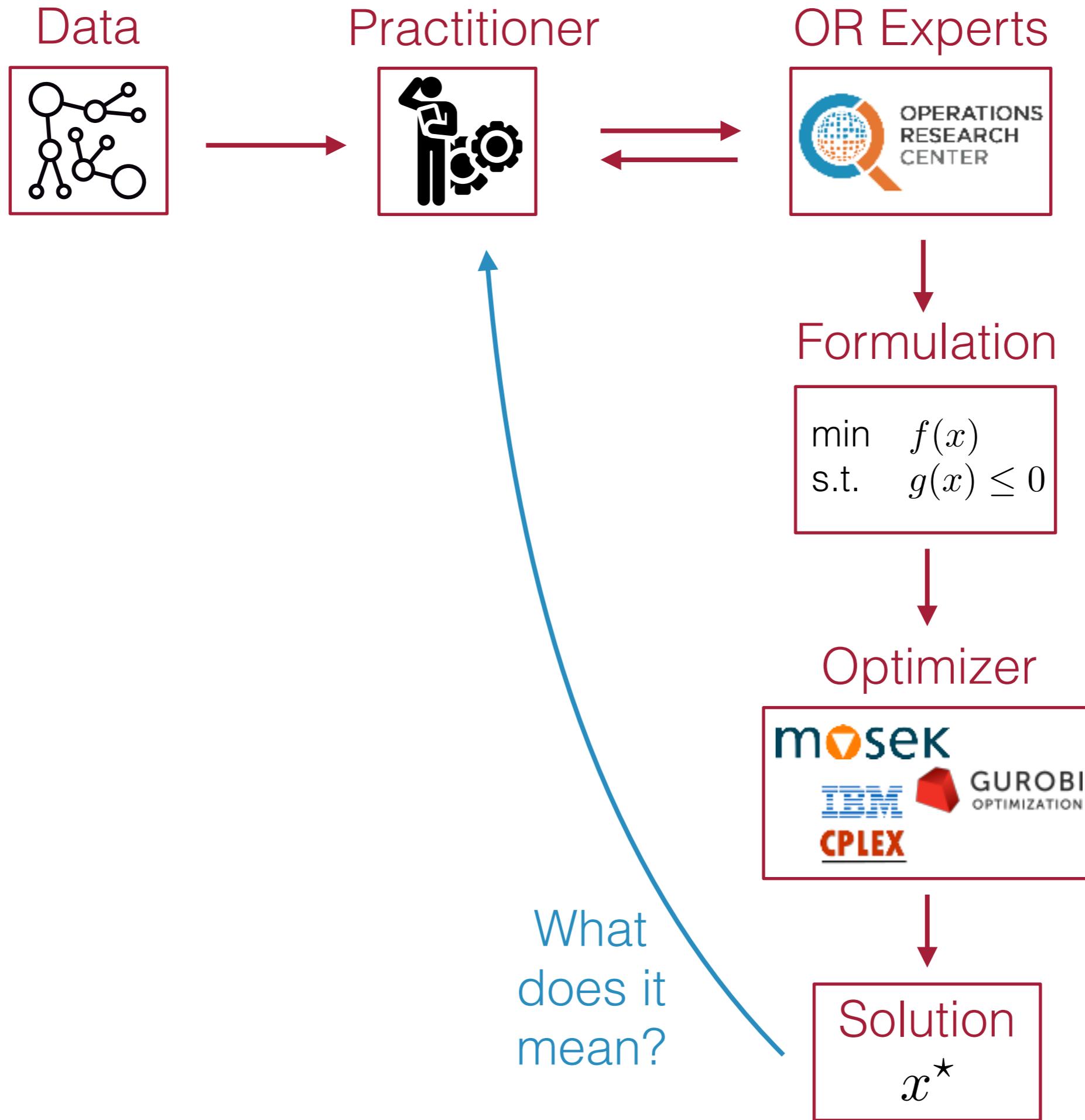
Formulation

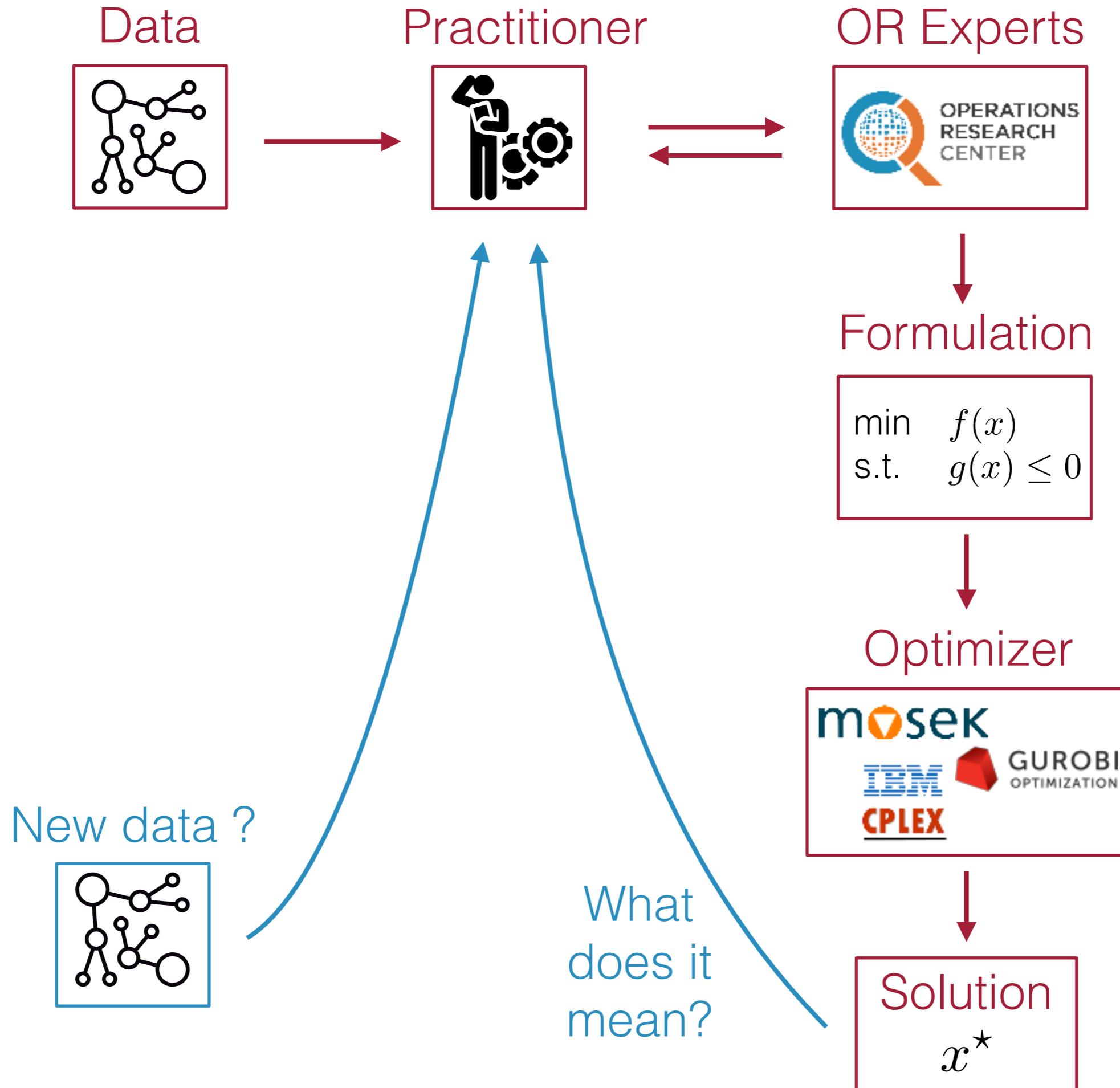
$$\begin{aligned} \min \quad & f(x) \\ \text{s.t.} \quad & g(x) \leq 0 \end{aligned}$$

Optimizer



Solution  
 $x^*$









Black box

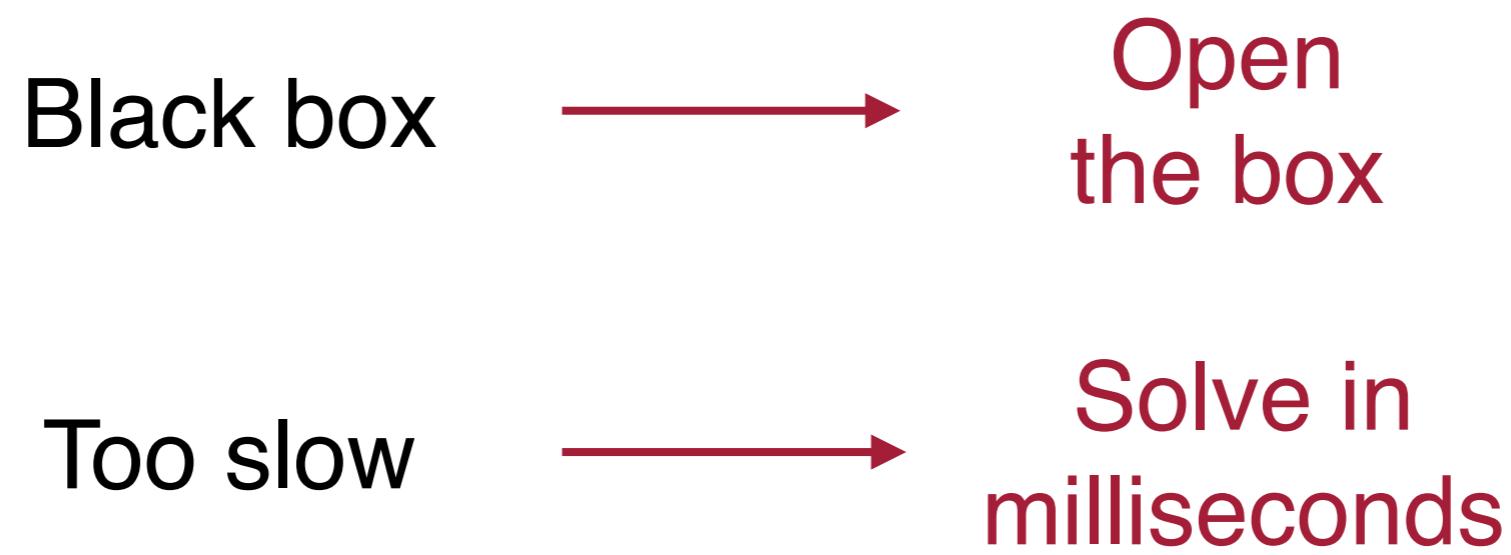


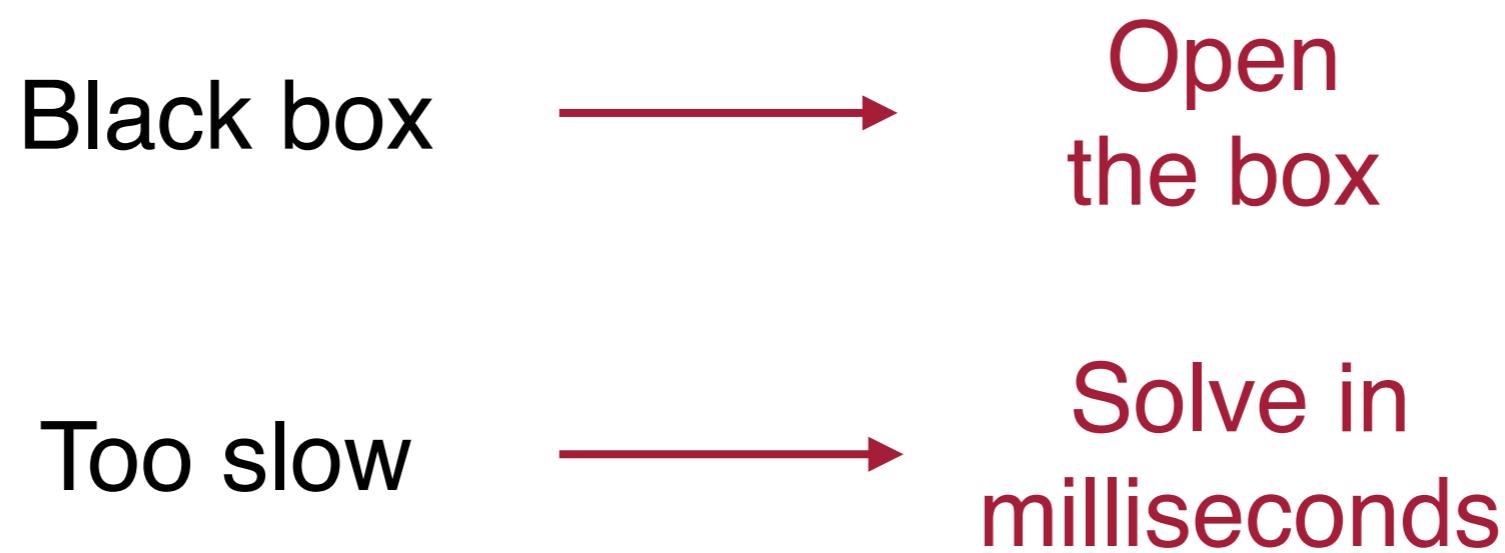
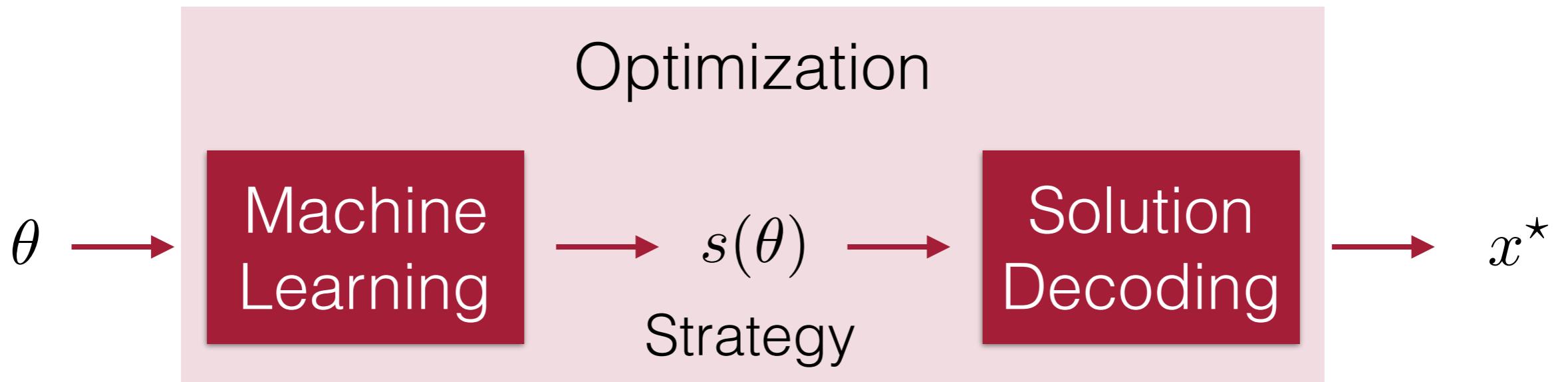
Black box → Open  
the box



Black box → Open  
the box

Too slow





# Optimal strategies

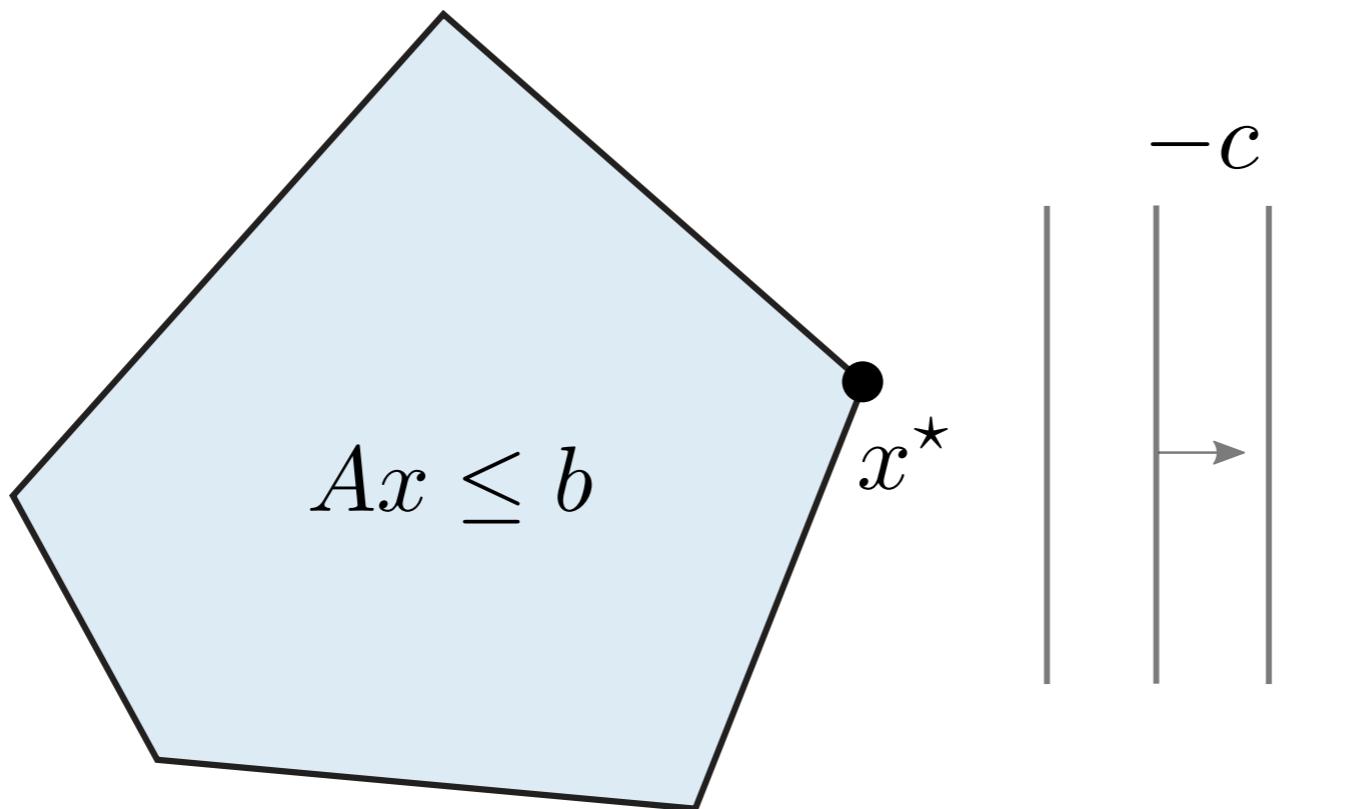
# Strategy

*The complete information to efficiently get  
the optimal solution*

$$s(\theta)$$

# Linear optimization

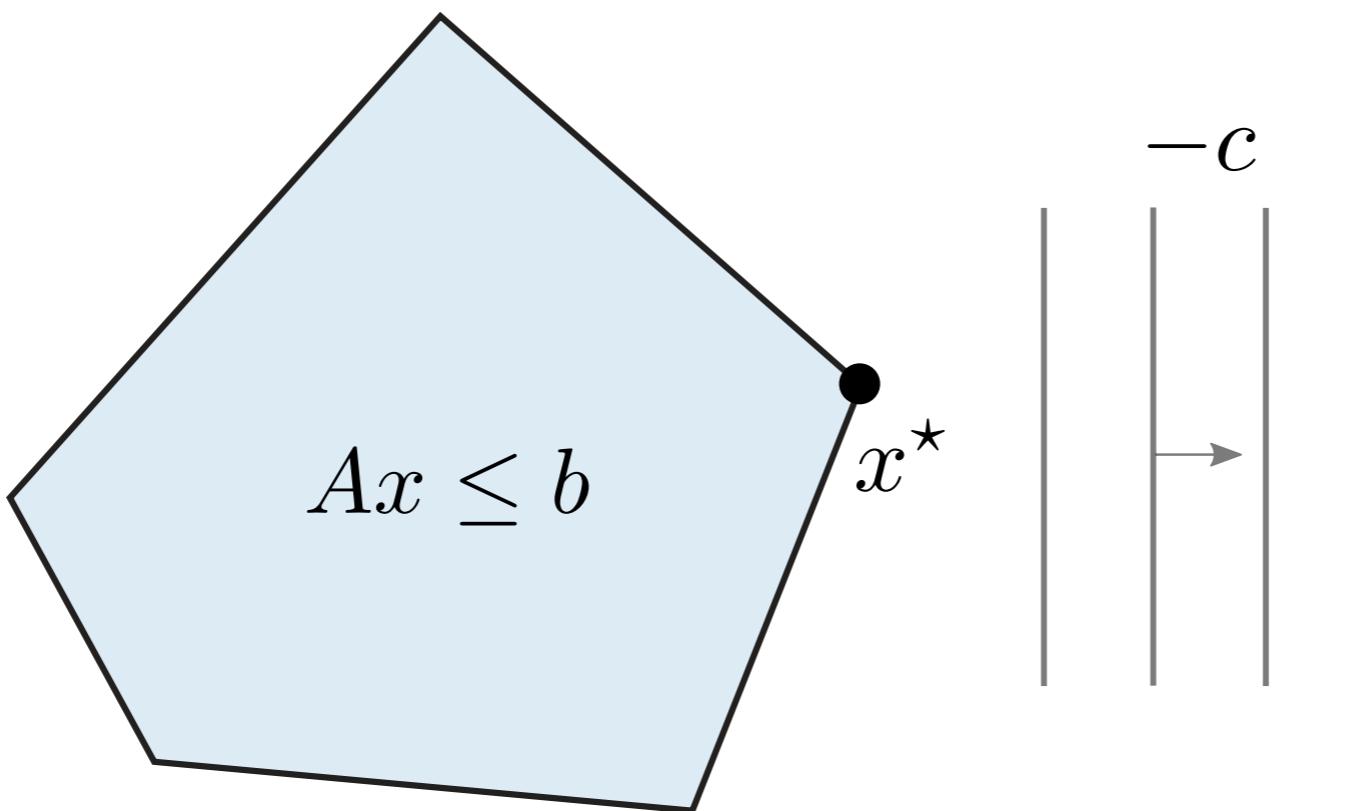
$$\begin{array}{ll}\text{minimize} & c(\theta)^T x \\ \text{subject to} & A(\theta)x \leq b(\theta)\end{array}$$



# Linear optimization

$$\begin{array}{ll}\text{minimize} & c(\theta)^T x \\ \text{subject to} & A(\theta)x \leq b(\theta)\end{array}$$

Parameters

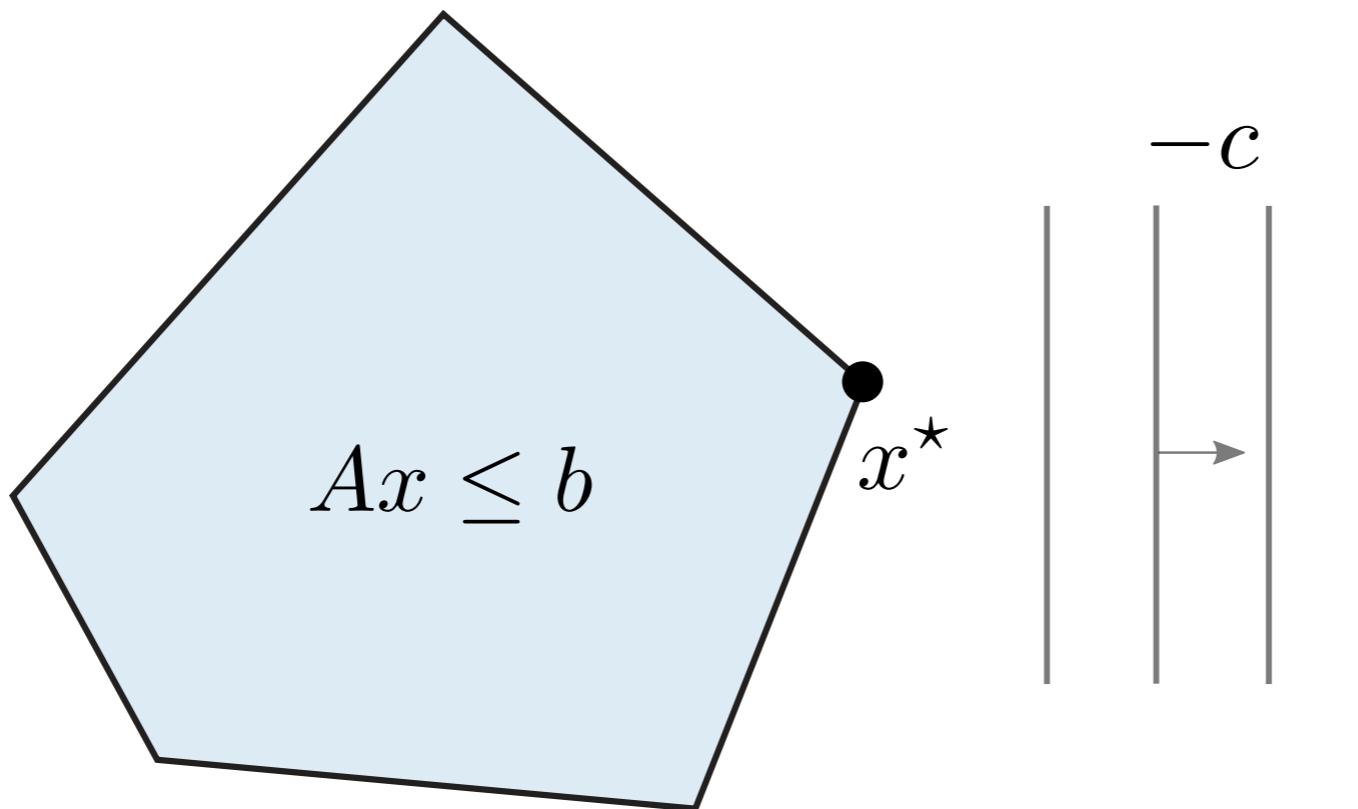


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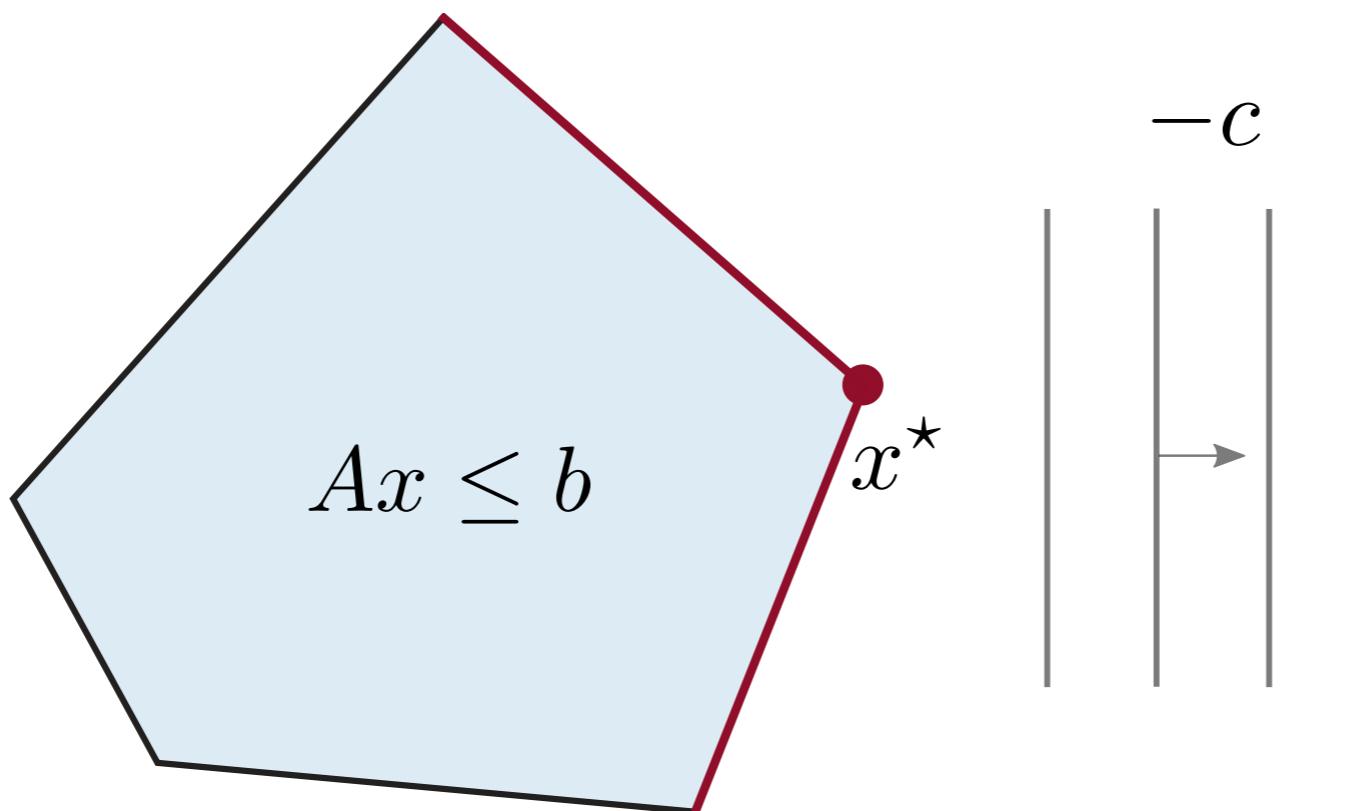
Parameters

Variables



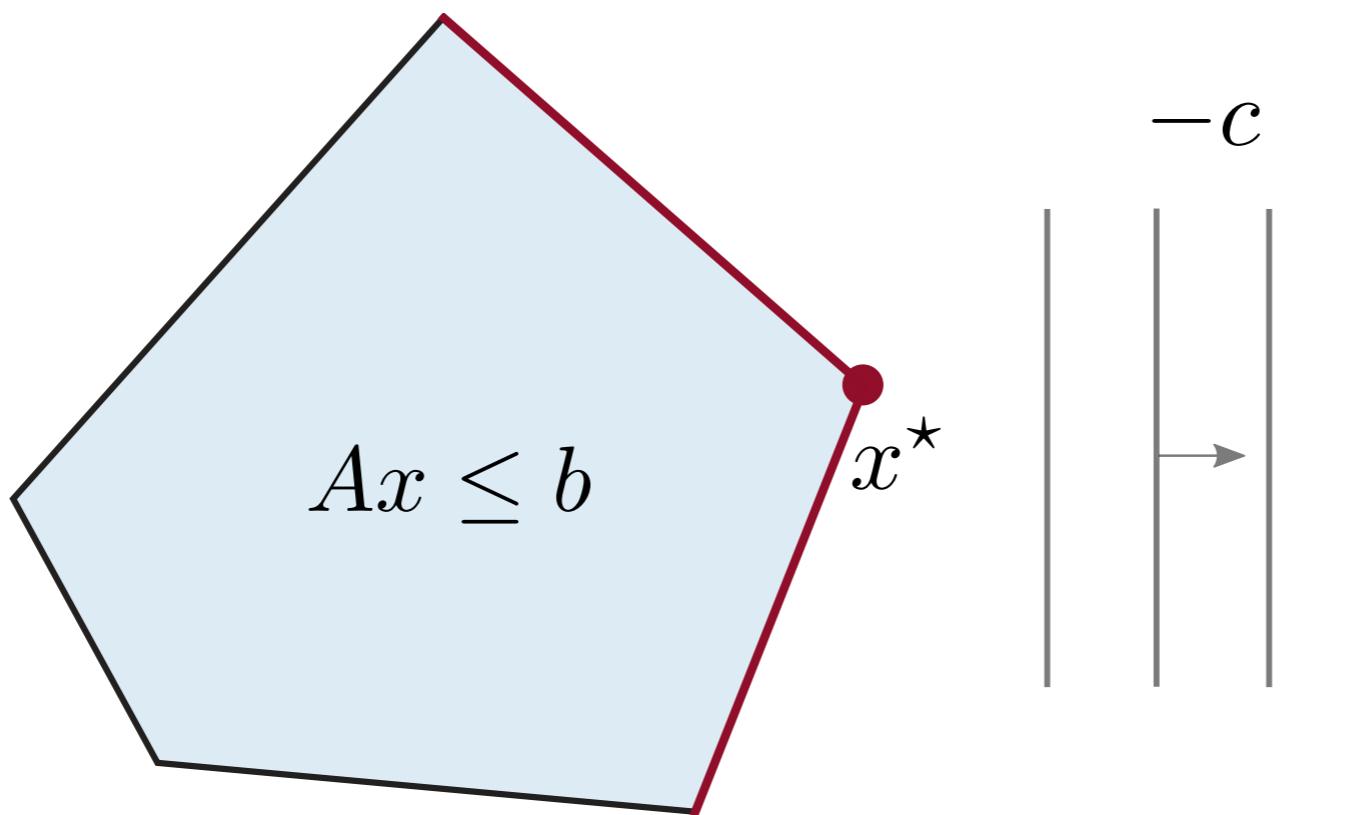
# Tight constraints

$$\mathcal{T}(\theta) = \{i \mid A_i(\theta)x^* = b_i(\theta)\}$$



# Tight constraints

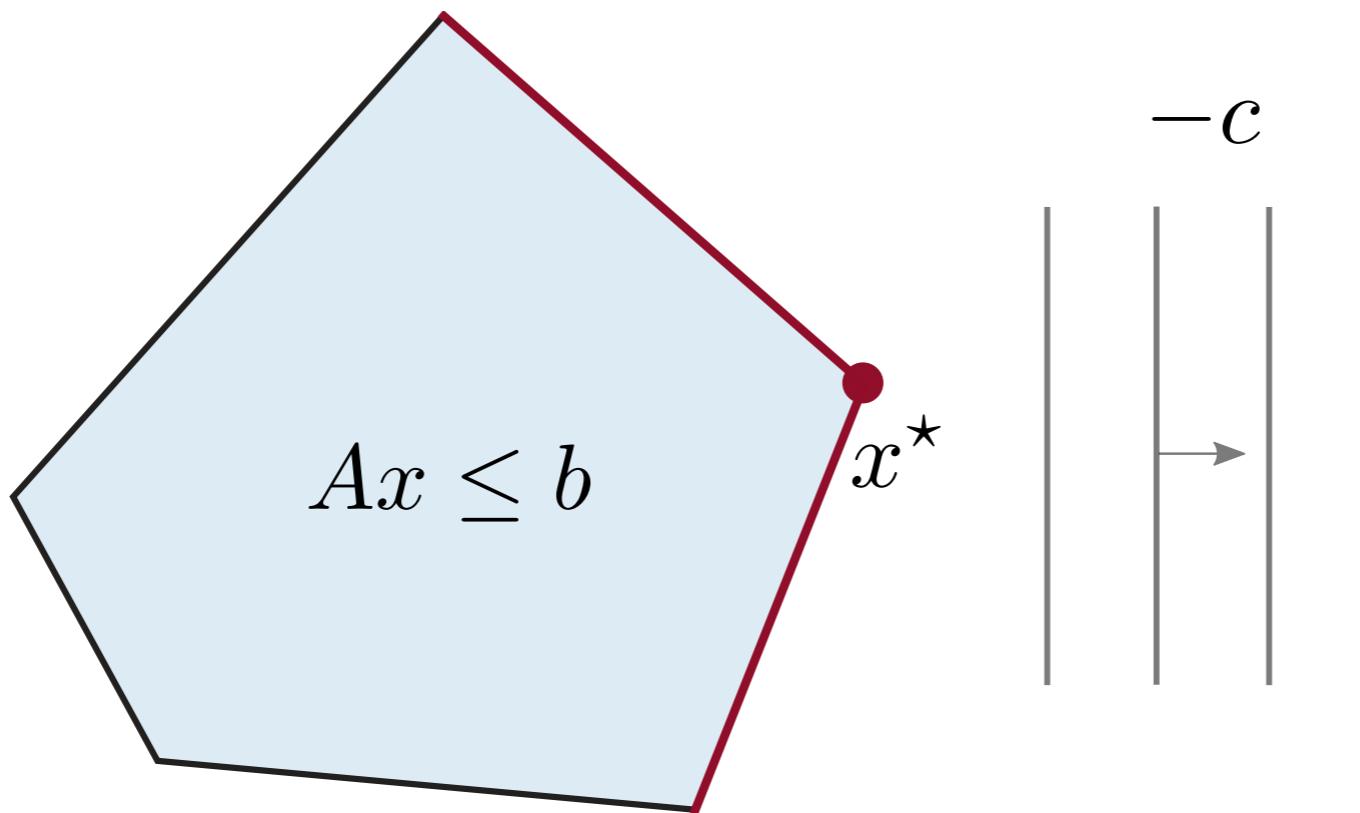
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$|\mathcal{T}(\theta)| = \# \text{ variables}$  if non-degenerate

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$|\mathcal{T}(\theta)| = \# \text{ variables}$

if non-degenerate  
in general

$|\mathcal{T}(\theta)| \ll \# \text{ constraints}$

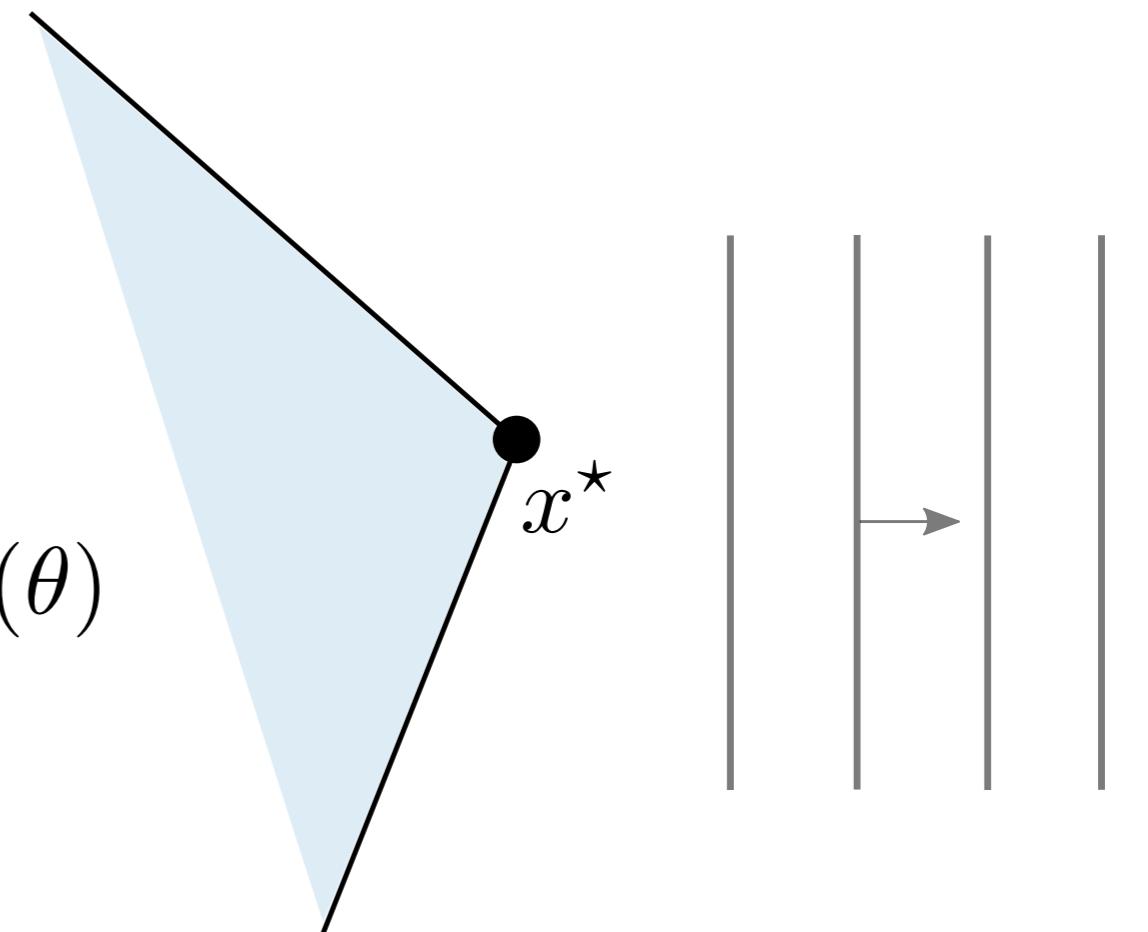
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the optimal solution*

$$s(\theta) = \mathcal{T}(\theta)$$



minimize  $c(\theta)^T x$   
subject to  $A_i(\theta)x = b_i(\theta) \quad \forall i \in \mathcal{T}(\theta)$





# Inventory management

$$\begin{aligned} & \text{minimize} && \sum_{t=0}^{T-1} h(x_t) + o(u_t) \\ & \text{subject to} && x_{t+1} = x_t + u_t - d_t \\ & && x_0 = x_{\text{init}} \\ & && 0 \leq u_t \leq M \end{aligned}$$



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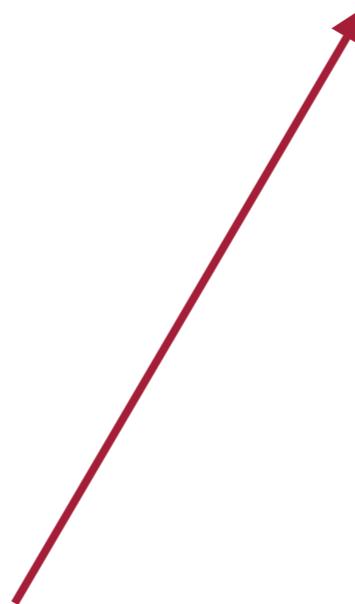
A diagram illustrating the components of the inventory management model. A vertical teal arrow points upwards from the word "Parameters" to the initial condition  $x_0 = x_{\text{init}}$ . A diagonal teal arrow points upwards and to the right from the word "Parameters" to the constraint  $0 \leq u_t \leq M$ .



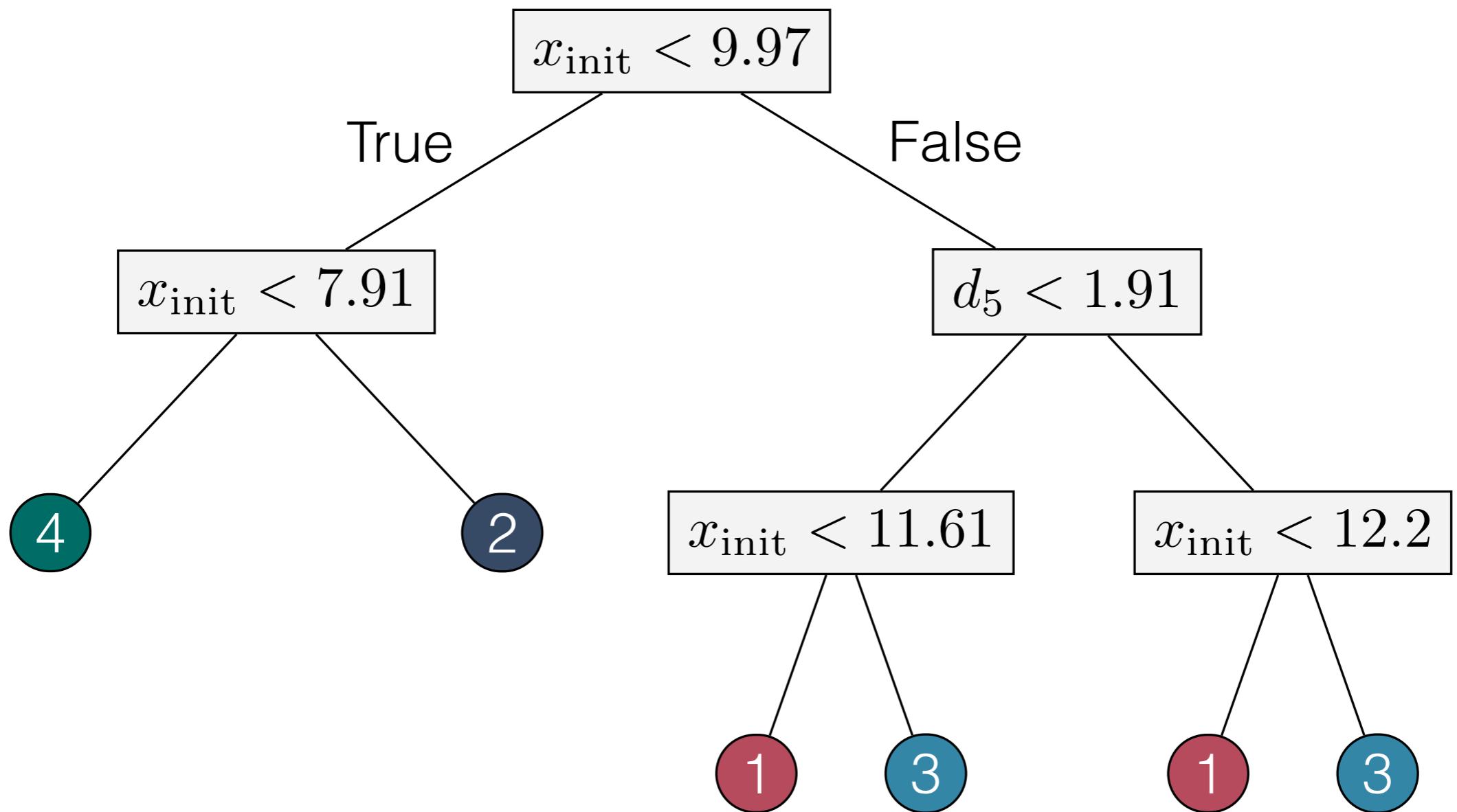
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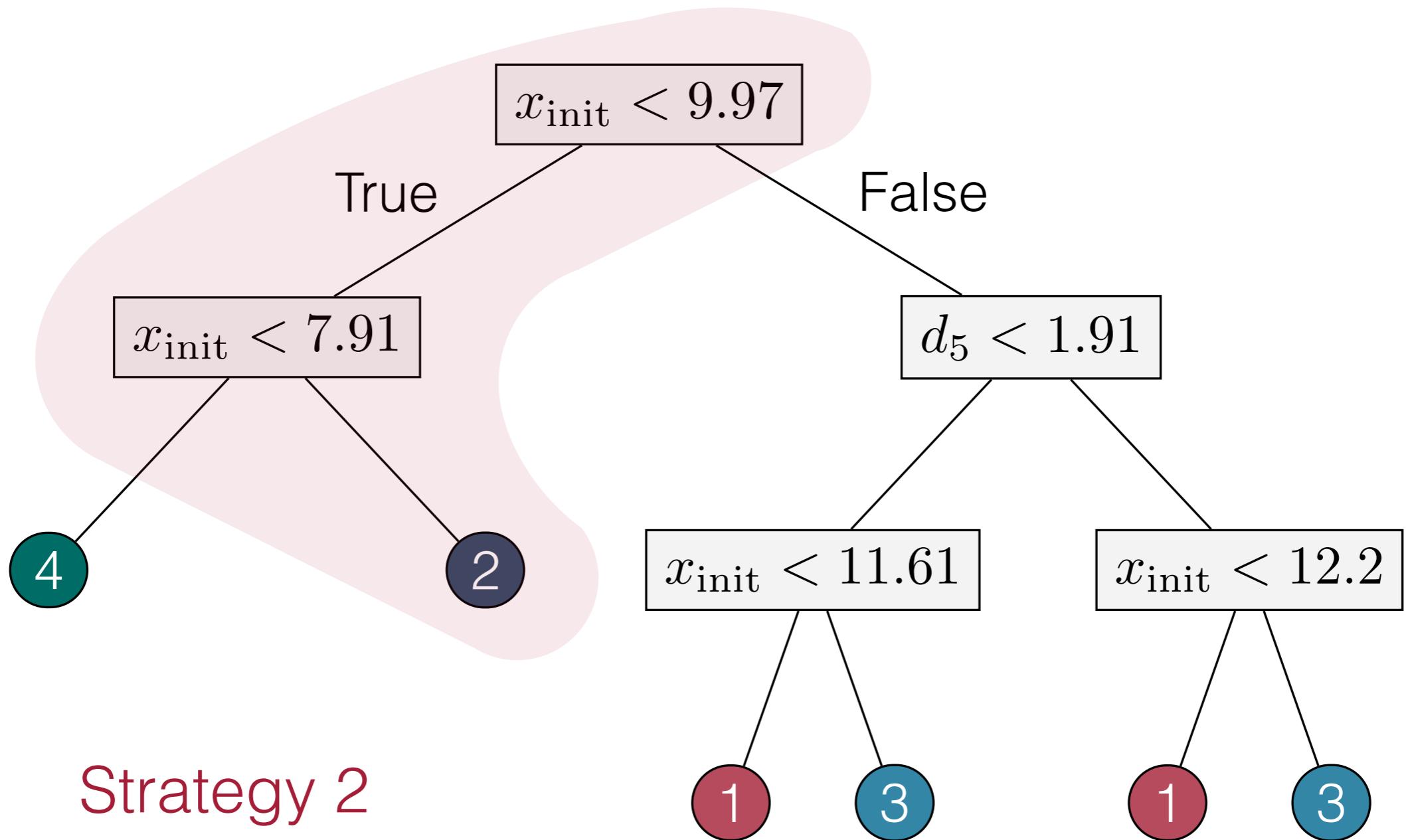
Inequalities



# Strategy selection



# Strategy selection



Strategy 2

$$u_t = 0 \quad t \leq 4$$

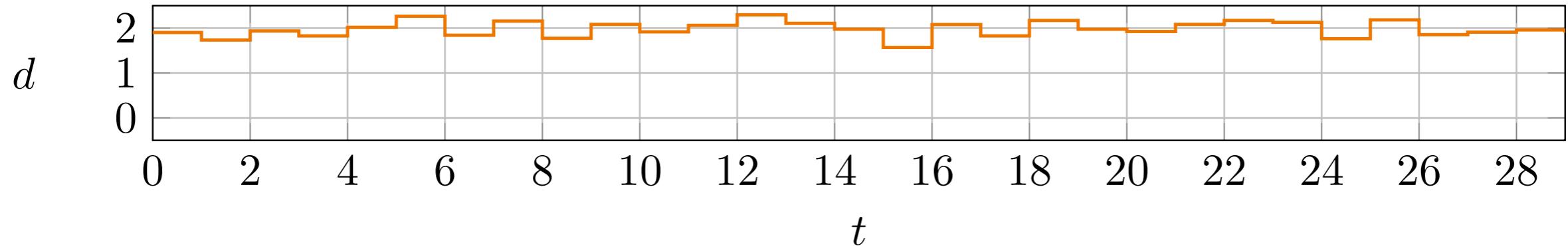
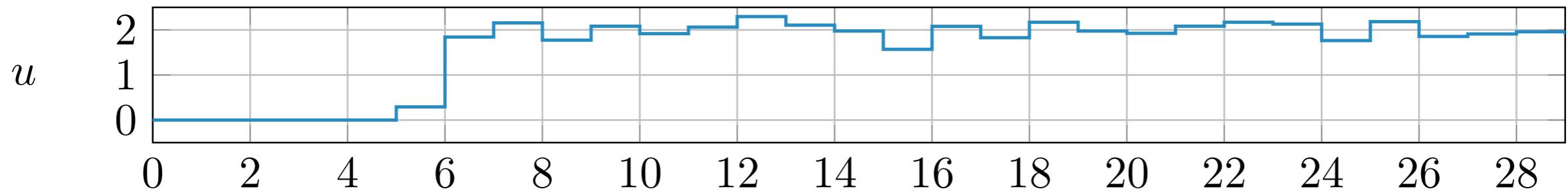
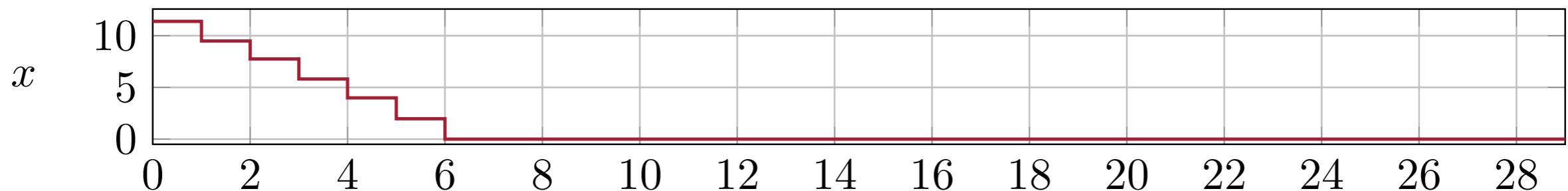
$$0 \leq u_t \leq M \quad t > 4$$

# Strategies for inventory

Strategy 2

$$u_t = 0 \quad t \leq 4$$

$$0 \leq u_t \leq M \quad t > 4$$



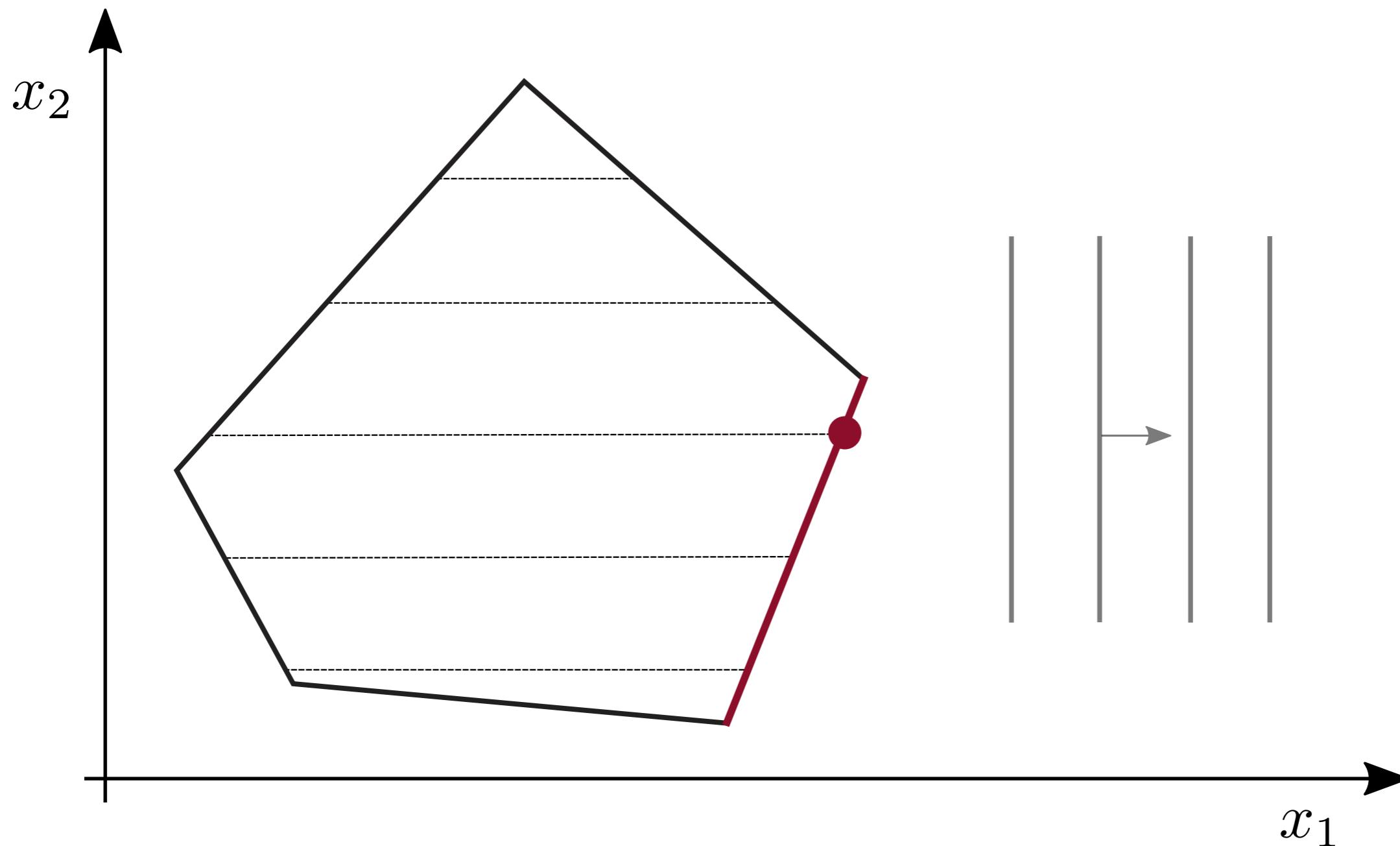
# Mixed-integer optimization

$$\begin{array}{ll}\text{minimize} & c(\theta)^T x \\ \text{subject to} & A(\theta)x \leq b(\theta) \\ & x_{\mathcal{I}} \in \mathbf{Z}^d\end{array}$$

# Mixed-integer optimization

$$\begin{array}{ll}\text{minimize} & c(\theta)^T x \\ \text{subject to} & A(\theta)x \leq b(\theta) \\ & x_{\mathcal{I}} \in \mathbf{Z}^d \quad \text{Integers}\end{array}$$

# Tight constraints are not enough

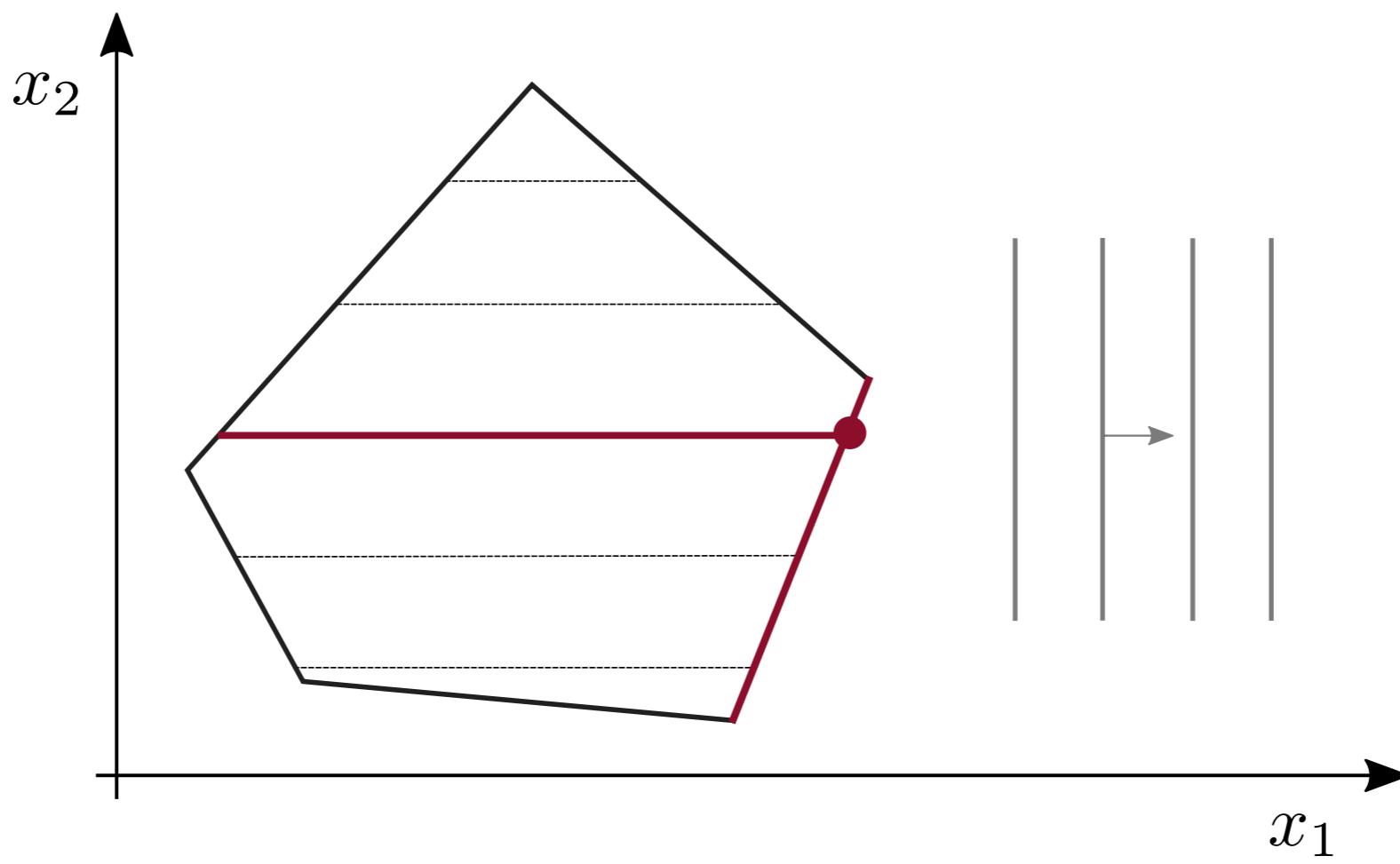


# Strategy

$$s(\theta) = (\mathcal{T}(\theta), x_{\mathcal{I}}^*(\theta))$$



Integer variables





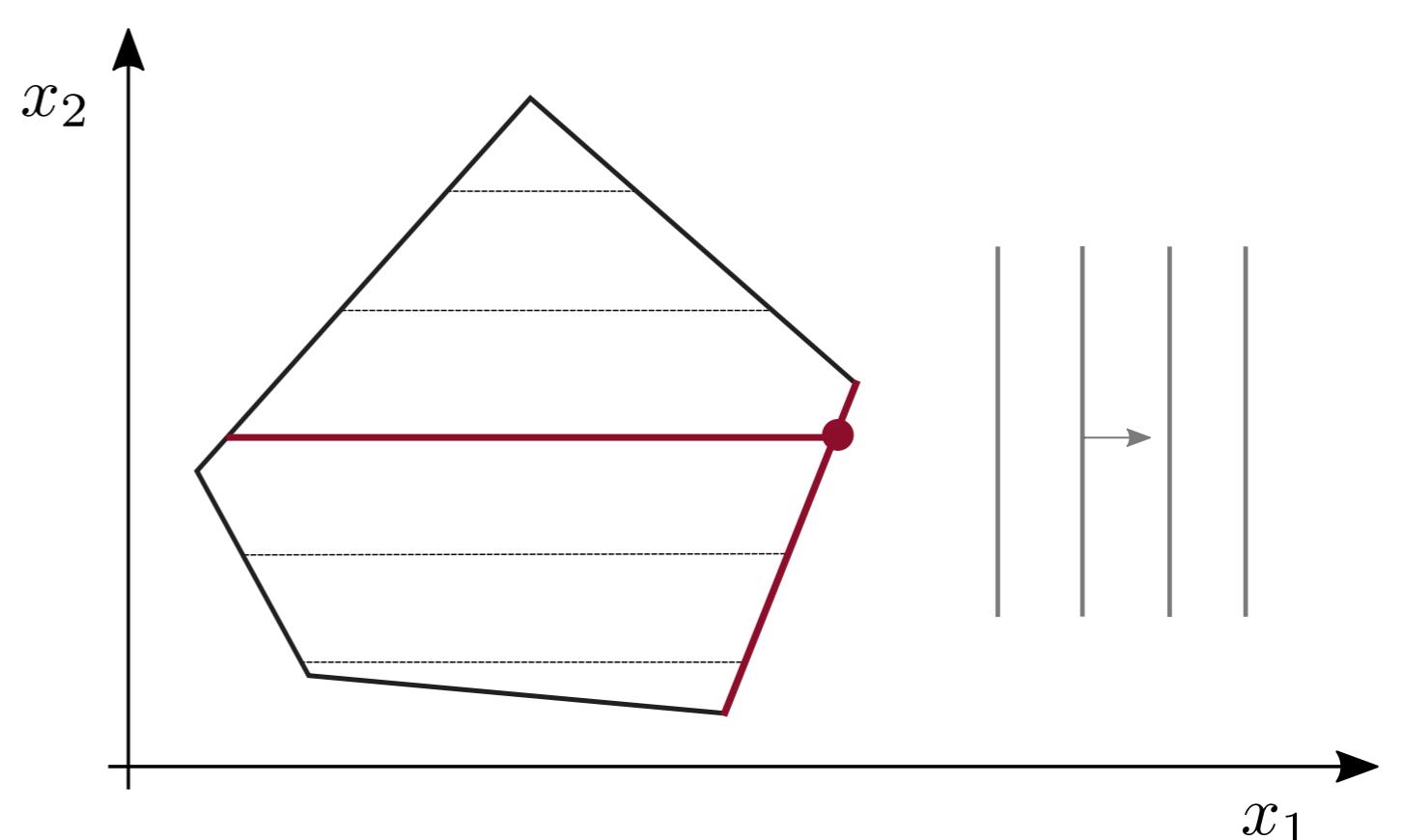
minimize  
subject to

$$c(\theta)^T x$$

$$A_i(\theta)x = b_i(\theta) \quad \forall i \in \mathcal{T}(\theta)$$

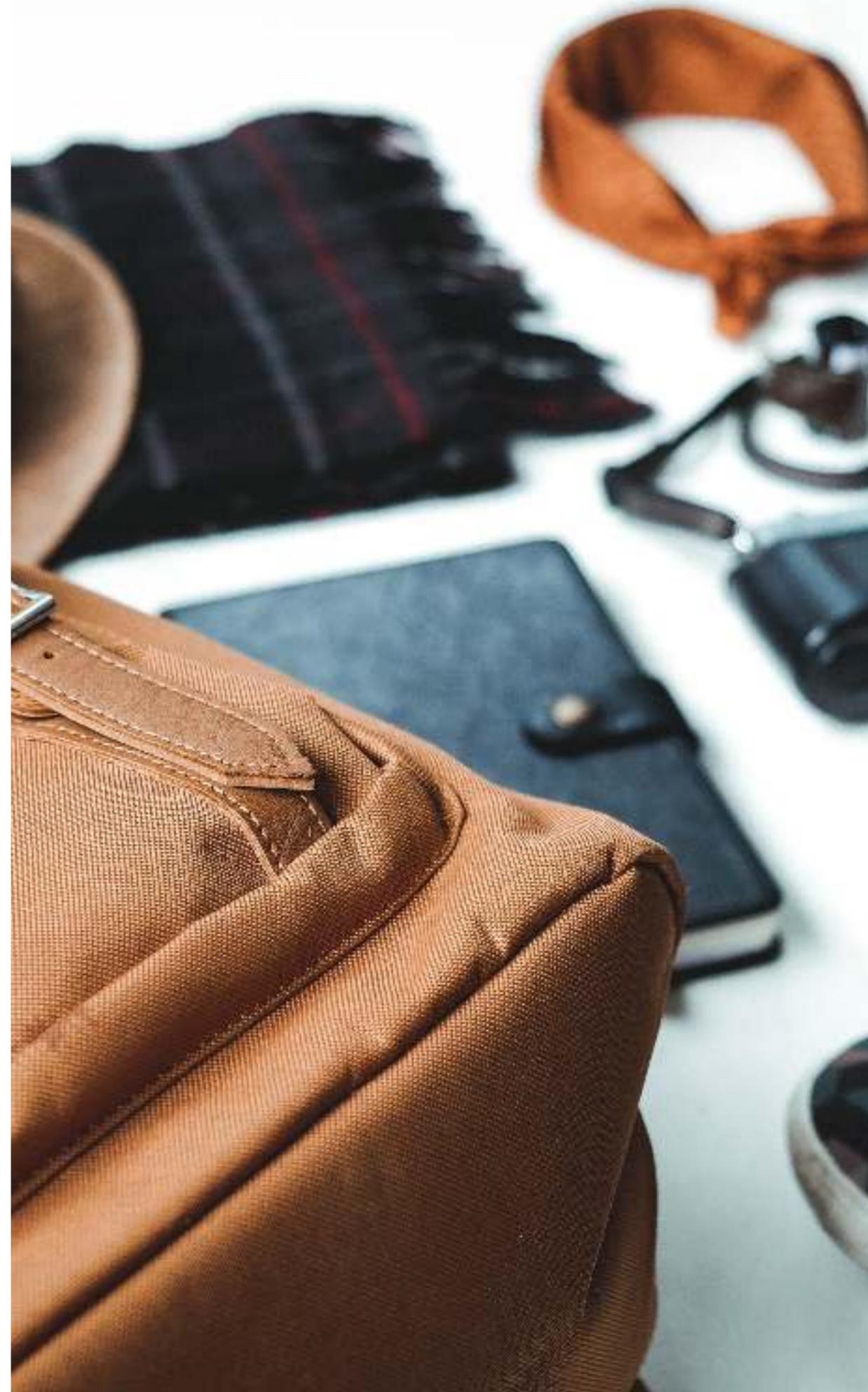
$$x_{\mathcal{I}} = x_{\mathcal{I}}^*(\theta)$$

↑  
Integers



# Knapsack

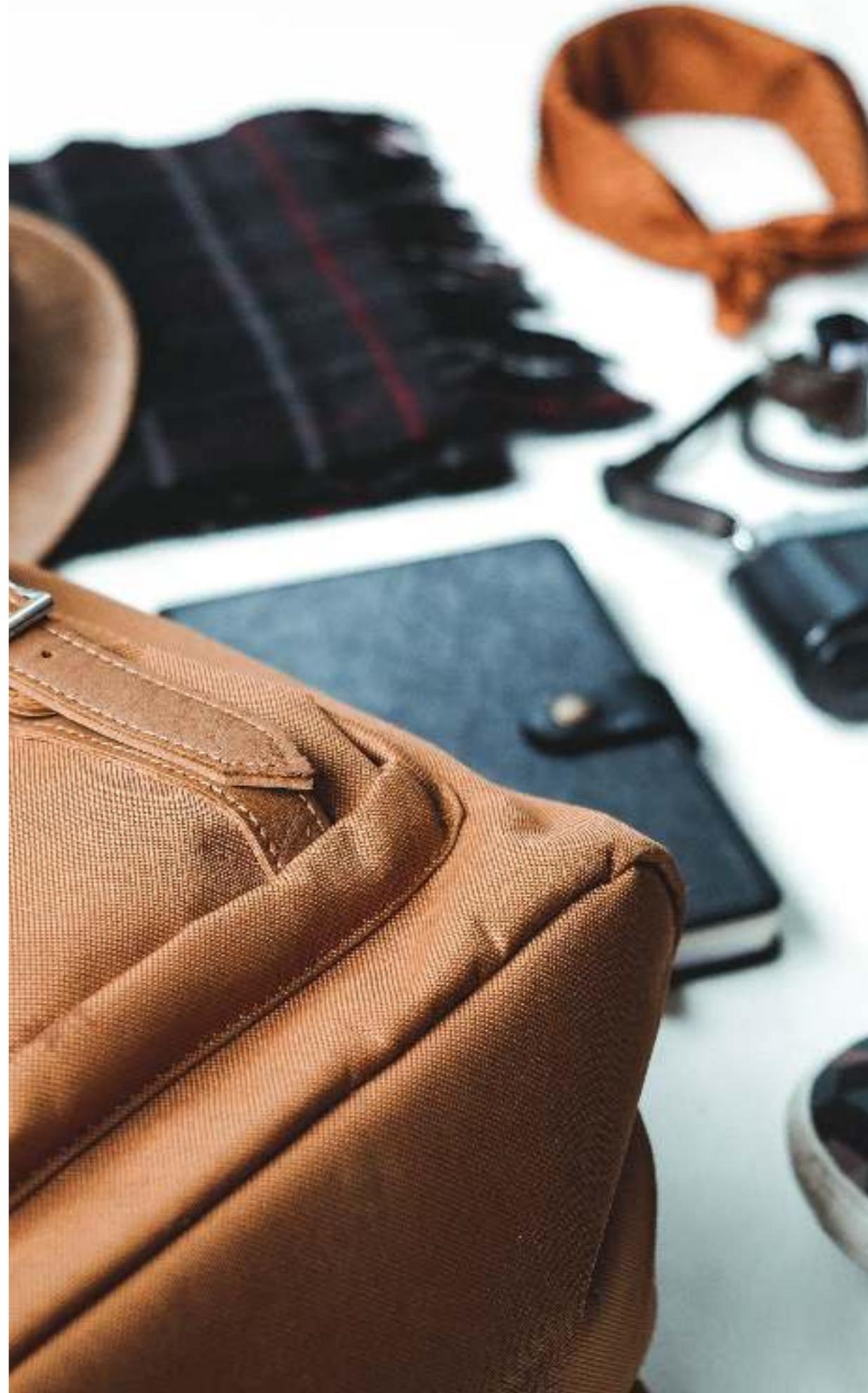
maximize       $c^T x$   
subject to      $a^T x \leq b$   
                   $0 \leq x \leq u$   
                   $x \in \mathbf{Z}^n$



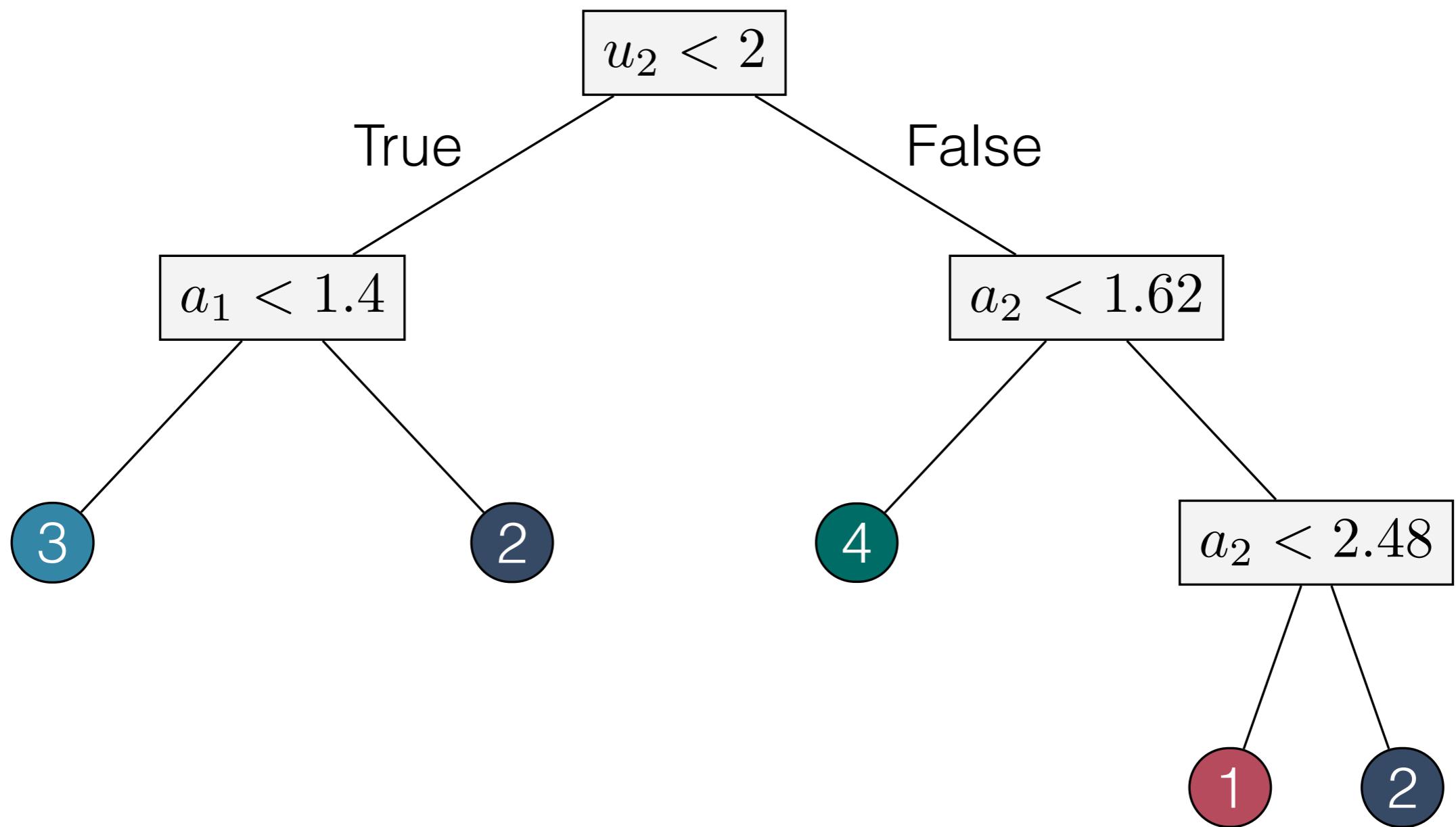
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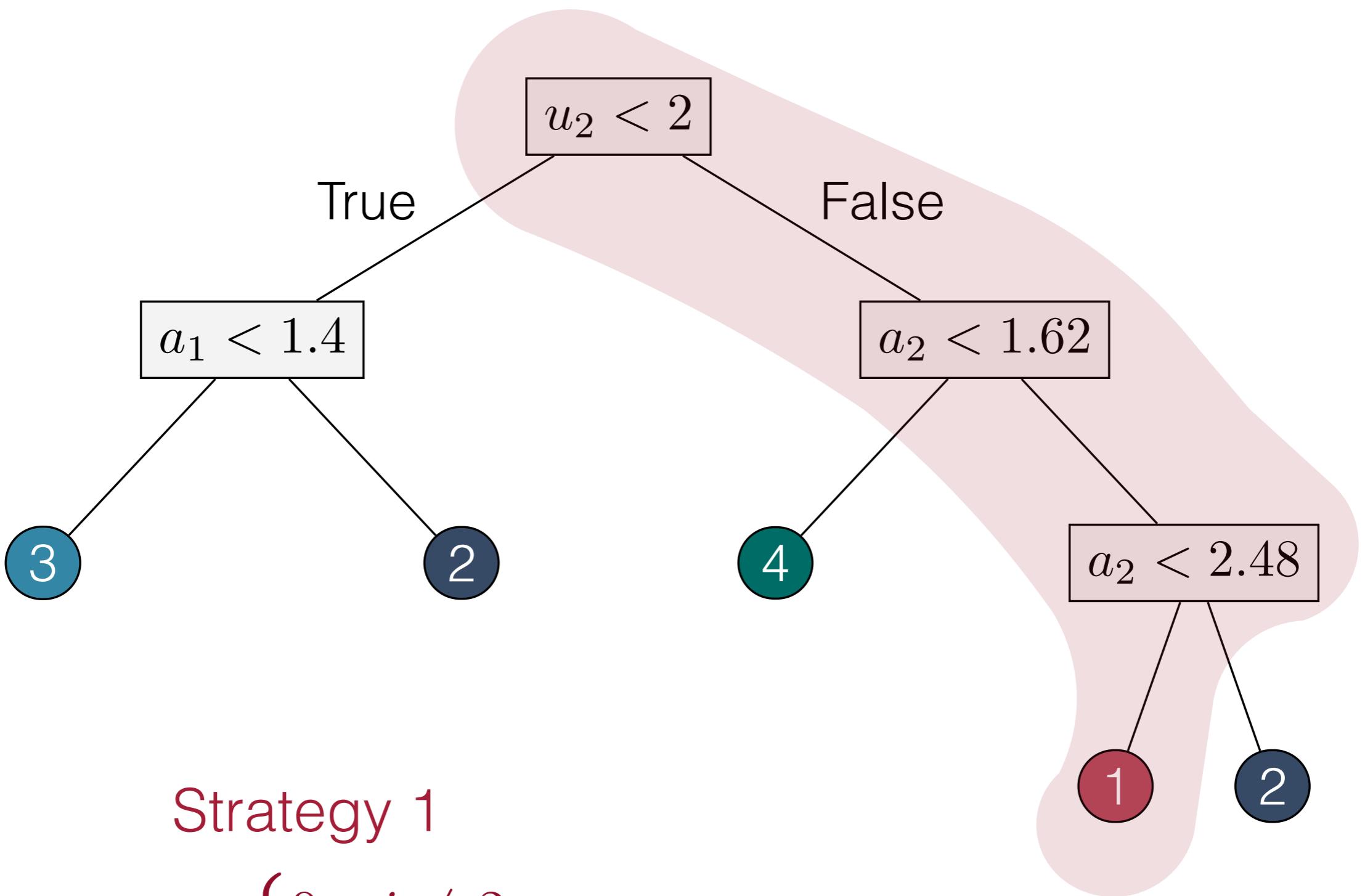
Parameters



# Strategy selection



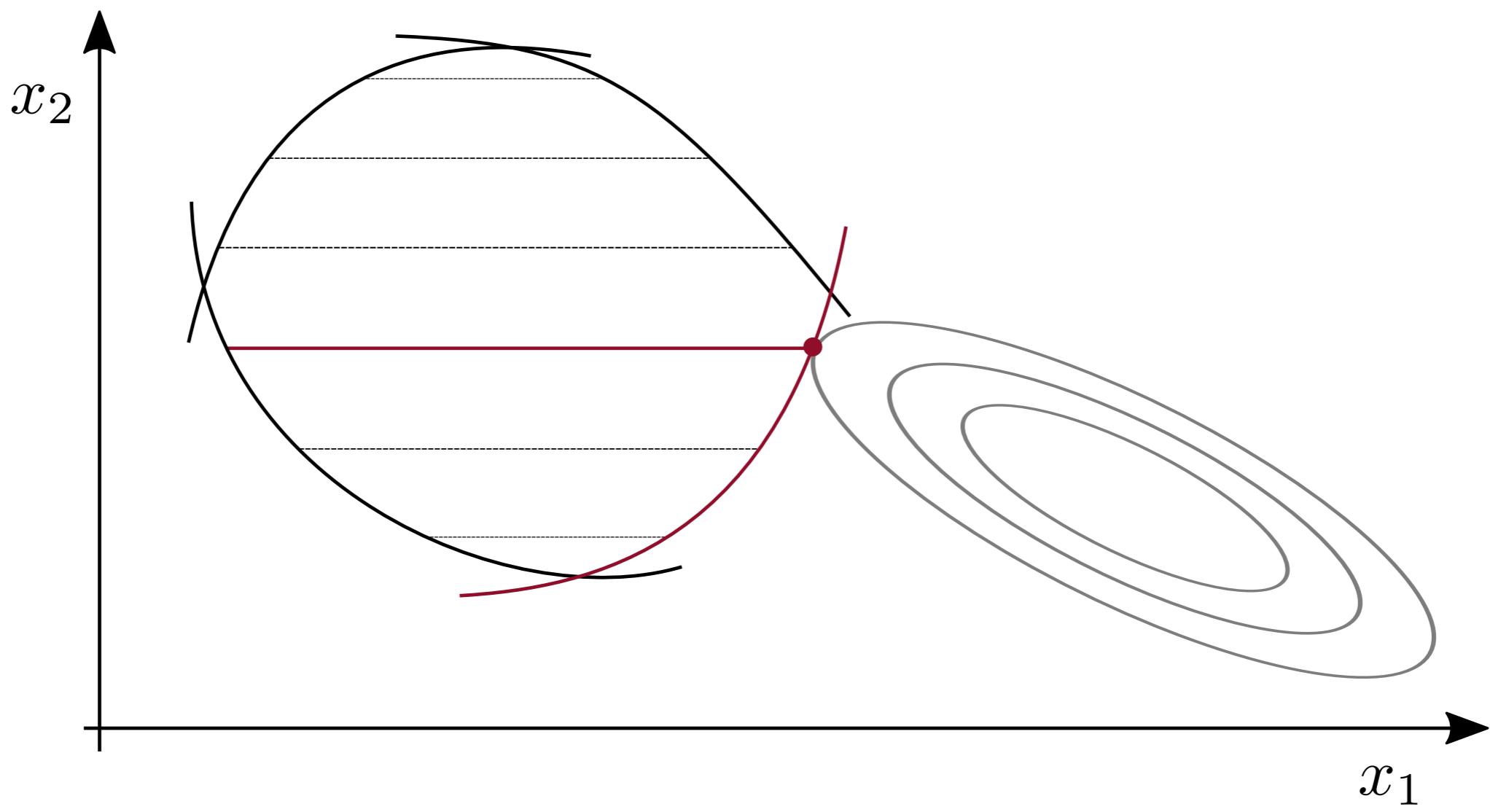
# Strategy selection



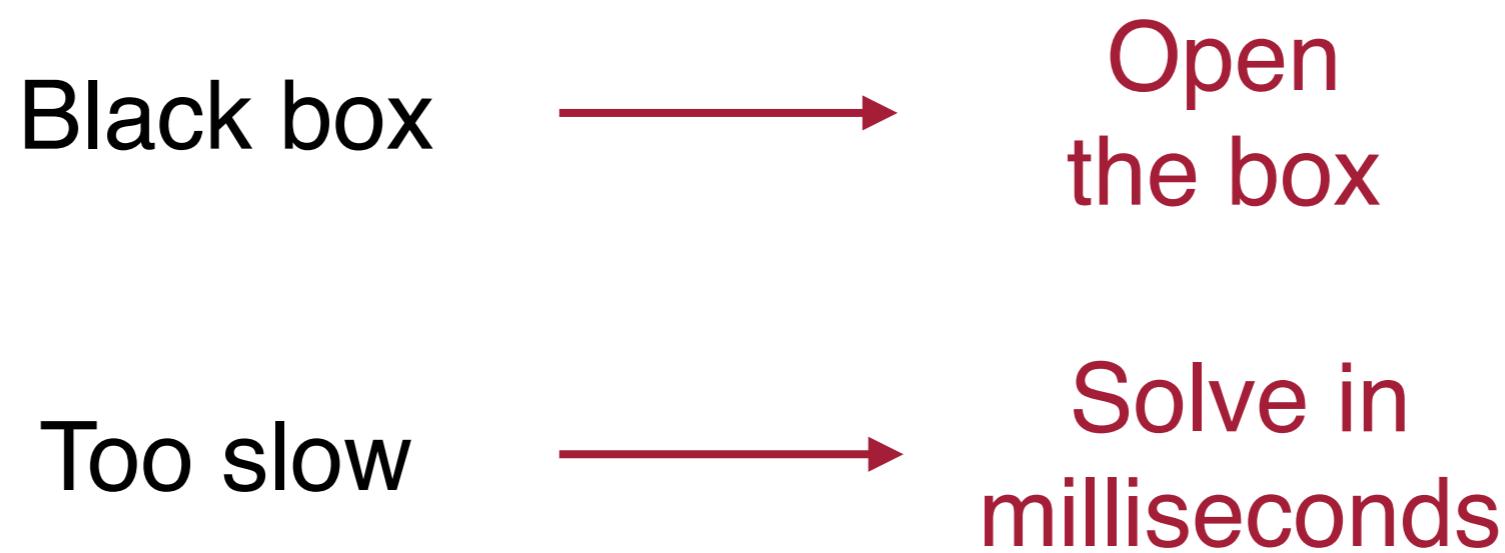
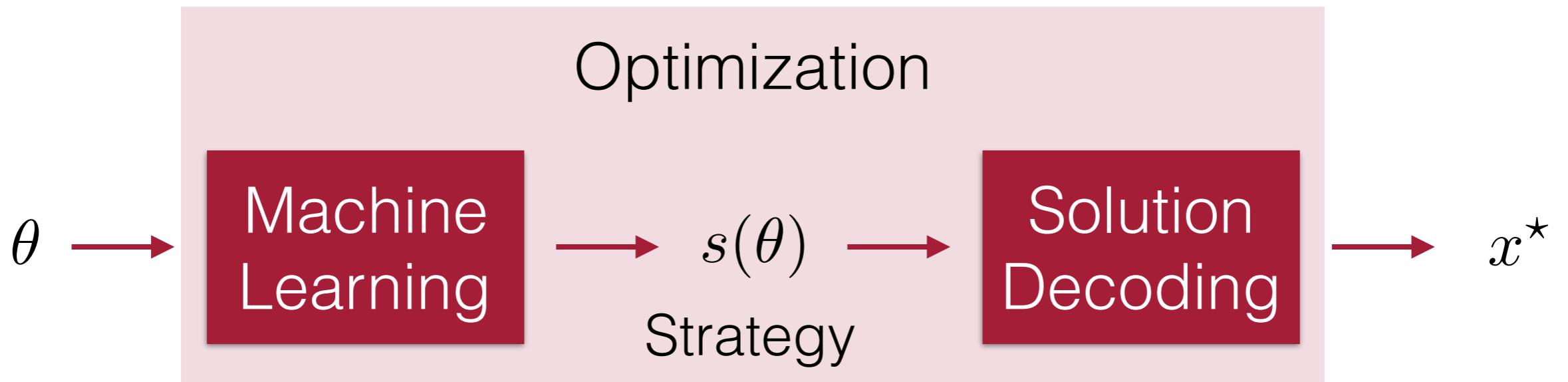
$$x_i = \begin{cases} 0 & i \neq 2 \\ 2 & i = 2 \end{cases}$$

# Mixed-integer convex optimization

$$\begin{array}{ll}\text{minimize} & f(\theta, x) \\ \text{subject to} & g(\theta, x) \leq 0 \\ & x_{\mathcal{I}} \in \mathbf{Z}^d\end{array}$$



# Learning the strategies



# Classification problem

$N$  data

$(\theta_i, s(\theta_i))$

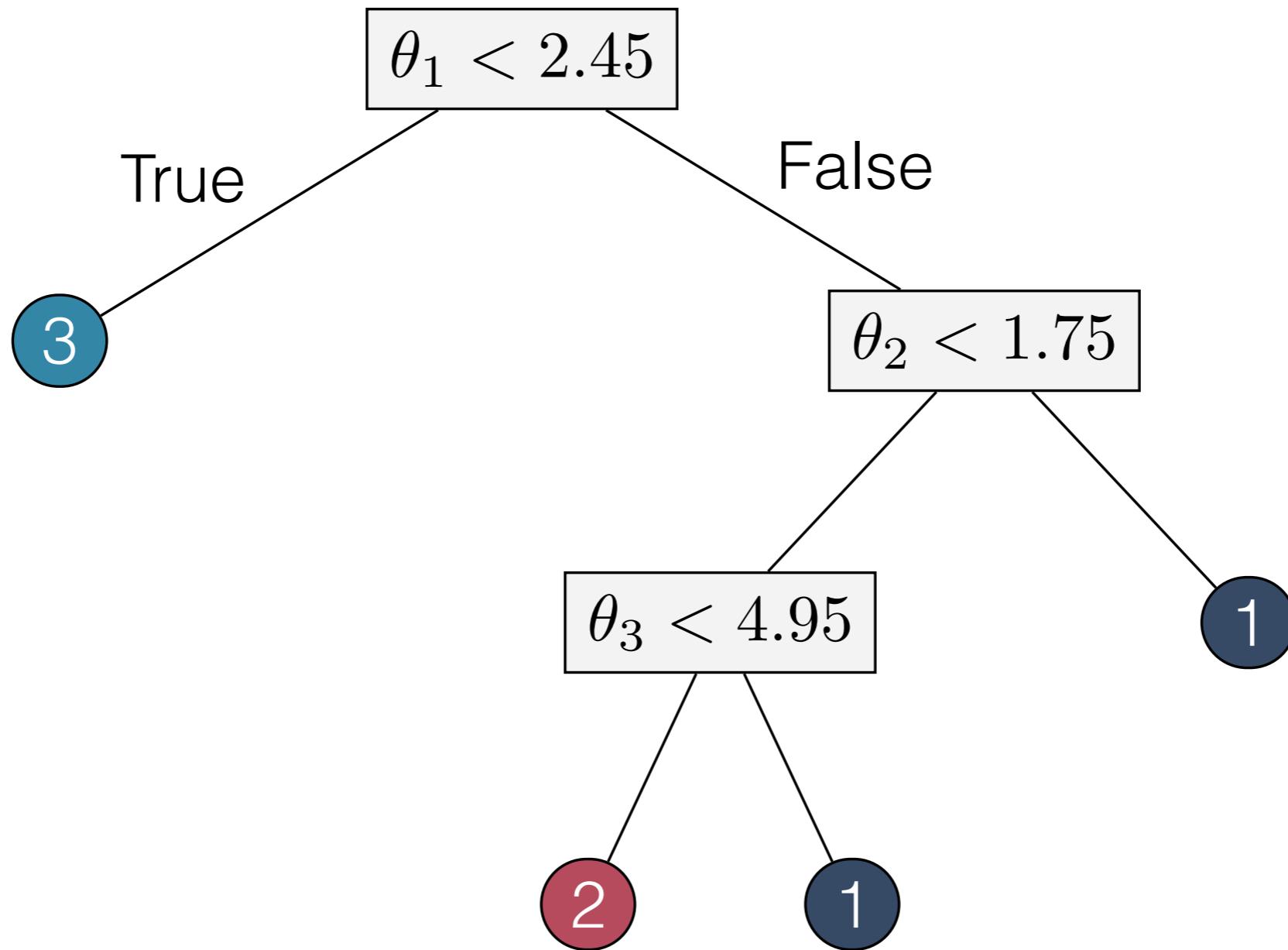
$M$  labels

Strategies  $\mathcal{S}$

Learn multiclass classifier

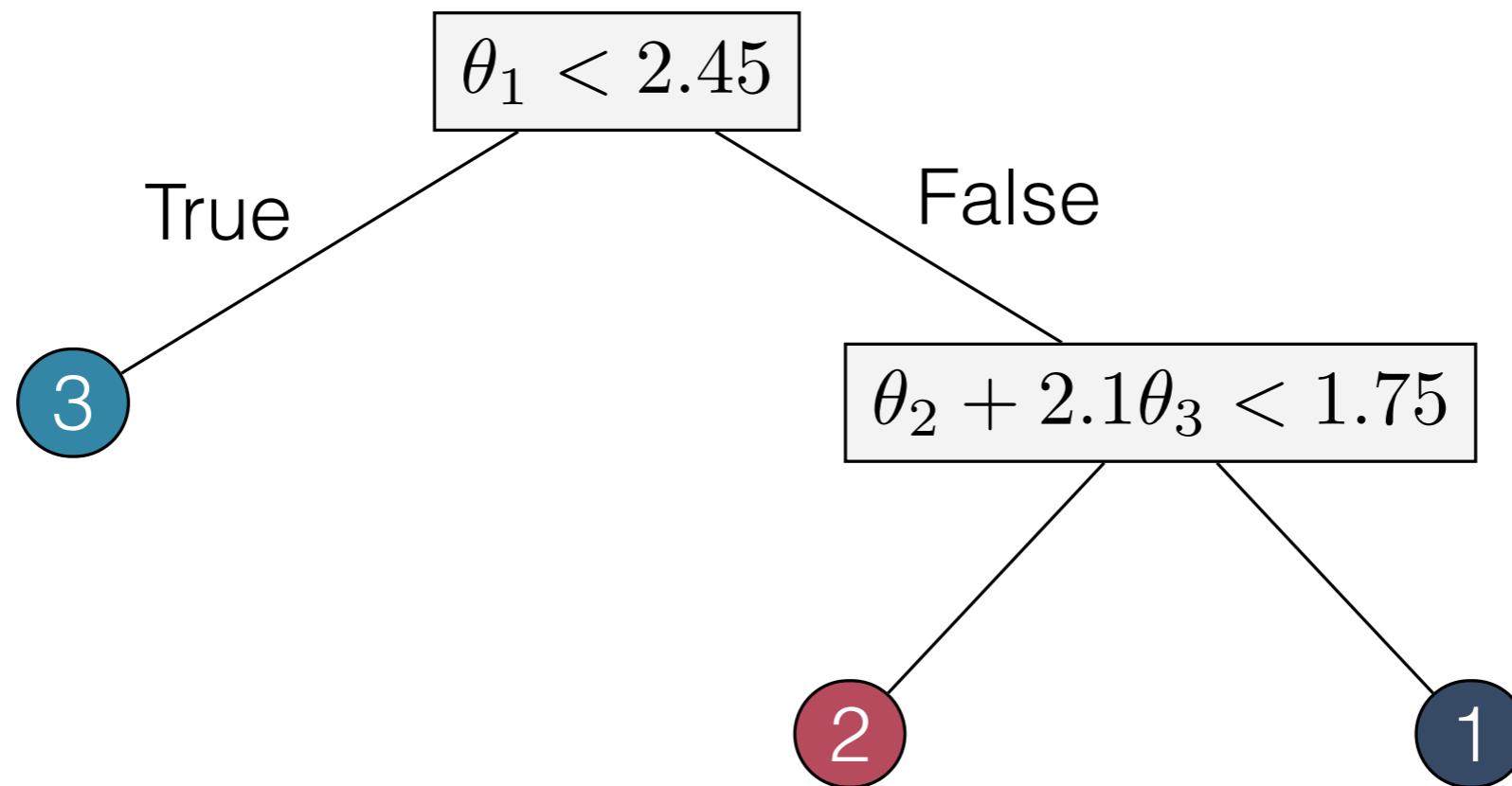


# Optimal Classification Trees (OCT)



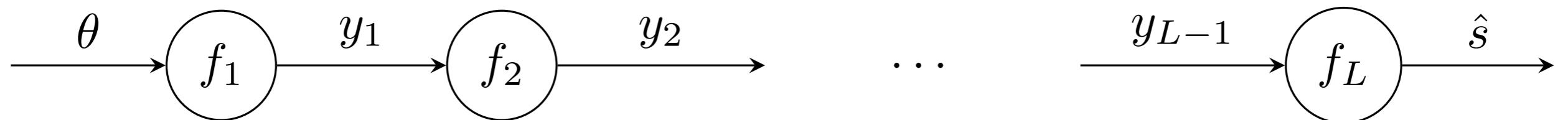
**Interpretable AI**  
[Bertsimas and Dunn (2017)]

# Optimal Classification Trees with Hyperplanes (OCT-H)



Interpretable AI  
[Bertsimas and Dunn (2017)]

# Neural Networks (NN)



Single layer

$$y_l = f(y_{l-1}) = (W_l y_{l-1} + b_l)_+$$

 PyTorch

# Strategies exploration

# Classification problem

$N$  data

$(\theta_i, s(\theta_i))$

$M$  labels

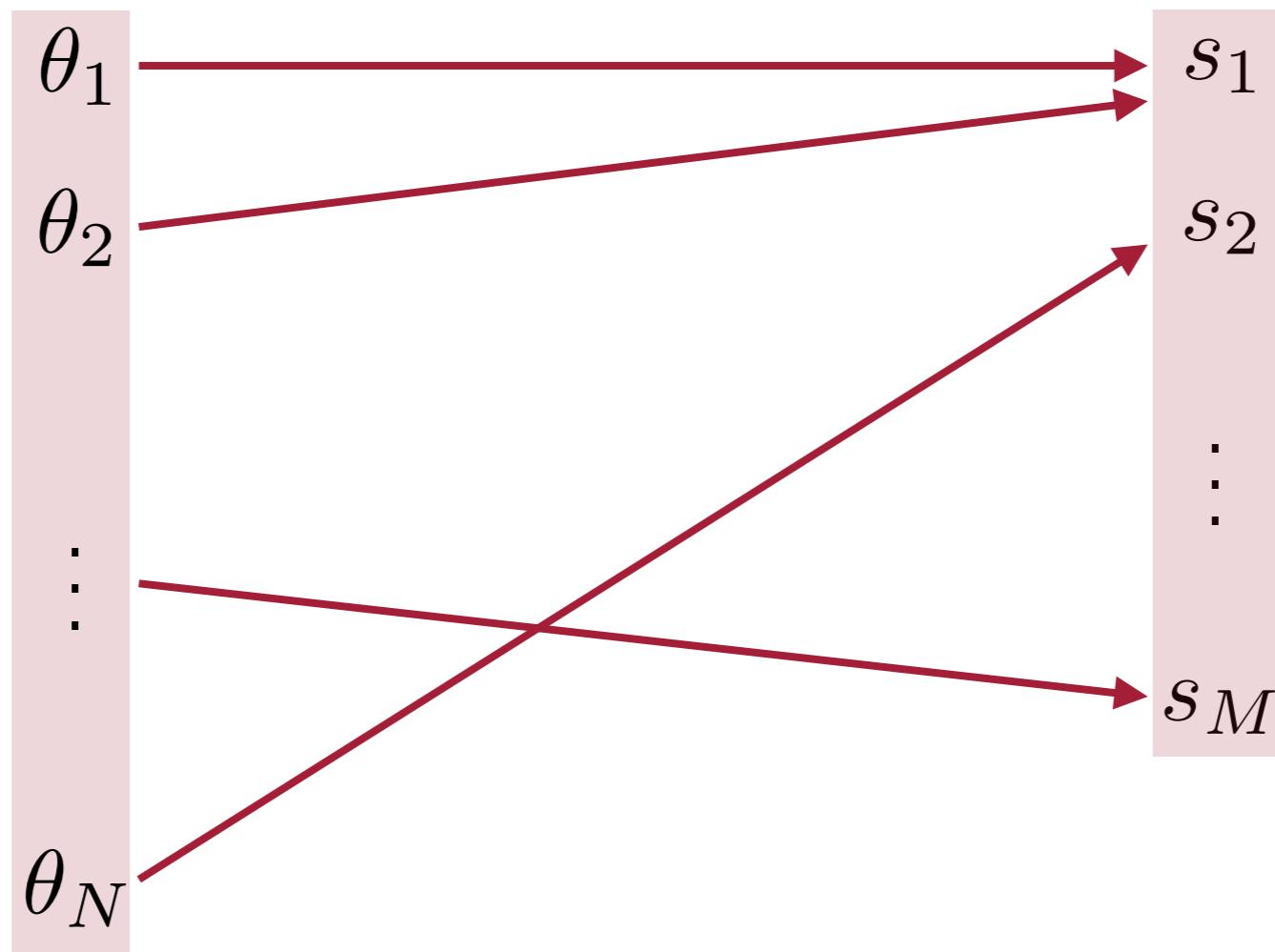
Strategies  $\mathcal{S}$



Have we seen enough data?

# Will we find new strategies?

Parameters



$\theta_{N+1}?$

Alan Turing already knew that

Decoded the Enigma Machine  
in  
World War II

Only some words (labels) needed



# Good-Turing estimator

$s_1$	12 times
$s_2$	45 times
:	:
$s_M$	2 times

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$s_1$	12 times
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Probability of unseen strategies

$$GT = \frac{N_1}{N}$$

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Probability of unseen strategies

$$GT = \frac{N_1}{N} \quad \text{\# strategies appeared once}$$

# Good-Turing estimator

$s_1$	12 times
$s_2$	45 times
:	:
$s_M$	2 times

Probability of unseen strategies

$$GT = \frac{N_1}{N} \quad \begin{array}{l} \text{\# strategies appeared once} \\ \text{\# samples} \end{array}$$

# Sampling scheme

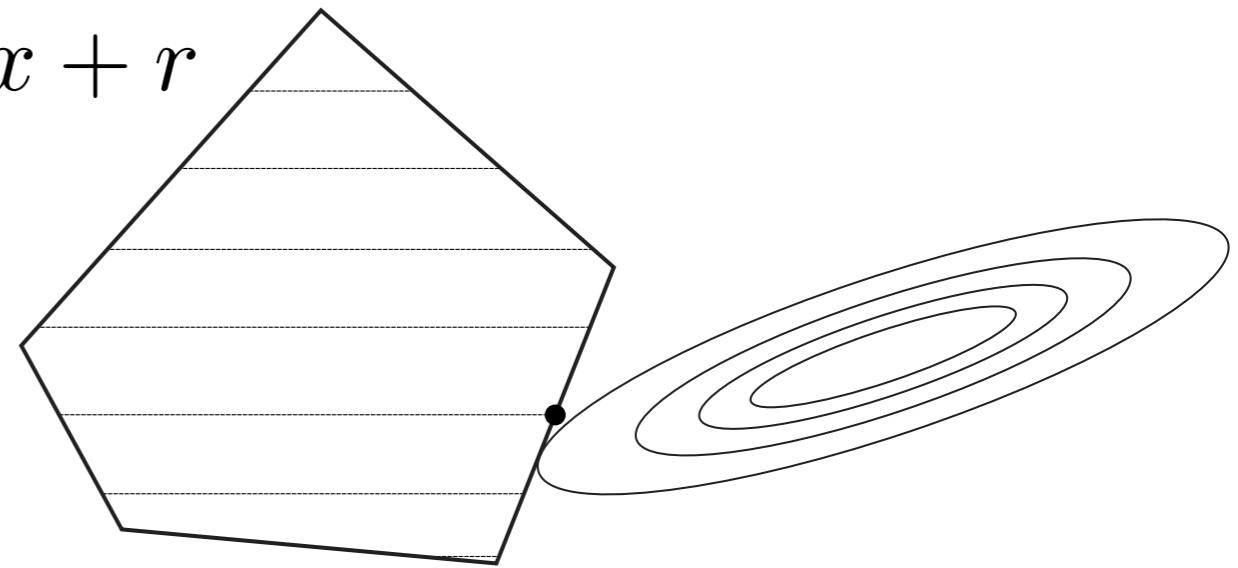
Repeat until  $GT \leq \epsilon$

1. sample  $\theta_i$
2. compute  $s(\theta_i)$
3. update estimator  $GT = \frac{N_1}{N}$

# Speedups

# Mixed-integer quadratic optimization

$$\begin{array}{ll}\text{minimize} & (1/2)x^T Px + q^T x + r \\ \text{subject to} & Ax \leq b \\ & x_{\mathcal{I}} \in \mathbf{Z}^d\end{array}$$

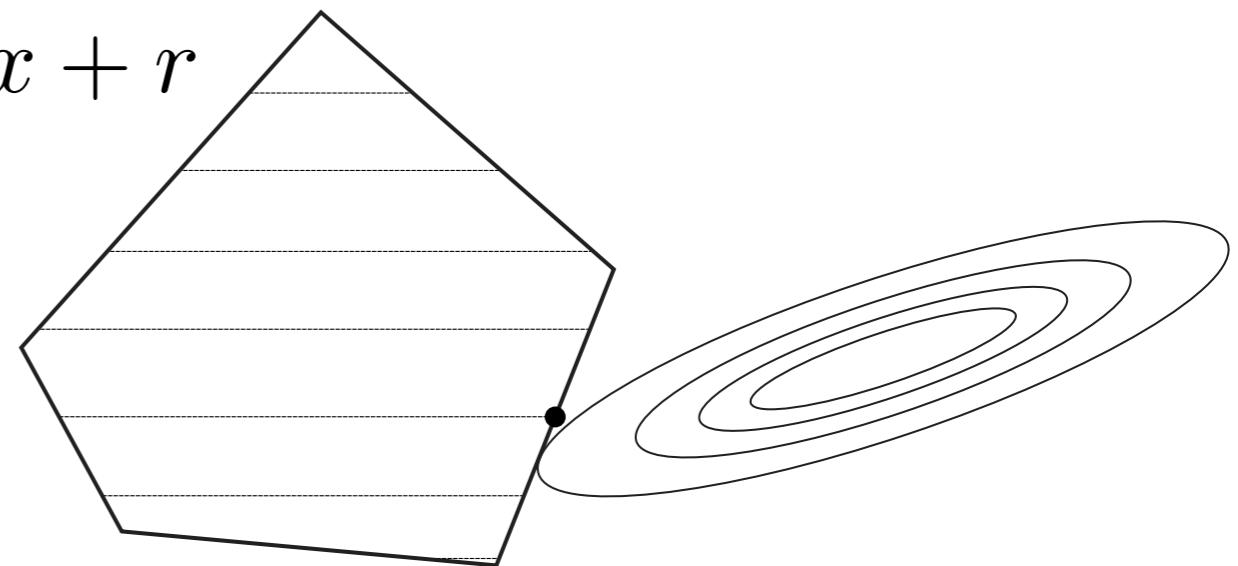


# Mixed-integer quadratic optimization

$$\begin{array}{ll}\text{minimize} & (1/2)x^T Px + q^T x + r \\ \text{subject to} & Ax \leq b \\ & x_{\mathcal{I}} \in \mathbf{Z}^d\end{array}$$



Strategy



$$\begin{array}{ll}\text{minimize} & (1/2)x^T Px + q^T x + r \\ \text{subject to} & A_{\mathcal{T}(\theta)}x = b_{\mathcal{T}(\theta)} \\ & x_{\mathcal{I}} = x_{\mathcal{I}}^*(\theta)\end{array}$$

Equality  
Constrained  
QP

# Quick solution

## KKT System

$$\begin{bmatrix} P & A_{\mathcal{T}(\theta)}^T & I_{\mathcal{I}}^T \\ A_{\mathcal{T}(\theta)} & 0 & \\ I_{\mathcal{I}} & & \end{bmatrix} \begin{bmatrix} x \\ \nu \end{bmatrix} = \begin{bmatrix} -q \\ b_{\mathcal{T}(\theta)} \\ x_{\mathcal{I}}^*(\theta) \end{bmatrix}$$

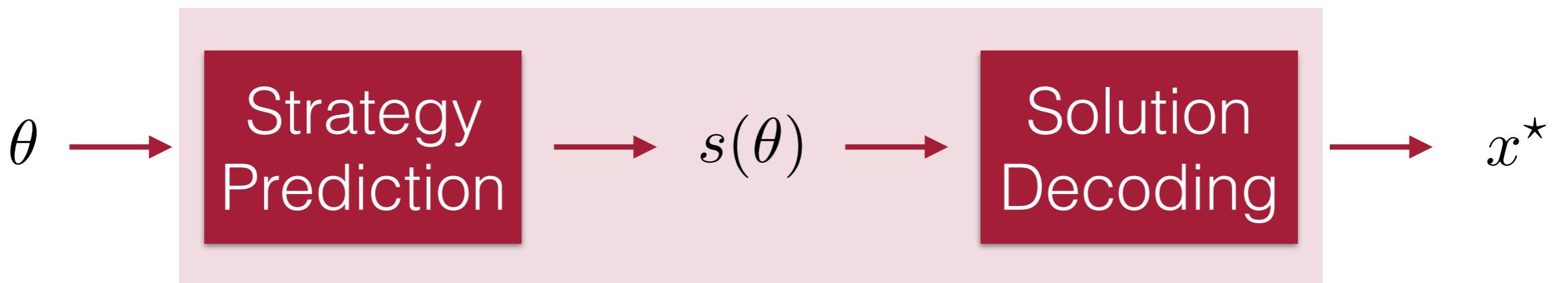


# MLOPT

## Offline Learning

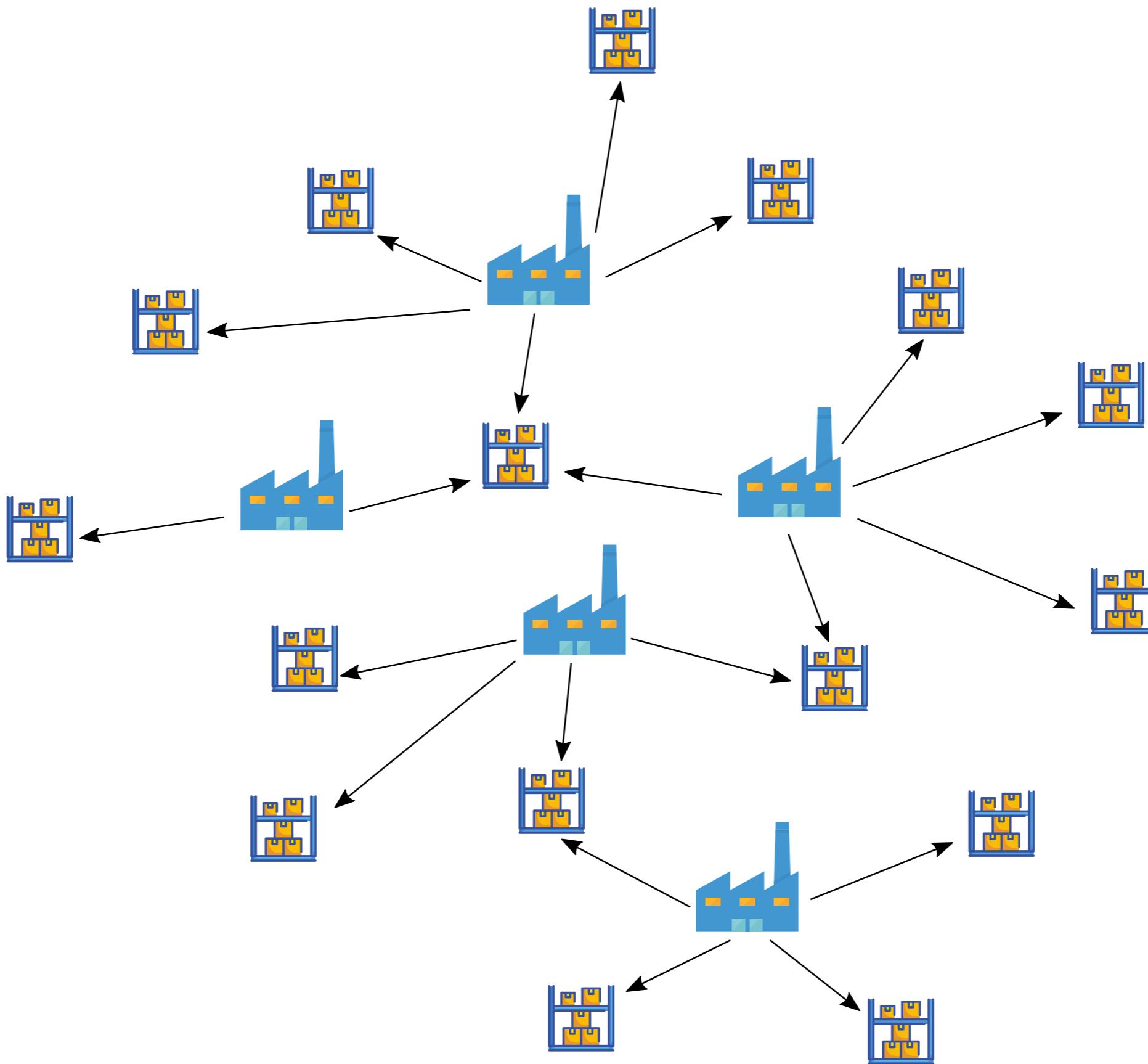


## Online Optimization



# Benchmarks

# Facility location



# Facility location

$$\text{minimize} \quad \sum_{i \in F} \sum_{j \in W} c_{ij} x_{ij} + \sum_{i \in I} f_i y_i$$

$$\begin{aligned} \text{subject to} \quad & \sum_{i \in F} x_{ij} \geq d_j \quad \forall j \in W \\ & \sum_{j \in W} x_{ij} \leq s_i y_i, \quad \forall i \in F \\ & x_{ij} \geq 0, y_i \in \{0, 1\} \end{aligned}$$

# Facility location

$$\begin{array}{ll}\text{minimize} & \sum_{i \in F} \sum_{j \in W} c_{ij} x_{ij} + \sum_{i \in I} f_i y_i \\ \text{subject to} & \sum_{i \in F} x_{ij} \geq d_j \quad \forall j \in W \\ & \sum_{j \in W} x_{ij} \leq s_i y_i, \quad \forall i \in F \\ & x_{ij} \geq 0, y_i \in \{0, 1\}\end{array}$$

↑  
Facility

Facility  
cost

# Facility location

$$\text{minimize} \quad \sum_{i \in F} \sum_{j \in W} \text{Shipment cost } c_{ij} x_{ij} + \sum_{i \in I} \text{Facility cost } f_i y_i$$

$$\begin{aligned} \text{subject to} \quad & \sum_{i \in F} x_{ij} \geq d_j \quad \forall j \in W \\ & \sum_{j \in W} x_{ij} \leq s_i y_i, \quad \forall i \in F \\ & x_{ij} \geq 0, y_i \in \{0, 1\} \end{aligned}$$

Shipment  $\nearrow$  Facility

# Facility location

$$\begin{array}{ll} \text{minimize} & \text{Shipment cost} + \text{Facility cost} \\ & \sum_{i \in F} \sum_{j \in W} c_{ij} x_{ij} + \sum_{i \in I} f_i y_i \\ \text{subject to} & \sum_{i \in F} x_{ij} \geq d_j \quad \forall j \in W \quad \text{Demand} \\ & \sum_{j \in W} x_{ij} \leq s_i y_i, \quad \forall i \in F \\ & x_{ij} \geq 0, y_i \in \{0, 1\} \end{array}$$

Shipment

Facility

A red arrow points from the word "Shipment" to the term  $c_{ij}x_{ij}$  in the objective function. A blue arrow points from the word "Facility" to the term  $f_iy_i$  in the objective function.

# Facility location

$$\begin{array}{ll} \text{minimize} & \text{Shipment cost} + \text{Facility cost} \\ & \sum_{i \in F} \sum_{j \in W} c_{ij} x_{ij} + \sum_{i \in I} f_i y_i \\ \text{subject to} & \sum_{i \in F} x_{ij} \geq d_j \quad \forall j \in W \\ & \sum_{j \in W} x_{ij} \leq s_i y_i, \quad \forall i \in F \\ & x_{ij} \geq 0, y_i \in \{0, 1\} \end{array}$$

Demand Supply

Shipment Facility

# Facility location

minimize      Shipment cost      Facility cost

$$\sum_{i \in F} \sum_{j \in W} c_{ij} x_{ij} + \sum_{i \in I} f_i y_i$$

subject to

$$\sum_{i \in F} x_{ij} \geq d_j \quad \forall j \in W$$

Demand

$$\sum_{j \in W} x_{ij} \leq s_i y_i, \quad \forall i \in F$$

Supply

Parameters

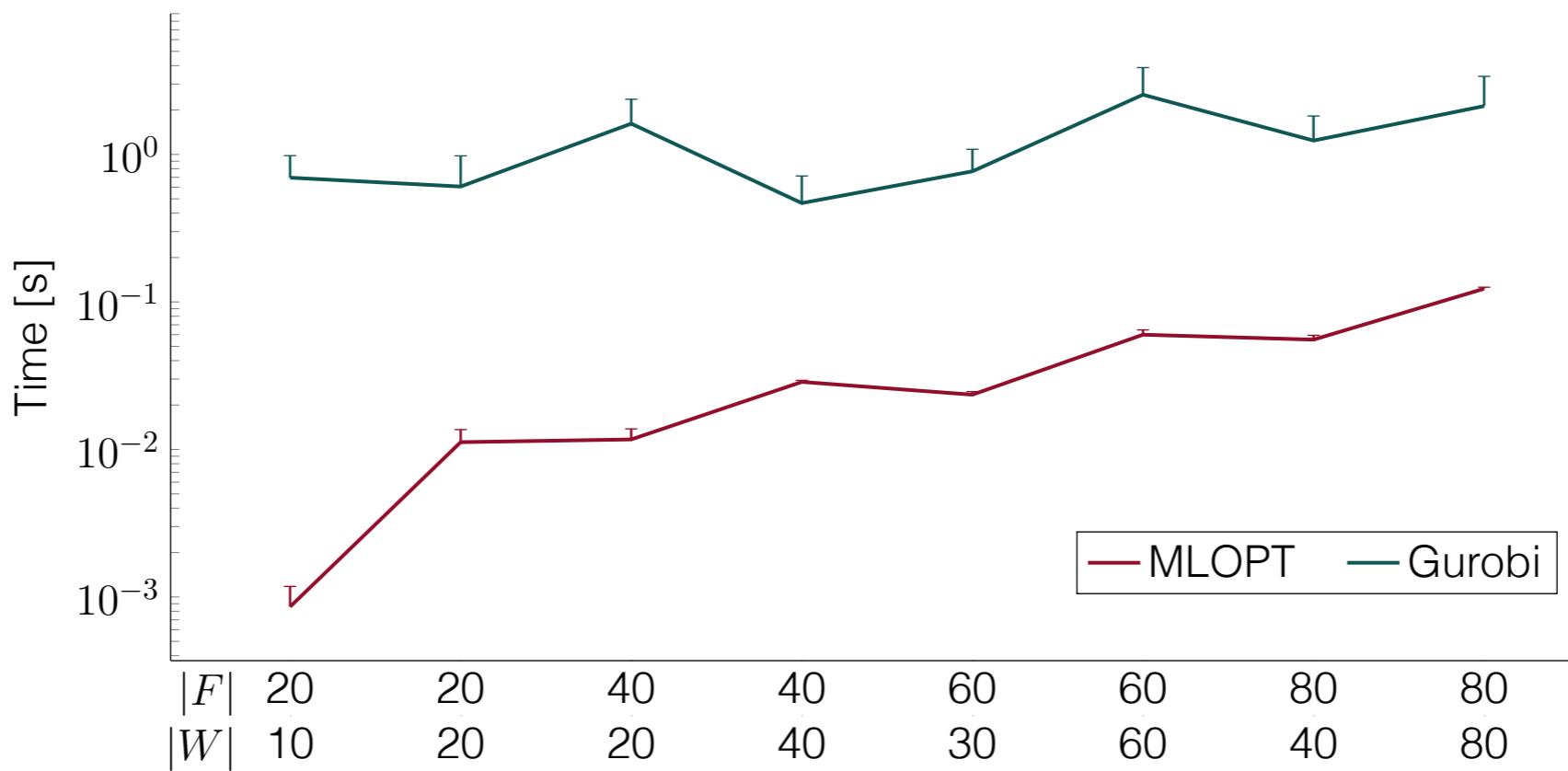
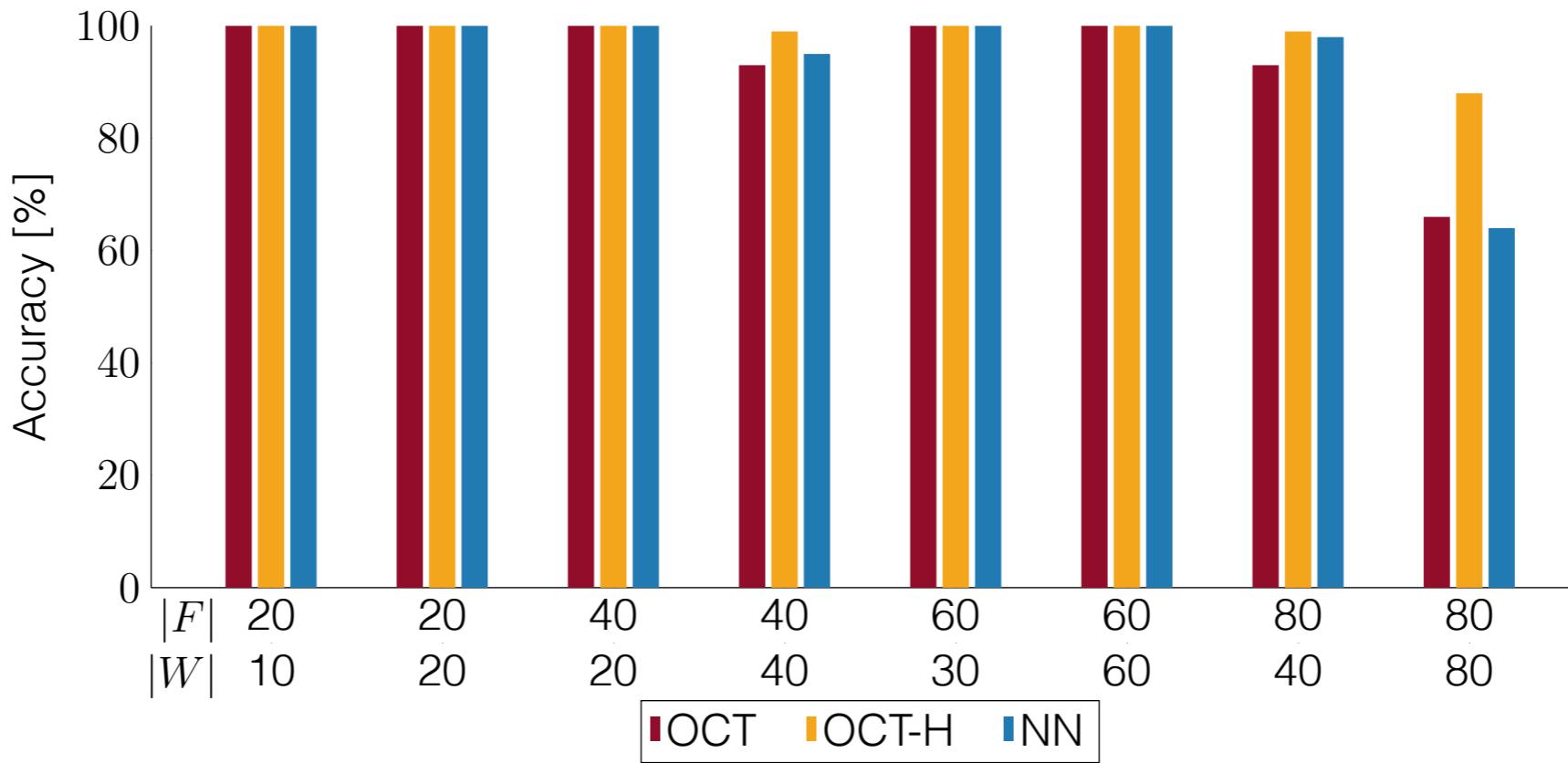
$$x_{ij} \geq 0, y_i \in \{0, 1\}$$

Shipment

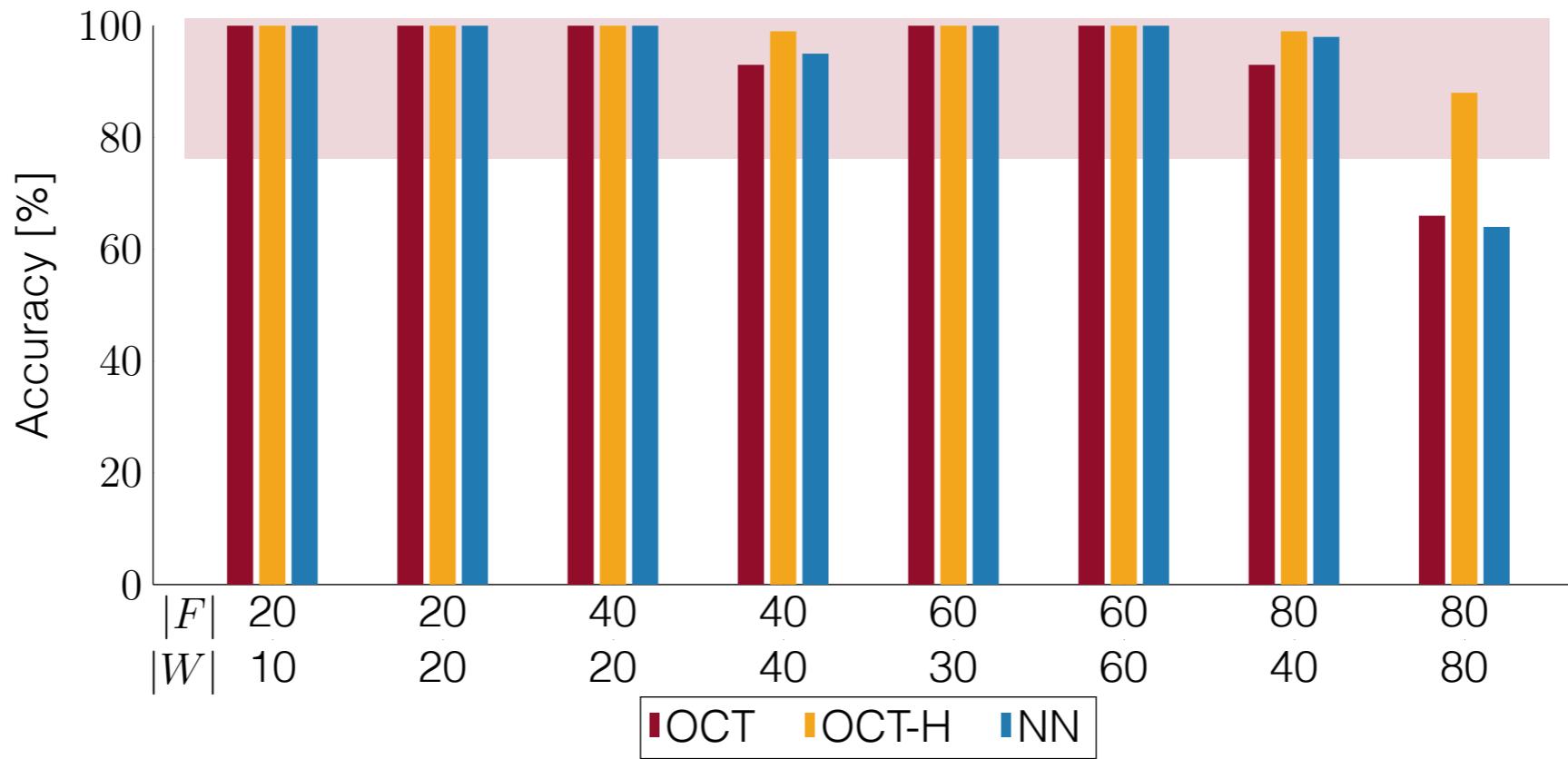
Facility

```
graph TD; Parameters -- green arrow --> F_Sum; Shipment -- red arrow --> Cij_xij; Facility -- blue arrow --> fi_yi;
```

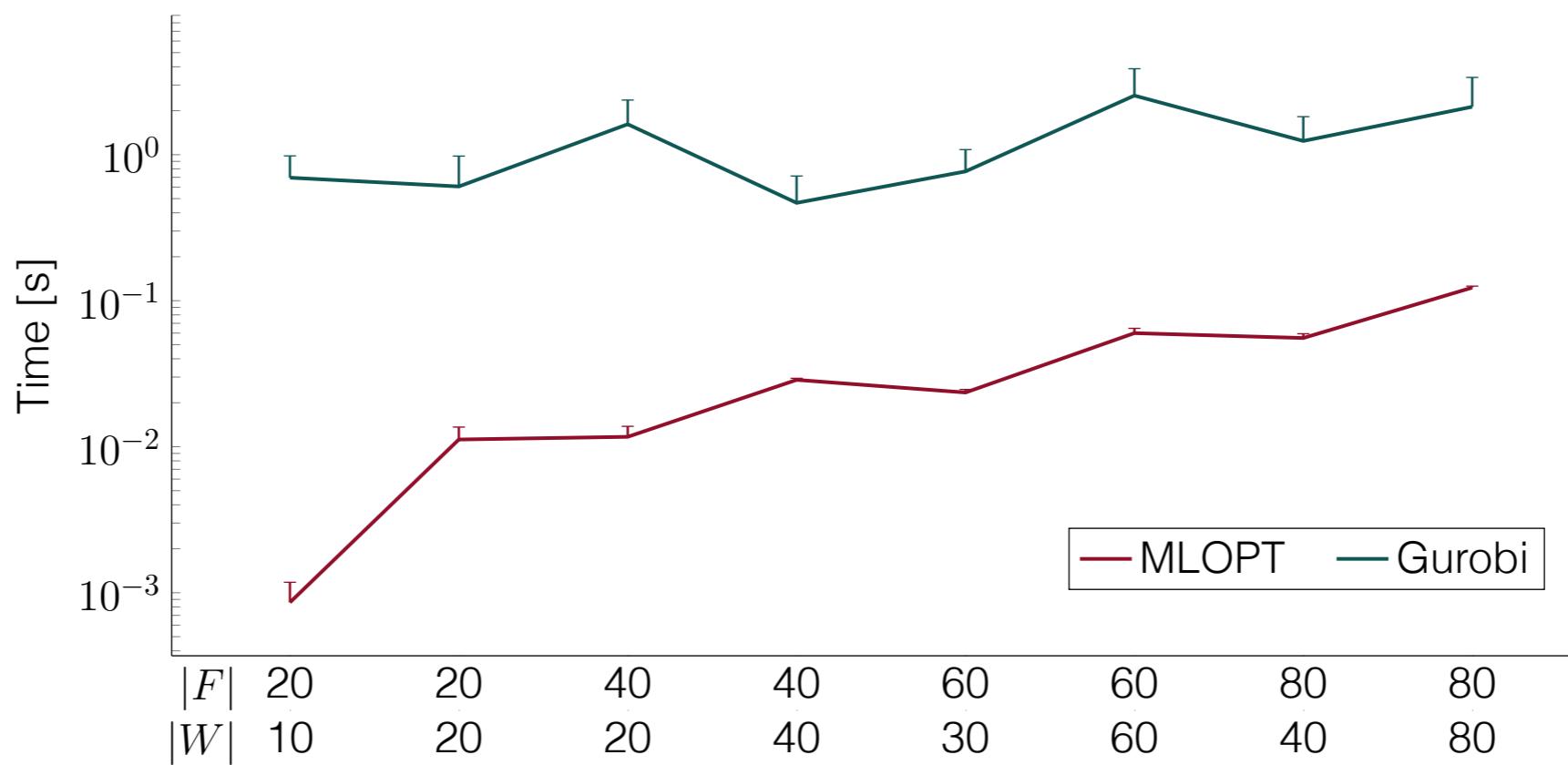
# Facility location results



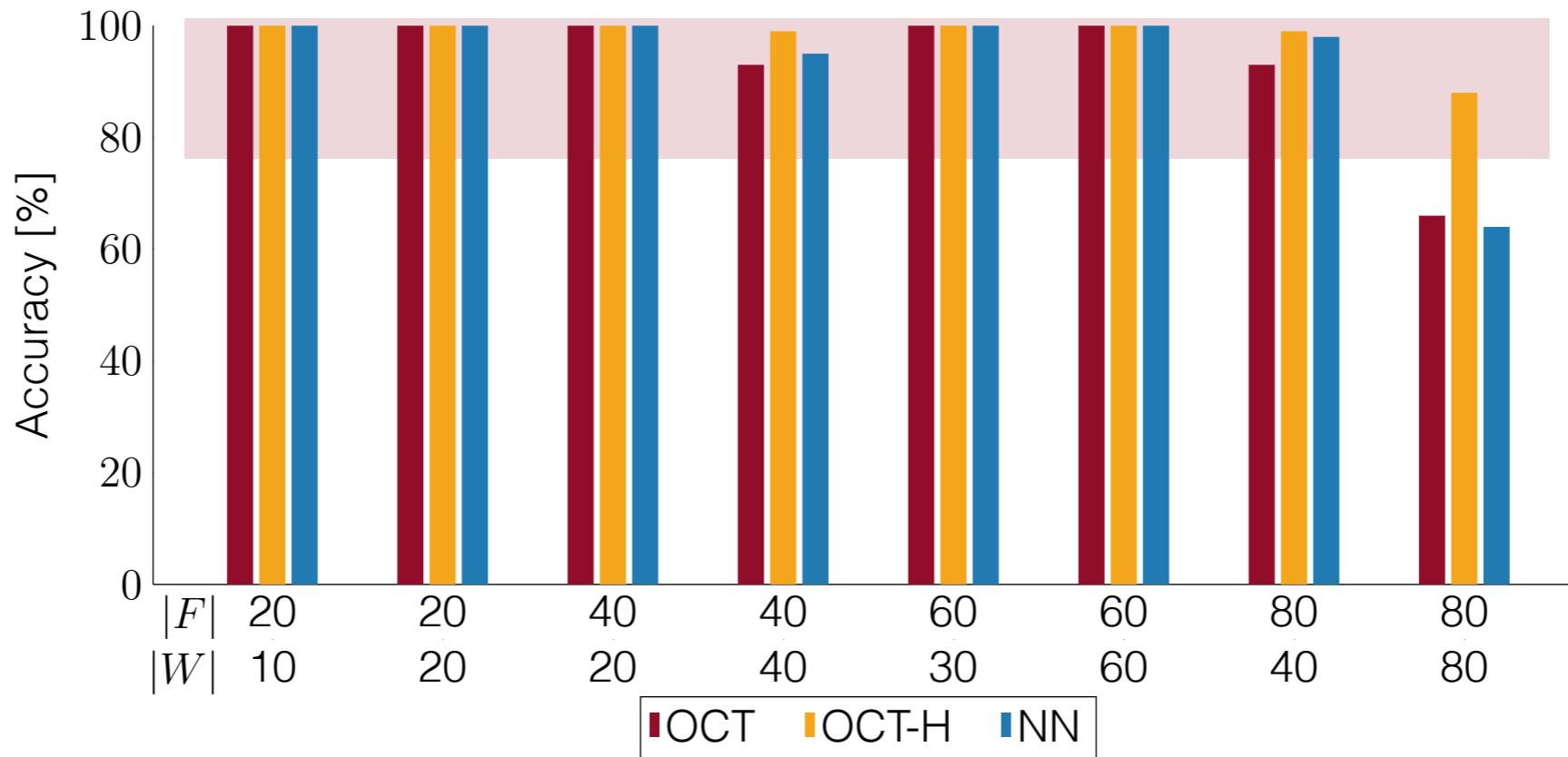
# Facility location results



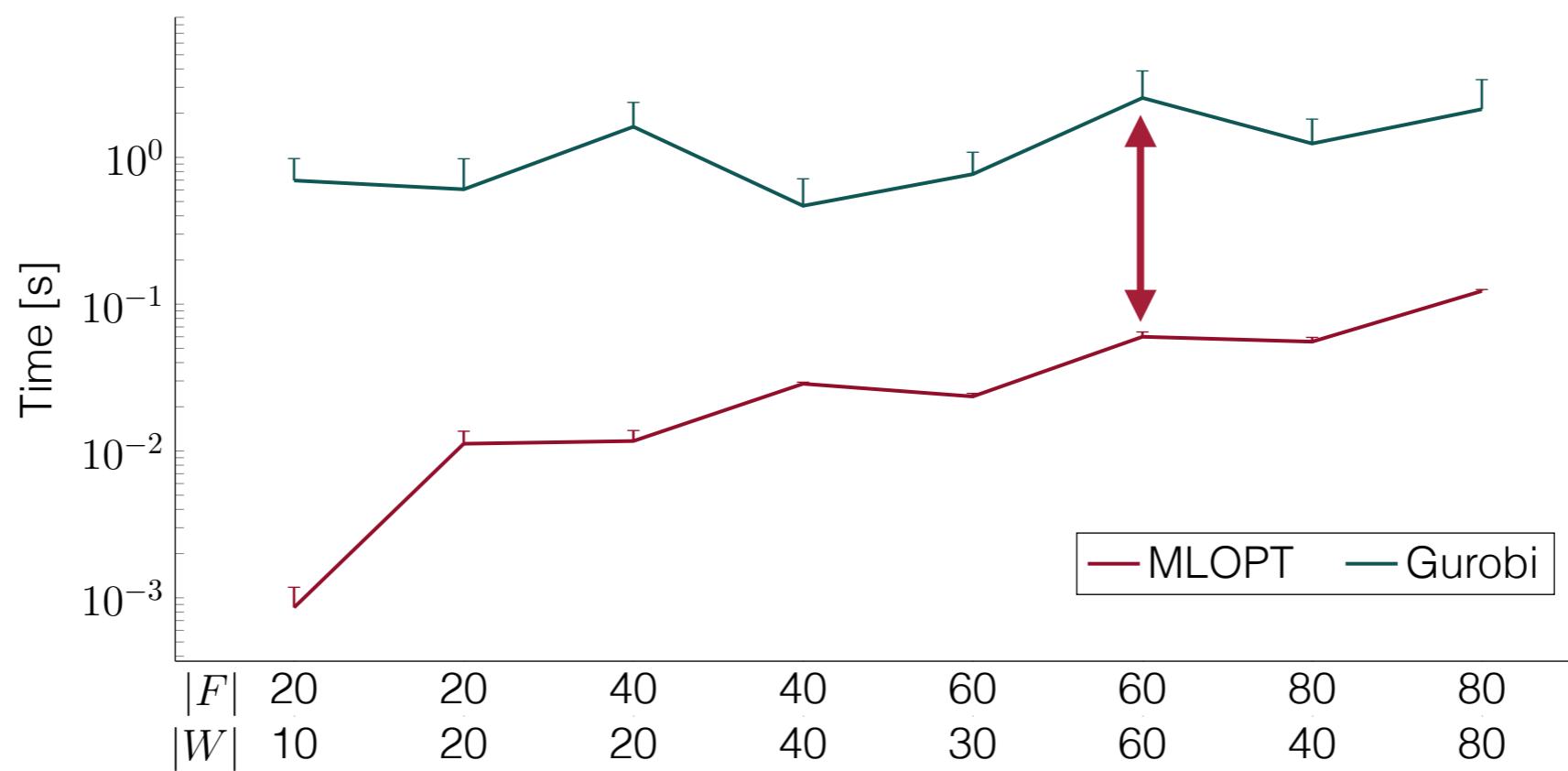
High  
accuracy



# Facility location results

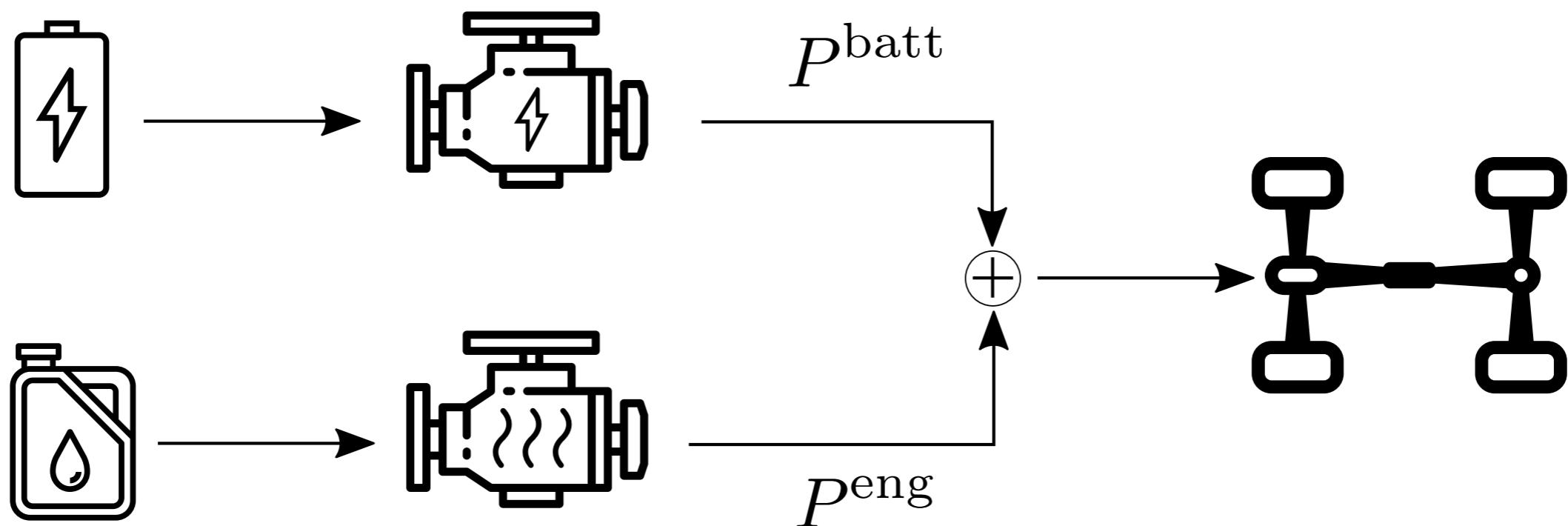


High  
accuracy

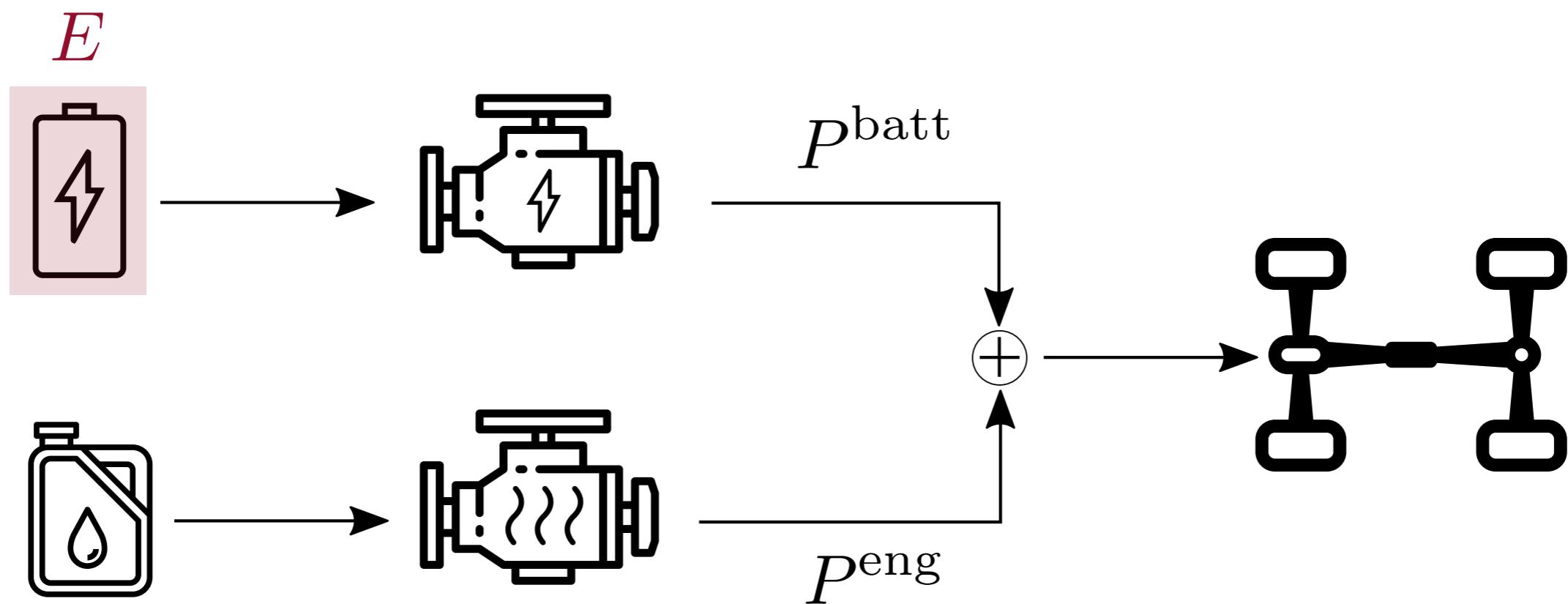


100x  
faster

# Hybrid-vehicle control



# Hybrid-vehicle control



# Hybrid-vehicle control

$$\begin{aligned} & \text{maximize} && \sum_{t=0}^{T-1} f(P_t^{\text{eng}}, z_t) + \delta(z_t - z_{t-1})_+ \\ & \text{subject to} && P_t^{\text{batt}} + P_t^{\text{eng}} \geq P_t^{\text{des}} \end{aligned}$$

$$\begin{aligned} E_{t+1} &= E_t - \tau P_t^{\text{batt}} \\ E_0 &= E_{\text{init}} \end{aligned}$$

$$z_t \in \{0, 1\}$$

# Hybrid-vehicle control

maximize 
$$\sum_{t=0}^{T-1} f(P_t^{\text{eng}}, z_t) + \delta(z_t - z_{t-1}) +$$

subject to 
$$P_t^{\text{batt}} + P_t^{\text{eng}} \geq P_t^{\text{des}}$$
 Desired Power

$$E_{t+1} = E_t - \tau P_t^{\text{batt}}$$
$$E_0 = E_{\text{init}}$$

$$z_t \in \{0, 1\}$$

# Hybrid-vehicle control

maximize 
$$\sum_{t=0}^{T-1} f(P_t^{\text{eng}}, z_t) + \delta(z_t - z_{t-1})_+$$

subject to  $P_t^{\text{batt}} + P_t^{\text{eng}} \geq P_t^{\text{des}}$  Desired Power

$$\begin{aligned} E_{t+1} &= E_t - \tau P_t^{\text{batt}} \\ E_0 &= E_{\text{init}} \end{aligned}$$
 Battery Dynamics

$$z_t \in \{0, 1\}$$

# Hybrid-vehicle control

maximize 
$$\sum_{t=0}^{T-1} f(P_t^{\text{eng}}, z_t) + \delta(z_t - z_{t-1})_+$$

subject to  $P_t^{\text{batt}} + P_t^{\text{eng}} \geq P_t^{\text{des}}$  Desired Power

$$\begin{aligned} E_{t+1} &= E_t - \tau P_t^{\text{batt}} \\ E_0 &= E_{\text{init}} \end{aligned}$$
 Battery Dynamics

$$z_t \in \{0, 1\}$$
 Engine Switch

# Hybrid-vehicle control

maximize  $\sum_{t=0}^{T-1} f(P_t^{\text{eng}}, z_t) + \delta(z_t - z_{t-1}) +$

subject to  $P_t^{\text{batt}} + P_t^{\text{eng}} \geq P_t^{\text{des}}$  Desired Power

$$E_{t+1} = E_t - \tau P_t^{\text{batt}}$$

$$E_0 = E_{\text{init}}$$

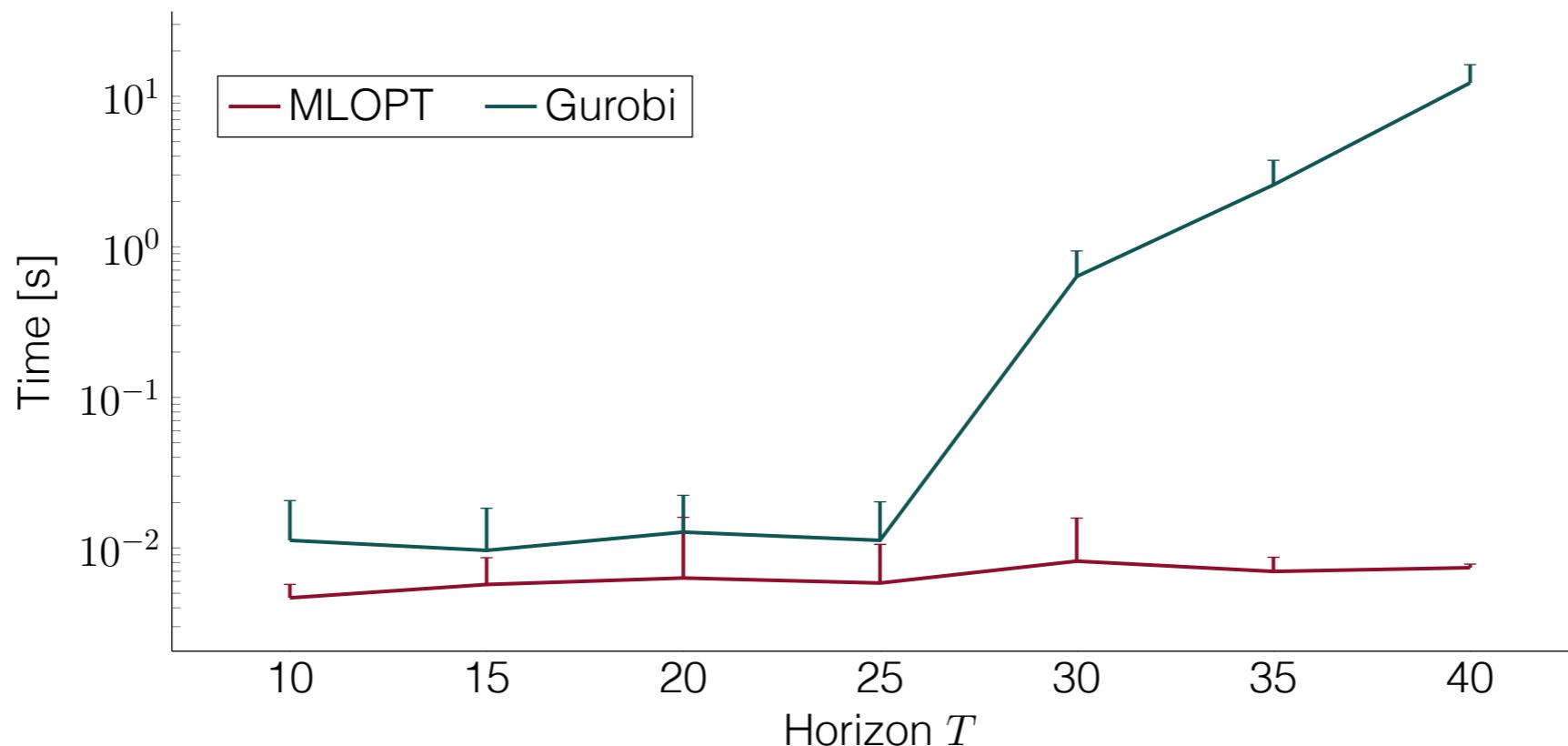
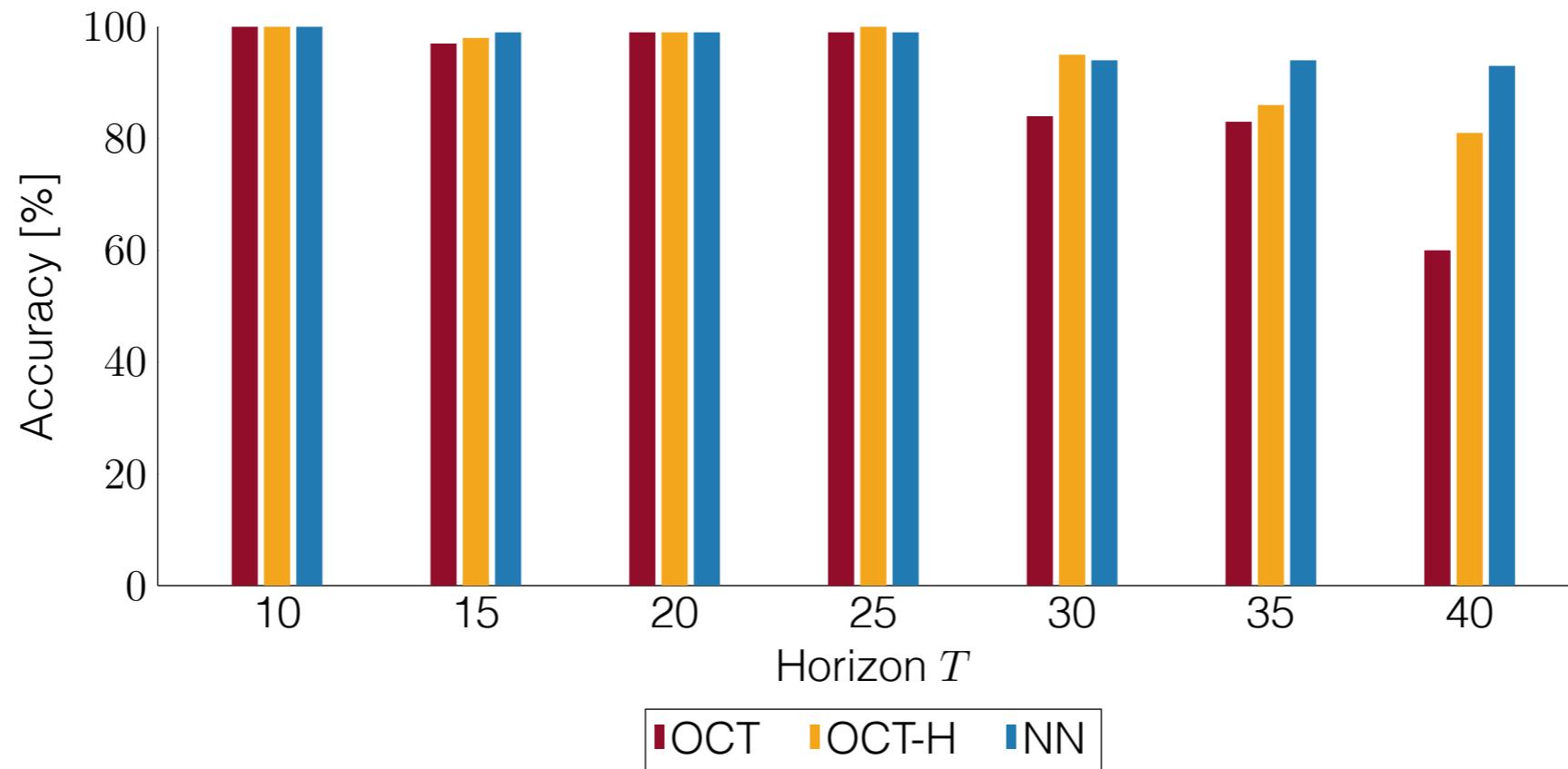
Battery  
Dynamics

$$z_t \in \{0, 1\}$$

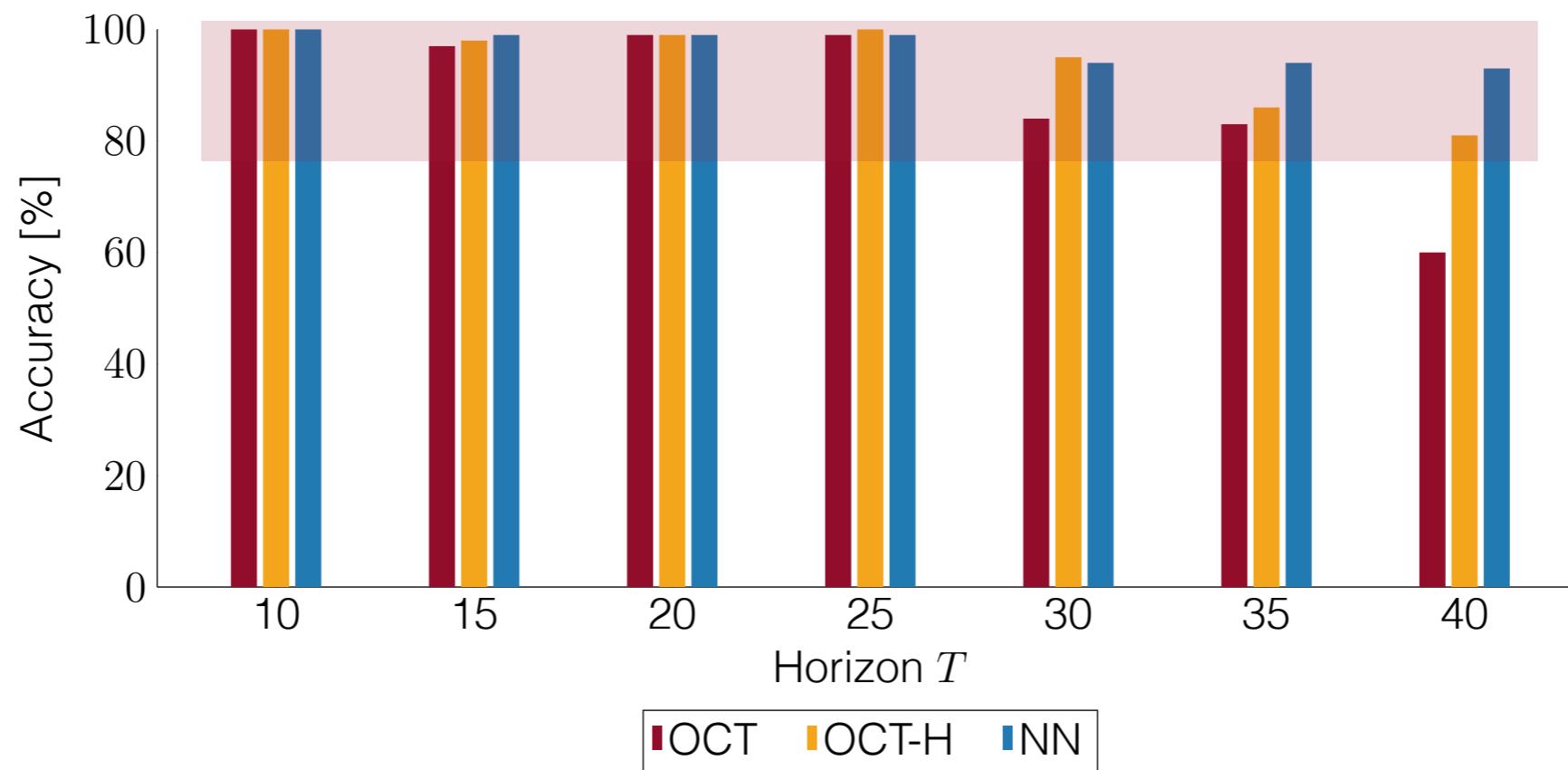
Engine Switch

Parameters

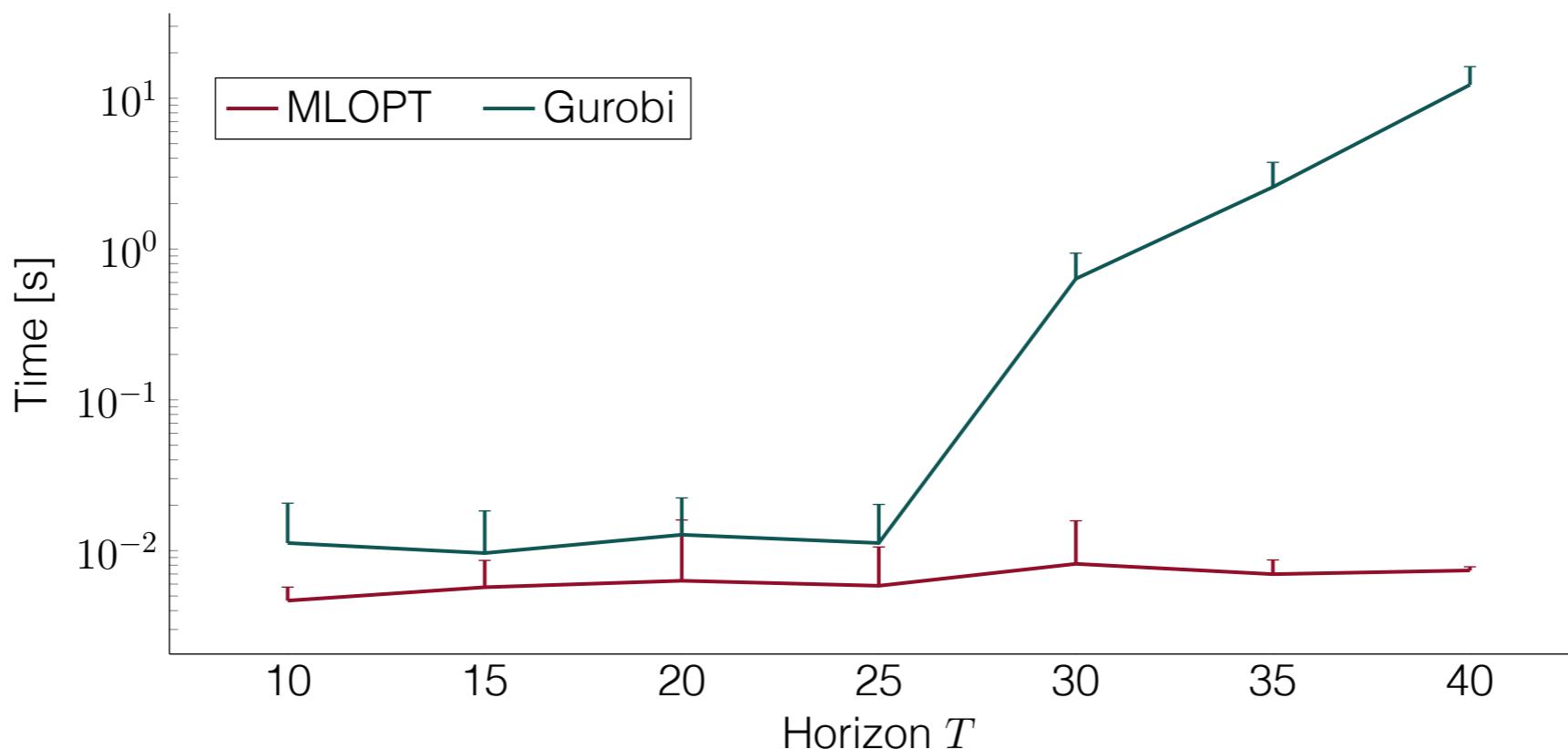
# Model-predictive control results



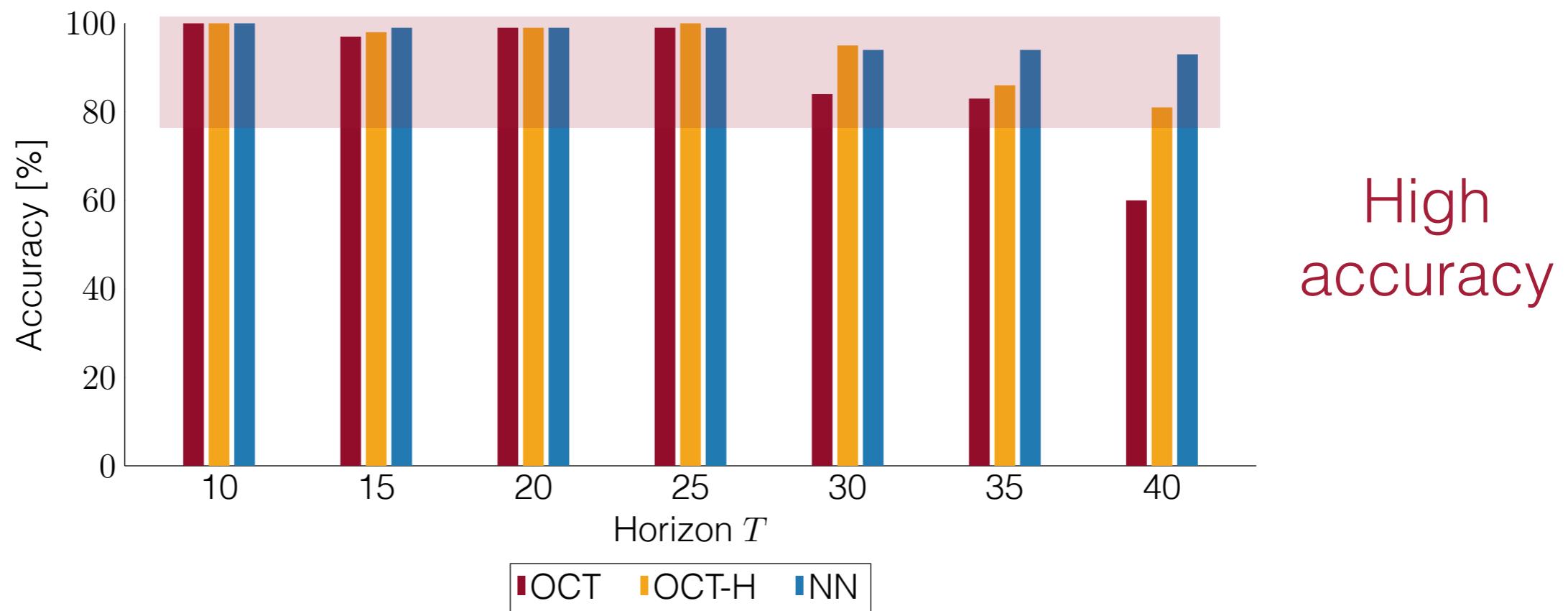
# Model-predictive control results



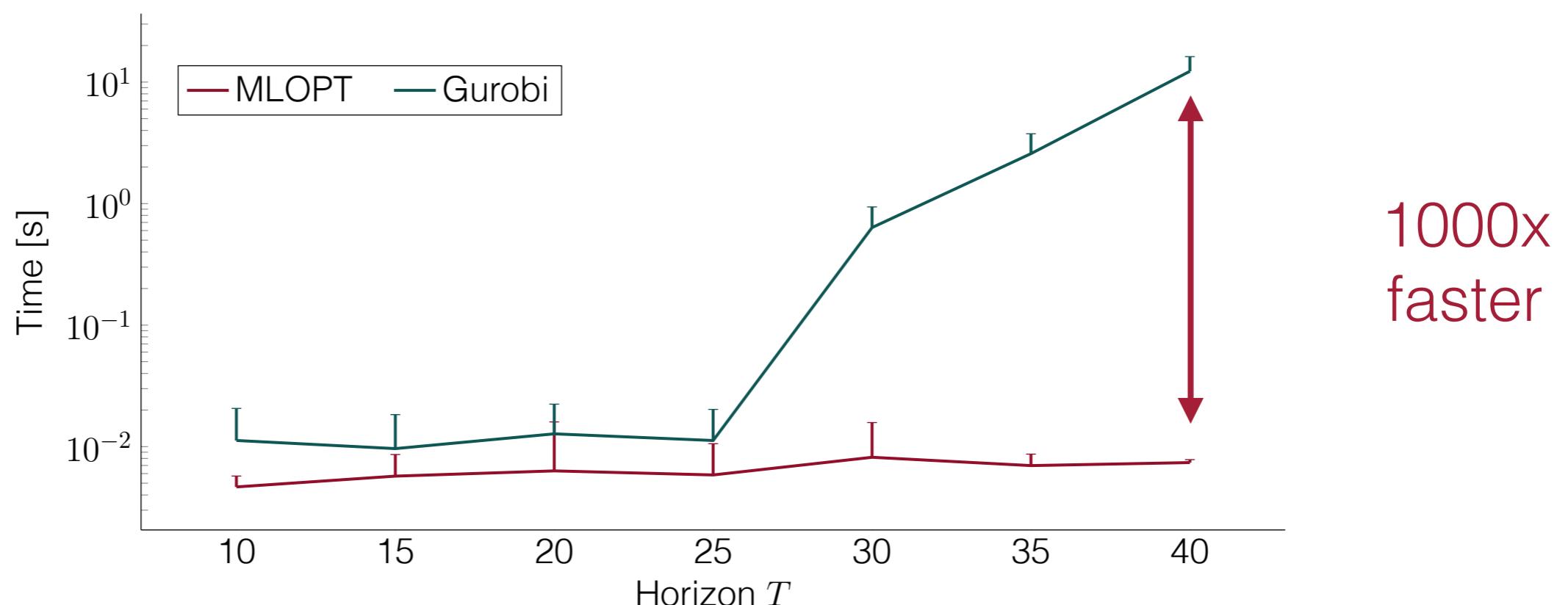
High  
accuracy



# Model-predictive control results



High accuracy



1000x faster

# Conclusions

Optimization  
is  
Parametric



Understand  
it using ML

Solve  
it really fast

# Backup

# Good-Turing estimator

Approximation

$$GT = \frac{N_1}{N} \approx \mathbf{P}(\theta_{N+1} \in \mathbf{R}^p \mid s(\theta_{N+1}) \notin \mathcal{S}(\Theta_N))$$

Bound

$$\mathbf{P}(\theta_{N+1} \in \mathbf{R}^p \mid s(\theta_{N+1}) \notin \mathcal{S}(\Theta_N)) \leq GT + c\sqrt{(1/N) \ln(3/\beta)}$$

# Sampling scheme

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## Algorithm 1 Strategies exploration

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```
1: given  $\epsilon, \beta, \Theta = \emptyset, \mathcal{S} = \emptyset, u = \infty$ 
2: for  $k = 1, \dots, \text{do}$ 
3:   Sample  $\theta_k$  and compute  $s(\theta_k)$             $\triangleright$  Sample parameter and strategy.
4:    $\Theta \leftarrow \Theta \cup \{\theta_k\}$                    $\triangleright$  Update set of samples.
5:   if  $s(\theta_k) \notin \mathcal{S}$  then
6:      $\mathcal{S} \leftarrow \mathcal{S} \cup \{s(\theta_k)\}$        $\triangleright$  Update strategy set if new strategy found
7:   end if
8:   if  $G + c\sqrt{(1/k) \ln(3/\beta)} \leq \epsilon$  then     $\triangleright$  Break if bound less than  $\epsilon$ 
9:     break
10:  end if
11: end for
12: return  $k, \Theta, \mathcal{S}$ 
```

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